Course Objectives:

- Understand the internal storage structures using different file and indexing techniques which will help in physical DB design.
- Learn the fundamentals of data models and to conceptualize and depict a database system using ER diagram.
- Know the fundamental concepts of transaction processing- concurrency control techniques and recovery procedure.
- Possess the knowledge about the emerging trends in the area of distributed DB- OO DB- Data mining and Data Warehousing and XML.

Learning Outcomes:

On successful completion of this module, the student should:

- Have gained knowledge and understanding of what is involved in the design of a database.
- Have gained knowledge and understanding of the models used for structuring data in database systems.
- Be able to implement a database and report on the process.
- Be able to query a database.

UNIT-I Introduction and Conceptual Modeling

Introduction to File and Database systems- Database system structure –Introduction and concept Modeling-Database user Data Models – Introduction to Network and Hierarchical Models – ER model – Relational Model – Relational Algebra and Calculus.

UNIT-II Relational Model

SQL – Data definition- Queries in SQL- Updates- Views – Integrity and Security – Relational Database design-Relational Models-Design issues – Functional dependences and Normalization for Relational Databases (up to BCNF).

UNIT- III Data Storage and Query Processing

Record storage and Primary file organization- Secondary storage Devices- Operations on Files- Heap File-Sorted Files- Hashing Techniques – Index Structure for files –Different types of Indexes- B-Tree - B+Tree

– Query Processing. Record storage and Primary file organization- Secondary storage Devices- Operations on Files- Heap File- Sorted Files- Hashing Techniques – Index Structure for files –Different types of Indexes- B-Tree - B+Tree – Query Processing.

UNIT- IV Transaction Management

Transaction Processing – Introduction- Need for Concurrency control- Desirable properties of Transaction- Schedule and Recoverability- Serializability and Schedules – Concurrency Control – Types of Locks- Two Phases locking- Deadlock- Time stamp based concurrency control – Recovery Techniques – Concepts- Immediate Update- Deferred Update - Shadow Paging.

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UNIT- V **Current Trends**

Object Oriented Databases – Need for Complex Data types- OO data Model- Nested relations- Complex Types- Inheritance Reference Types - Distributed databases-Homogenous and Heterogenous- Distributed data Storage – XML – Structure of XML- Data-XML Document- Schema- Querying and Transformation. - Data Mining and Data Warehousing.

Total Hours: 45 +15 =60

Text Books:

- 1. Abraham Silberschatz, Henry F.Korth and S.Sudarshan" Database System Concepts", McGraw-Hill, 6th edition,2011.
- 2. Raghuramakrishnan and Johannes Gehrke, "Database Management Systems", 3rd Edition, Mc Graw Hill, 2002
- 3. Ramesh Elmasri, Shamkant B.Navathe, "Fundamentals of Database Systems", 5th Edition, Pearson Education, 2008

References:

- 1. Hector Garcia-Molina, Jeffrey D.Ullman and Jennifer Widom" Database System Implementation" Pearson Education, 2nd Edition, 2013
- 2. Peter Rob and Corlos Coronel "Database System, Design Implementation and Management",

Thompson Learning Course Technology, 11th Edition,2014.

Websites:

- 1. http://www.tutorialized.com/tutorial/DB2-Tutorial/
- 2. http://www.techtutorials.info/datadb2.html
- 3. http://www.firstsql.com/tutor.htm
- 4. http://sqlzoo.net/



KARPAGAM ACADEMY OF HIGHER EDUCATION FACULTY OF ENGINEERING DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

15BECS401

Database Management Systems

Lecture plan

| | UN | IIT I | | |
|---|---|-------|--------------------|----|
| | Introduction and C | | | |
| 1 | Introduction to File and Database systems- | 1 | T(1) Page no.3 | |
| 2 | Database system structure | 1 | T(1) Page no.28-29 | |
| 3 | Introduction and concept | 1 | T(1) Page no.26-28 | |
| 4 | Modeling-Database user Data Models | 1 | | BB |
| 5 | Introduction to Network and Hierarchical Models | 1 | T(1) Page no.43-47 | |
| 6 | ER model. | 1 | T(1) Page no.55-56 | |
| 7 | Relational Model | | T(1) Page no.60-65 | |
| 8 | Tutorial Hour - Introduction to Network and Hierarchical Models | 1 | | |

| | Relational Algebra and Calculus | 1 | T(1) Page no.65-67 | | |
|----|--|--------|------------------------|----------|--|
| 9 | | | | | |
| | UN | IIT-II | | | |
| | Relation | | | | |
| 10 | SQL – Data definition- | 1 | T(1) Page no 68-70 | | |
| 11 | Queries in SQL- Updates- Views | 1 | T(1) Page no.76-78 | | |
| 12 | Integrity and Security | 1 | T(1) Page no.158-168 | | |
| 13 | Tutorial Hour - Queries in SQL- Updates- Views | 1 | | BB | |
| 14 | Relational Database design. | 1 | R(1) page no 3.25-3.32 | & PPT | |
| 15 | Relational Models-Design issues | 1 | T(1) Page no.193-199 | | |
| 16 | Functional dependences and Normalization for Relational Databases (up to BCNF) | 1 | R(1) page no 4.13-4.18 | | |
| 17 | Tutorial Hour - Relational Models-Design issues | 1 | | | |

| | UNIT –III | | | |
|----|---|------------|-----------------------|---|
| | Data Storage and (| Query Proc | cessing | |
| | | | T(1) Page no.245-251 | 9 |
| 18 | Record storage and Primary file organization | 1 | R(1) page no 5.2-5.7 | |
| | Secondary storage Devices- Operations on File | | T(1) Page no.252-255 | |
| 19 | | 1 | R(1) page no 5.8-5.10 | |

| | Heap File-Sorted Files- Hashing Techniques | 1 | T(1) Page no.256-261 | | |
|----|--|---------|---|-----|--|
| 20 | | - | R(1) page no 5.10-5.20 | | |
| 21 | Tutorial Hour - Secondary storage Devices- Operations on File | 1 | | | |
| 22 | Index Structure for files –Different types of Indexes | 1 | R(1) page no 5.21-5.23, T(1) Page no.266-267 | | |
| 23 | B-Tree - B+Tree Query Processing. | 1 | R(1) page no 6.2-6.8 | | |
| 24 | Record storage and Primary file organization- | 1 | R(1) page no 6.8-6.10 T(1) Page no.282-283 | | |
| 25 | Tutorial Hour - B-Tree - B+Tree Query Processing | 1 | | BB | |
| 26 | Secondary storage Devices- Operations on File | 1 | R(1) page no 6.21-6.29 T(1) Page no.288-296 | РРТ | |
| 27 | Heap File- Sorted Files- Hashing Techniques | 1 | R(1) page no 6.31-6.33 | _ | |
| 28 | Index Structure for files –Different types of Indexes- | 1 | R(1) page no 6.33-6.34 | | |
| 29 | Tutorial Hour - Secondary storage Devices- Operations on Files | 1 | | | |
| | UNIT | IV | | | |
| | Transaction | Manager | ment | | |
| | Transaction Processing- Introduction- Need | | R(1) page no 7.2-7.9 | - | |
| 30 | for Concurrency control | 1 | T(1) Page no.319-324 | | |
| 31 | Desirable properties of Transaction- Schedule and Recoverability | 1 | | | |
| | - Serializability and Schedules – Concurrency Control | | R(1) page no 7.11-7.19 | | |
| 32 | | 1 | T(1) Page no.327-339 | | |
| 33 | Tutorial Hour - Test is to be conducted for the topic 4.1-4.4 | | | | |

| 34 | Types of Locks- Two Phases locking- Deadlock | 1 | R(1) page no 7.20-7.23 | |
|----|--|--------|--|----------|
| 35 | Time stamp based concurrency control – Recovery Techniques – | 1 | R(1) page no 8.2-8.7 T(1) Page no.373-381 | ВВ |
| 36 | Concepts- Immediate Update- Deferred Update - Shadow Paging. | 1 | R(1) page no 8.8-8.11 | & PPT |
| 37 | Tutorial Hour - Concepts- Immediate Update- Deferred Update - Shadow Paging. | 1 | | |
| | UNIT | -V | | |
| | Current | Trends | | |
| 53 | Object Oriented Databases – Need for Complex | 1 | R(1) page no 8.24-8.27 | |
| 54 | Data types- OO data Model | | R(1) page no 8.29-8.35 | BB |
| | | 1 | T(1) Page no.419-121 | & PPT |
| 55 | Nested relations- Complex Types | 1 | T(1) Page no.429-431 R(1) page no 8.35, | |
| 56 | Tutorial Hour - Nested relations- Complex Types | 1 | | |
| 57 | Inheritance Reference Types - Distributed databases | 1 | R(1) page no 9.12-9.15 R(1) page no 9.21-9.22 | |
| 58 | Homogenous and Heterogenous- Distributed data Storage – | 1 | T(1) Page no.456-461 | |
| 59 | XML – Structure of XML- Data- XML Document | 1 | R(1) page no 9.27-9.35 | PPT |
| 60 | Tutorial Hour - Nested relations- Complex Types | 1 | | BB |

| 61 | Schema- Querying and Transformation | 1 | R(1) page no 9.36-9.37 | & |
|----|---|---|--|-----|
| 62 | Data Mining and Data Warehousing | 1 | R(1) page no 9.37-9.38 T(1) Page no.466-467 | PPT |
| 6 | Tutorial Hour - Querying and Transformation. | 1 | | |

Hours allocated

Lecture

hours:45

Tutorial

hours:15

Total

hours:60

Text Books:

- 4. Abraham Silberschatz, Henry F.Korth and S.Sudarshan" Database System Concepts", McGraw-Hill, 6th edition,2011.
- 5. Raghuramakrishnan and Johannes Gehrke, "Database Management Systems", 3rd Edition, Mc Graw Hill, 2002
- 6. Ramesh Elmasri, Shamkant B.Navathe, "Fundamentals of Database Systems", 5th Edition, Pearson Education, 2008

References:

- Hector Garcia–Molina, Jeffrey D.Ullman and Jennifer Widom" Database System Implementation" Pearson Education, 2nd Edition,2013
 Peter Rob and Corlos Coronel "Database System, Design Implementation and
- Peter Rob and Corlos Coronel "Database System, Design Implementation and Management", Thompson Learning Course Technology, 11th Edition,2014

UNIT I

INTRODUCTION AND CONCEPTUAL MODELING

Introduction to File and Database systems- Database system structure –Introduction and concept Modeling-Database user Data Models – Introduction to Network and Hierarchical Models – ER model – Relational Model – Relational Algebra and Calculus

1. INTRODUCTION

Data: Known facts that can be recorded that have implicit meaning. E.g. Student roll no, names, address etc

Database: collection of inter-related data organized meaningfully for a specific purpose.

DBMS: DBMS is a collection of interrelated data and a set of program to access those data. The primary goal of a DBMS is to provide a way to store and retrieve database information that is both *convenient* and *efficient*.

Database System: Database and DBMS collectively known as database system. INTRODUCTION TO FILE AND DATABASE SYSTEMS:

Database Applications

- Banking: all transactions
- Airlines: reservations, schedules
- Universities: registration, grades
- Sales: customers, products, purchases
- Online retailers: order tracking, customized recommendations
- Manufacturing: production, inventory, orders, supply chain
- Human resources: employee records, salaries, tax deductions
- Credit card transactions
- Telecommunications & Finance

2. PURPOSE OF DATABASE SYSTEMS

Drawbacks of Conventional File Processing System

i. Data redundancy and inconsistency

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Since the files and application programs are created by different programmers over a long period of time, the files have different formats and the programs may be written in several

programming language. The same piece of information may be duplicated in several files.

For Example: The address and phone number of particular customer may appear in a file that consists of personal information and in saving account records file also. This redundancy

leads to data consistency that is, the various copies of the same data may no longer agree.

For example: a changed customer address may be reflected in personal information file, but not in saving account records file.

ii. Difficulty in accessing data

Conventional file processing environments do not allow needed data to be retrieved in a convenient and efficient manner.

For Example: Suppose that bank officer needs to find out the names of all customers who live within the city's 411027 zip code. The bank officer has now two choices: Either get the

list of customers and extract the needed information manually, or ask the data processing department to have a system programmer write the necessary application program. Both alternatives are unsatisfactory.

iii. Data isolation

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Since, data is scattered in various files, and files may be in different formats, it is difficult to write new application programs to retrieve appropriate data.

iv. Concurrent access anomalies

In order to improve the overall performance of the system and obtain a faster response time many systems allow multiple users to update the data simultaneously. In such environment,

interaction of concurrent updates may results in inconsistent data.

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For Example: Consider bank account A, with \$500.If two customers with draw funds (say \$50 and \$100 resp) from account A at the same time, the result of the concurrent executions

\$400, rather than \$350. In order to guard against this possibility, some form of supervision must be maintained in the system.

v. Atomicity Problem

System failure will lead to atomicity problem.

For Example: Failure during transfer of fund from system A to A. It will be debited from A but not credited to B leading to wrong transaction.

vi. Concurrent Access Anomalies

In order to improve the overall performance of the system and obtain a faster response time many systems allow multiple users to update the data simultaneously. In such environment,

interaction of concurrent updates may result in inconsistent data.

For Example: Consider bank account *A*, containing \$500. If two customers withdraw funds say \$50 and \$100 respectively) from account *A* at about the same time, the result of the

concurrent executions may leave the account in an incorrect (or inconsistent) state. Balance

will be \$400 instead of \$350. To protect against this possibility, the system must maintain some form of supervision.

vii. Security problems

Not every user of the database system should be able to access all the data. System should be protected using proper security.

For Example: In a banking system, pay roll personnel should be only given authority to see the part of the database that has information about the various bank employees. They do not

need access to information about customer accounts.

Since application programs added to the system in an ad-hoc manner, it is difficult to enforce such security constraints.

viii. Integrity problems

The data values stored in the database must satisfy certain types of consistency constrains.

For Example: The balance of a bank account may never fall below a prescribed amount (say \$100). These constraints are enforced in the system by adding appropriate code in the various

application programs.

Advantages of Database

Data base is a way to consolidate and control the operational data centrally. It is a better way to control the operational data. The advantages of having a centralized control of data are:

i. Redundancy can be reduced

In non-database systems, each application or department has its own private files resulting in considerable amount of redundancy of the stored data. Thus storage space is wasted. By having a centralized database most of this can be avoided.

ii. Inconsistency can be avoided

When the same data is duplicated and changes are made at one side, which is not propagated to the other site, it gives rise to inconsistency. Then the two entries regarding the same data will not agree. So, if the redundancy is removed, chances of having inconsistent data are also removed.

iii. The data can be shared

The data stored from one application, can be used for another application. Thus, the data of database stored for one application can be shared with new applications.

iv. Standards can be enforced

With central control of the database, the DBA can ensure that all applicable standards are observed in the representation of the data.

v. Security can be enforced

DBA can define the access paths for accessing the data stored in database and he can define authorization checks whenever access to sensitive data is attempted.

vi. Integrity can be maintained

Integrity means that the data in the database is accurate. Centralized control of the data helps in permitting the administrator to define integrity constraints to the data in the database.

3. VIEW OF DATA

A major purpose of a database system is to provide users with an abstract view of the data. That is, the system hides certain details of how the data are stored and maintained.

Data abstraction

The Complexity is hidden from the users through several level of abstraction. There are three levels of data abstraction:

- Physical level: It is the lowest level of abstraction that describes <u>how the data are actually stored</u>. The physical level describes complex low-level data structures in details.
- ii. **Logical level:** It is the next higher level of abstraction that describes <u>what data are stored in the</u> <u>database</u> and what relationships exist among those data.
- iii. View level: It is the highest level of abstraction that describes only part of the entire database.

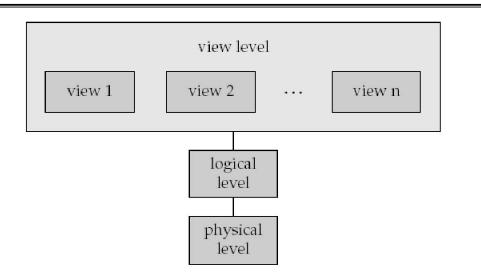


Figure The three levels of data abstraction

Data Independence

The ability to modify a scheme definition in one level without affecting a scheme definition in the next higher level is called data independence. There are two levels of data independence:

1. Physical data independence is the ability to modify the physical scheme without causing application programs to be rewritten. Modifications at the physical level are occasionally necessary in order to improve performance.

2. Logical data independence is the ability to modify the conceptual scheme without causing application programs to be rewritten. Modifications at the conceptual level are necessary whenever the logical structure of the database is altered.

Logical data independence is more difficult to achieve than physical data independence since application programs are heavily dependent on the logical structure of the data they access.

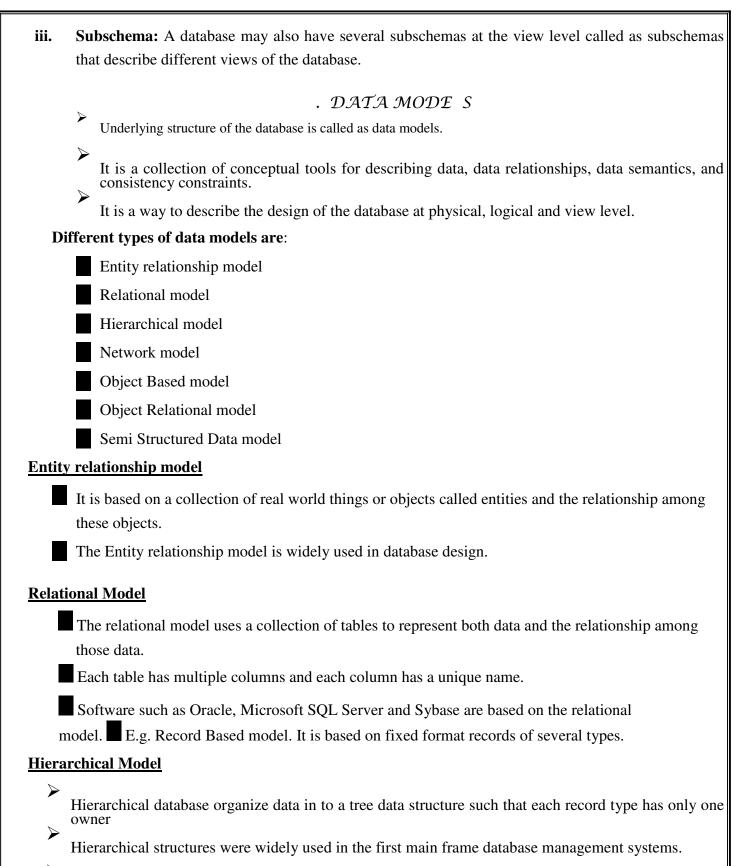
Instances and schemas

Database change over times as information is inserted and deleted. The collection of information stored in the database at a particular moment is called an **instance** of the database.

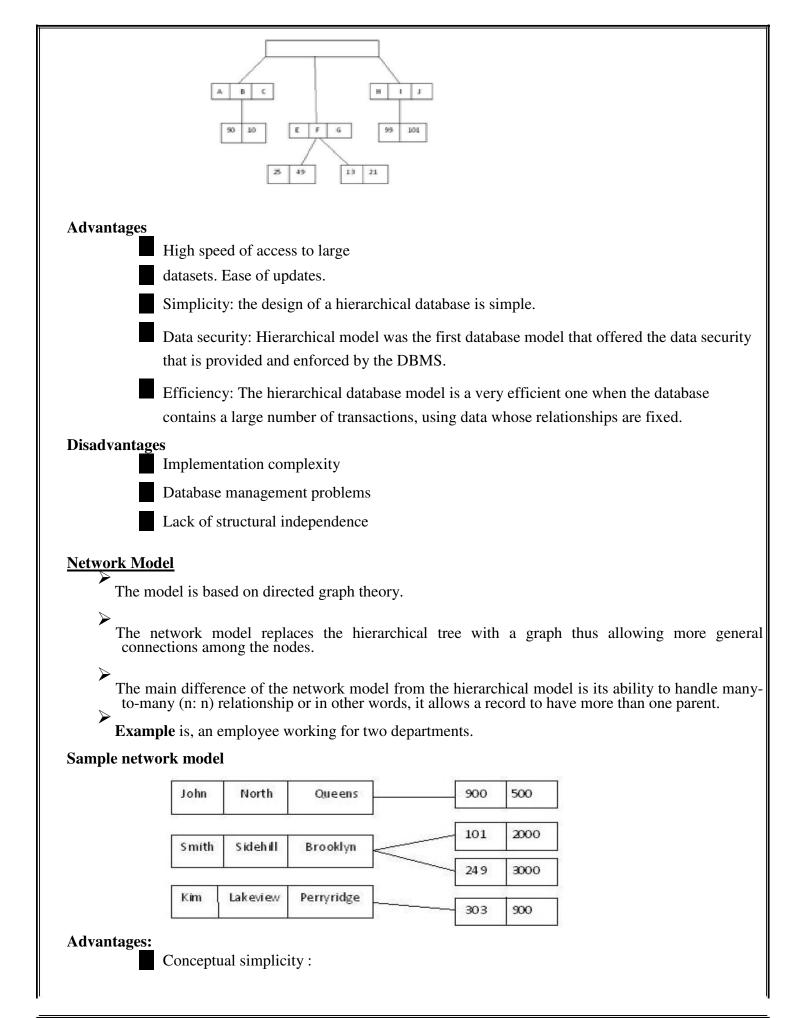
The overall design of the database is called the database schema.

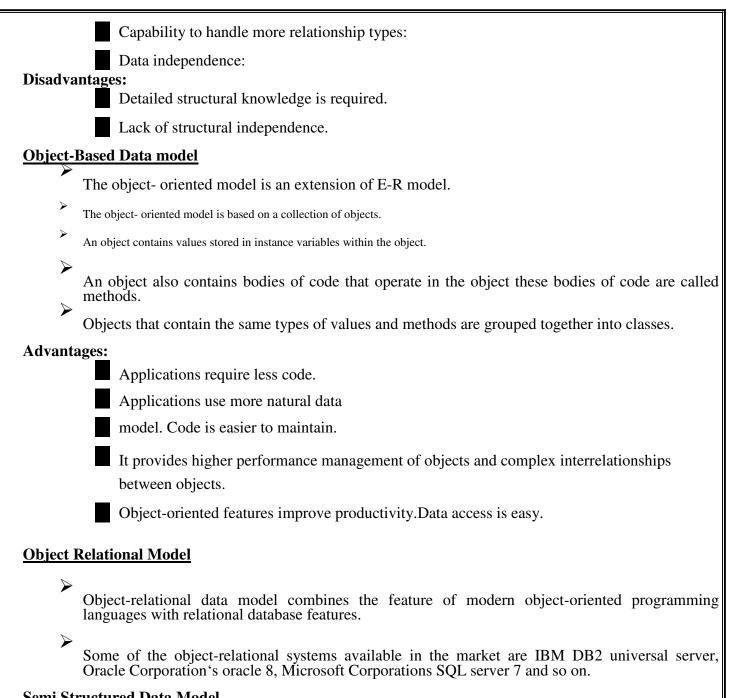
Types of database schemas

- i. **Physical schema**: It describes the database design at the physical level.
- **ii.** Logical schema: It describes the database design at the physical level.



Links are possible vertically but not horizontally or diagonally.





Semi Structured Data Model

This data model allows the individual data items of same type to have different sets of attributes.

Other data model allows a particular type of data item to have same set of attributes.

Extensible Markup Language (XML) is used to represent structured data.

5. DATABASE AN UA ES

A database system provides

A **Data Definition Language** to specify the database schema (DDL)

A **Data Manipulation Language** to express database queries and updates.

Data definition and data manipulation languages are not two separate languages but part of a single database language such as SQL language.

Data definition language

DDL specifies the database schema and some additional properties to data.

The storage structure and access methods are specified using specified using special type of DDL called s **data storage and data definition langu**age.

The data values stored in the database must satisfy certain **consistency constraints**. For example, suppose the balance on an account should not fall below \$100.

Database system concentrates on constraints that have less overload.

1. Domain Constraints:

Domain of possible value should be associated with every attributes.

E.g. integer type, character type, date/time type

Declaring attributes to a particular domain will act as a constraint on that value.

They are tested as and when values are entered in to database.

2. Referential Constraints:

In some cases there will be value that appears in one relation for a given set of attributes also appears for a certain set of attributes in some other relation. Such constraint is called Referential Constraints.

If any modification violates the constraints then the action that caused the violation should be rejected.

3. Assertions

It is a condition that database should always satisfy.

Domains and referential integrity are special form of assertion.

E.g. Every loan should have a customer whose account balance is minimum of \$1000.00 Modifications to database should not cause violation to assertion.

4. Authorization

The users are differentiated as per the access permit given to them on the different data of the database. This is known as authorization.

The most common authorizations are

i. Read authorization

Allows reading but no modification of data.

ii. Insert authorization

Allows insertion of new data but no modification of existing data.

iii. Update authorization

Allows modification but not deletion.

iv. Delete authorization

Allows deletion of data.

The users may be assigned with all, none or combination of these

types. The DDL gets some input and generates some output.

This output is placed in data dictionary which contains Meta

data. Meta data is data about data.

Data Dictionary is a special type of table which can only be accessed and updated by database system.

Database system consults the Data Dictionary before reading or modifying actual data.

Data Manipulation Language

A **data-manipulation language (DML)** is a language that helps users to access or manipulate data. A query is a statement requesting the retrieval of information.

The portion of DML that involves information retrieval is called as query language. There are basically two types of DML:

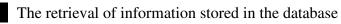
Procedural DMLs

User should specify what data are needed and how to get those data.

Declarative DMLs (nonprocedural DMLs)

User should specify *what* data are needed *without* specifying how to get those data. This is easier to learn and user than procedural DML.

Data manipulation that can be performed using DML are



The insertion of new information into the database

The deletion of information from the database

The modification of information stored in the database

. DATABASE SYSTEM ARC ITECTURE

A database system is partitioned into modules that deal with each of the responsibilities of the overall system. The functional components of a database system can be broadly divided into

- Storage Manager
- Query Processor

The database system architecture is influenced by the underlying computer

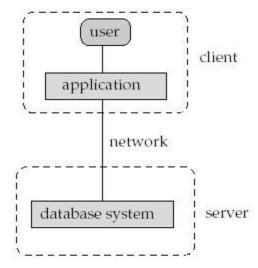
architecture. The database system can be centralized or client server.

Database systems are partitioned into two or three parts.

In **two tier architecture**, the application is partitioned into a component that resides at the client machine and invokes database functionality at the server machine through query language.

Application program interface standards like ODBC and JDBC are used for interaction between the client and the server.

Two tier architecture



In **three tier architecture**, the client machines act as a front end and do not contain any direct database calls.

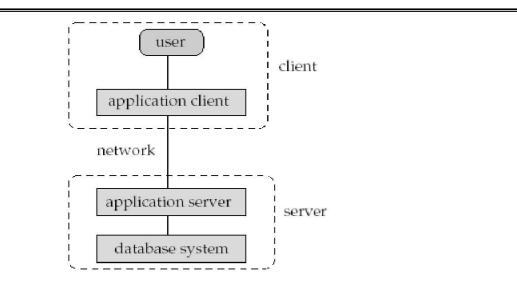
The client end communicates with the application servers through interface.

The application server interacts with database system to access data.

The business logic of application says what actions to be carried out under what condition.

Three tier is more appropriate for large applications.

Three tier architecture



Storage Manager

A storage manager is a program module that provides the interface between the low level data stored in the database and the application programs and queries submitted to the system.

The storage manager is responsible for the interaction with the file manager.

The storage manager translates the various DML statements into low-level file system commands. Thus, the storage manager is responsible for storing, retrieving, and updating data in the database.

Components of the storage manager are:

- 1. Authorization and integrity manager: It tests for satisfaction of various integrity constraints and checks the authority of users accessing the data.
- **2. Transaction manager**: It ensures that the database remains in a consistent state despite system failures, and concurrent executions proceed without conflicting.
- **3.** File manager: It manages the allocation of space on disk storage and the data structures used to represent information stored on disk.
- 4. Buffer manager: It is responsible for fetching data from disk storage into main memory and to decide what data to cache in main memory. It enables the database to handle data sizes that are much larger than the size of the main memory. The storage manager implements several data structures as part of physical system implementation.
 - i. **Data files**: which store the database itself.
 - ii. **Data dictionary**: It contains metadata that is data about data. The schema of a table is an example of metadata. A database system consults the data dictionary before reading and modifying actual data.

iii. Indices: Which provide fast access to data items that hold particular values.

The Query Processor

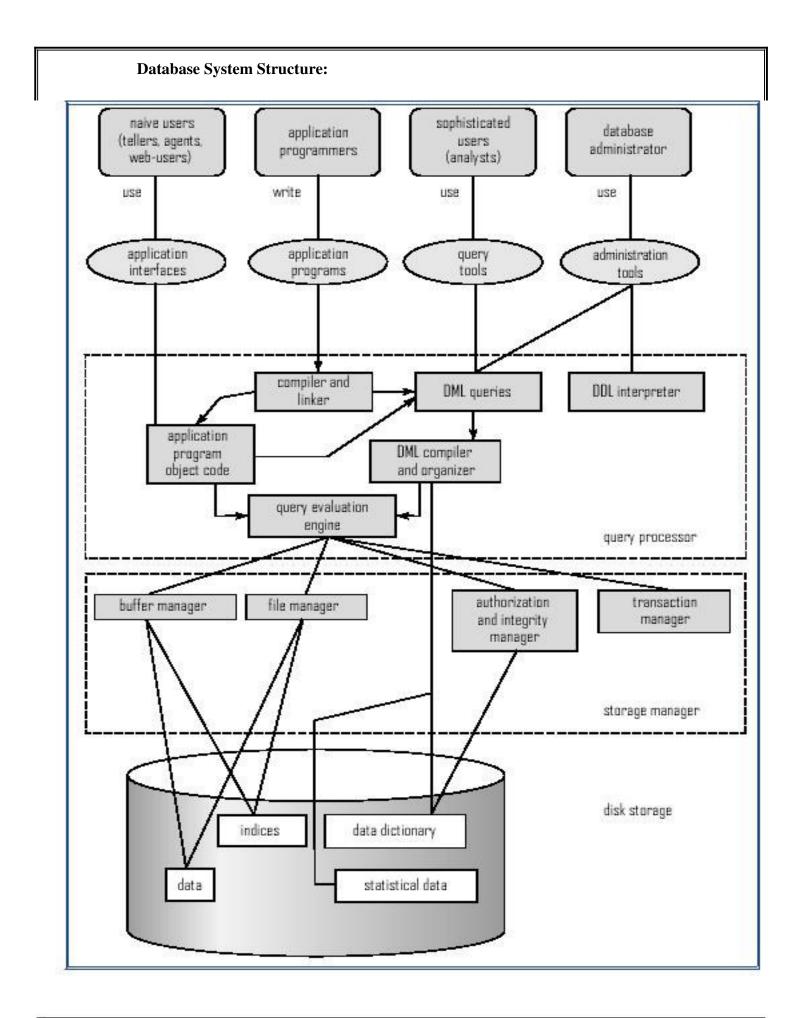
The query processor is an important part of the database system. It helps the database system to simplify and facilitate access to data. The query processor components include:

- 1. **DDL interpreter**, which interprets DDL statements and records the definitions in the data dictionary.
- **2. DML compiler**, which translates DML statements in a query language into an evaluation plan consisting of low-level instructions that the query evaluation engine understands.

A query can be translates into any number of evaluations plans that all give the same result.

The DML compiler also performs query optimization, that is, it picks up the lowest cost evaluation plan from among the alternatives.

3. Query evaluation engine, which executes low-level instructions generated by the DML compiler.



. DATABASE USERS AND ADMINISTRATOR

People who work with a database can be categorized as:

Database Users

Database administrators

7.1. DATABASE USERS

There are four types of database users, differentiated by the way they interact with the system.

1. Naive users

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Naive users interact with the system by invoking one of the application programs that have been written previously.

Naive users are typical users of form interface, where the user can fill in appropriate fields of the form.

Naive users may also simply read *reports* generated from the database.

2. Application Programmers

Application programmers are computer professionals who write application programs.

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Rapid application development (RAD) tools enable the application programmer to construct forms and reports without writing a program.

Special types of programming languages that combine control structures with data manipulation language. These languages, sometimes called *fourth-generation languages*.

3. Sophisticated users

Sophisticated users interact with the system without writing programs. Instead, they form their requests in a database query language.

They submit each such query to a **query processor** that the storage manager understands.

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Online analytical processing (OLAP) tools simplify analysis and **data mining** tools specify certain kinds of patterns in data.

4. Specialized users

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Specialized users are sophisticated users who write specialized database applications that do not fit into the traditional data-processing framework.

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The applications are computer-aided design systems, knowledge base and expert systems, systems that store data with complex data types

7.2. DATABASE ADMINISTRATORS

A person who has such central control over the system is called a **database administrator** (**DBA**).

The functions of a DBA include:

• Schema definition. The DBA creates the original database schema by executing a set of data definition statements in the DDL.

• Storage structure and access-method definition.

• Schema and physical-organization modification. The DBA carries out changes to the schema and physical organization to reflect the changing needs of the organization.

• Granting of authorization for data access. By granting different types of authorization, the database administrator can regulate which parts of the database various users can access.

Authorization information is kept in a special system structure that the database system consults whenever someone attempts to access the data in the system.

• Routine maintenance. Examples of the database administrator's routine maintenance activities are:

- 1. periodically backing up the database
- 2. Ensuring that enough free disk space
- 3. Monitoring jobs running on the database and ensuring that performance is not degraded by very expensive tasks submitted by some users.
- 4. Ensuring that performance is not degraded by very expensive tasks submitted by some users.

. ENTITY RE ATIONS IP MODE ER MODE

The E-R data model considers the real world consisting of a set of basic objects, called entities, and relationships among these objects.

The E-R data model employs three basic notions:

- 1. Entity sets
- 2. Relationship sets
- 3. Attributes

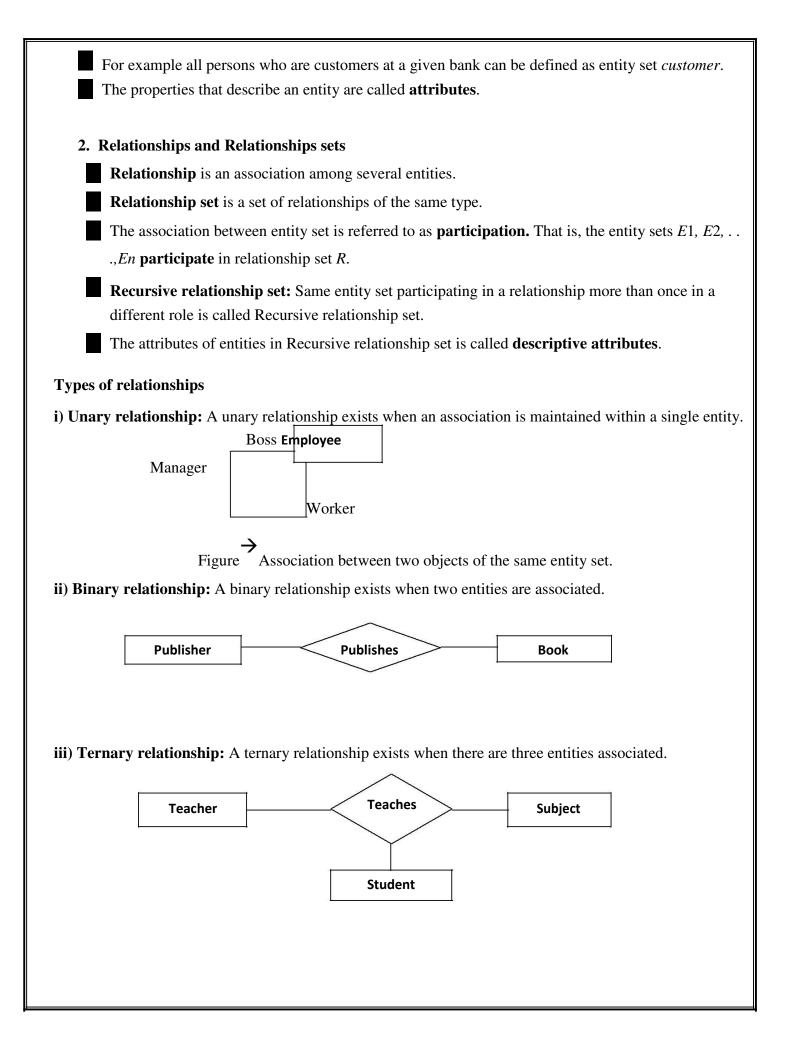
1. Entity Sets

An *entity* is _thing' or _object in the real world that is distinguishable from all other objects. For example, each person is an entity.

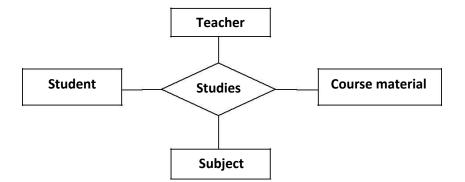
An entity has a set of properties, and the values for some set of properties may uniquely identify an entity.

For example, a customer with customer-id property with value C101 uniquely identifies that person. An entity may be concrete, such as person or a book, or it may be abstract, such as a loan, or a holiday.

An *entity set* is a set of entities of the same type that share the same properties, or attributes.



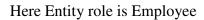
iv) Quaternary relationship: A quaternary relationship exists when there are four entities associated.



The number of entity set participating in a relationship is called **degree of the relationship set**. Binary relationship set is of degree 2; a tertiary relationship set is of degree 3.

Entity role: The function that an entity plays in a relationship is called that entity's role. A role is one end of an association.





3. Attributes

The properties that describes an entity is called attributes.

The attributes of customer entity set are customer_id, customer_name and city.

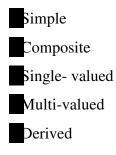
Each attributes has a set of permitted values called the domain or value set.

Each entity will have value for its attributes.

Example:

- Customer Name John
- Customer Id 321

Attributes are classified as



1) Simple attribute:

This type of attributes cannot be divided into sub parts.

Example: Age, sex, GPA

2) Composite attribute:

This type of attributes Can be subdivided.

Example: Address: street, city, state, zip

3) Single-valued attribute:

This type of attributes can have only a single value

Example: Social security number

4) Multi-valued attribute:

Multi-valued attribute Can have many values.

Example: Person may have several college degrees, phone numbers

5) Derived attribute:

Derived attribute Can be calculated or derived from other related attributes or entities.

Example: Age can be derived from D.O.B.

6) Stored attributes:

The attributes stored in a data base are called stored attributes.

An attribute takes a null value when an entity does not have a value for it.

Null values indicate the value for the particular attribute does not exists or unknown.

E.g. : 1. Middle name may not be present for a person (non existence case)

2. Apartment number may be missing or unknown.

CONSTRAINTS

An E-R enterprise schema may define certain constraints to which the contents of a database system must conform.

Three types of constraints are

- 1. Mapping cardinalities
- 2. Key constraints
- 3. Participation constraints

1. Mapping cardinalities

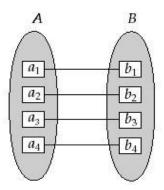
Mapping cardinalities express the number of entities to which another entity can be associated via a relationship set.

Cardinality in E-R diagram that is represented by two ways:

i) Directed line (→ ii) Undirected line ()—

There are 4 categories of cardinality.

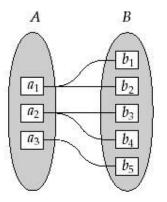
i) One to one: An entity in A is associated with at most one entity in B, and an entity in B is associated with at most one entity in A.



Example: A customer with single account at given branch is shown by one-toone relationship as given below



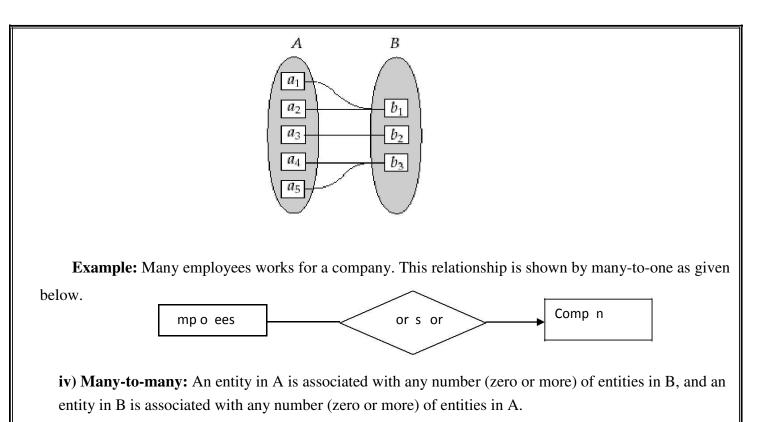
ii) One-to-many: An entity in A is associated with any number of entities (zero or more) in B. An entity in B, however, can be associated with at most one entity in A.

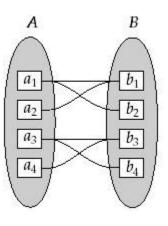


Example: A customer having two accounts at a given branch is shown by one-to-many relationship as given below.

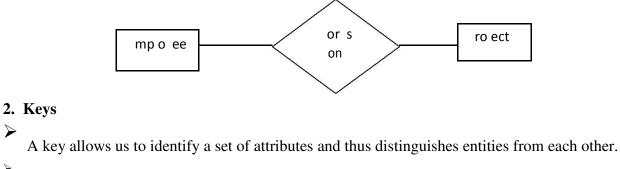


iii) **Many-to-one:** An entity in A is associated with at most one entity in B. An entity in B, however, can be associated with any number (zero or more) of entities in A.





Example: Employee works on number of projects and project is handled by number of employees. Therefore, the relationship between employee and project is many-to-many as shown below.



 \geq Keys also help uniquely identify relationships, and thus distinguish relationships from each other.

 \mathbf{F}

| Key Type | Definition |
|--------------------|------------|
| - 5 - 5 F - | |

| | Any attribute or combination of attributes that uniquely identifies a row in the table |
|---------------|---|
| | Any attribute or combination of attributes that uniquely identifies a row in the table. |
| Superkey | Example: Roll_No attribute of the entity set _student' distinguishes one student entity |
| | from another. Customer_name, Customer_id together is a Super key |
| | |
| | Minimal Superkey. A superkey that does not contain a subset of attributes that is itself a |
| | superkey. |
| Candiate Key | |
| | Example: Student_name and Student_street, are sufficient to uniquely identify one |
| | particular student. |
| | The candidate key selected to uniquely identify all rows. It should be rarely changed and |
| Primary Key | cannot contain null values. |
| | |
| | Example: Roll_No is a primary set of _student' entity set. |
| | An attribute (or combination of attributes) in one table that must either match the primary |
| | key of another table or be null |
| Foreign Key | |
| | Example: Consider in the staff relation the branch_no attribute exists to match staff to the |
| | branch office they work in. In the staff relation, branch_no is foreign key. |
| Secondary Key | An attribute or combination of attributes used to make data retrieval more efficient. |
| | |
| | |

3. Participation Constraint

Participation can be divided into two types. 1. Total 2. Partial

If every entity in the entity set E participates in at least one relationship in R. Then participation is called Total Participation

 \triangleright

If only some entities in the entity set E participate in relationships in R. Then the participation is called Partial Participation.

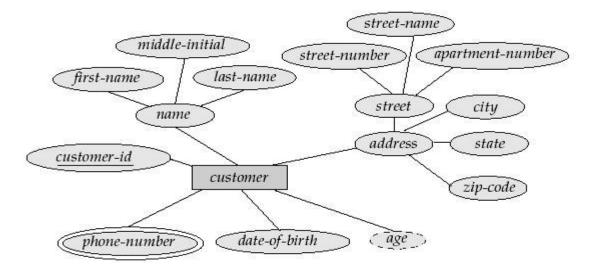
9. ENTITY-RELATIONSHIP(E-R) DIAGRAMS

E-R diagram can express the overall logical structure of a database graphically.

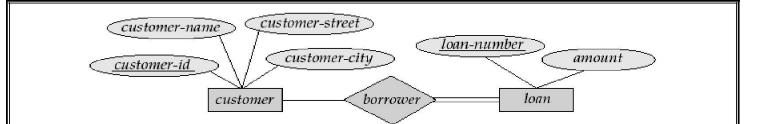
E-R diagram consists of the following major components:

| Component name | Symbol | Description |
|-------------------|---------------------|---|
| Rectangles | | represent entity sets |
| Ellipses | | represent attributes |
| Diamonds | \bigcirc | represent relationship sets |
| Lines | | link attributes to entity sets and entity sets to relationship sets |
| Double ellipses | | represent multivalued attributes |
| Dashed ellipses | $\langle A \rangle$ | represent derived attributes |
| Double lines | | Represent total participation of an entity in a relationship set |
| Double rectangles | represe | nt weak entity sets |

E-R diagram with composite, multivalued, and derived attributes.



Double lines are used in an E-R diagram to indicate that the participation of an entity set in a relationship set is total; that is, each entity in the entity set occurs in at least one relationship in that relationship set.



The number of time an entity participates in a relationship can be specified using complex **cardinalities.**

An edge between an entity set and binary relationship set can have an associated minimum and maximum cardinality assigned in the form of l..h.

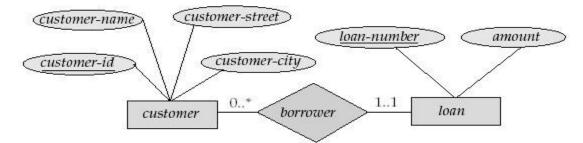
l - Minimum cardinalityh - Maximum cardinality

A minimum value of 1 indicates total participation of the entity set in the relationship

set. A maximum value of 1 indicates that the entity participates in at most one

relationship. A maximum value * indicates no limit.

A label 1... on an edge is equivalent to a double line.



0..* indicates a customer can have 0 or more loan

1..1 indicates a loan must have one associated customer

Strong and Weak entity sets

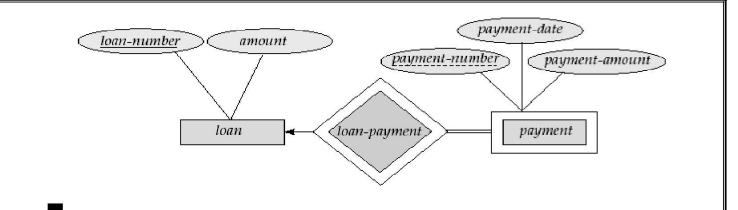
An entity set may not have sufficient attributes to form a primary key. Such an entity set is termed a **weak entity set**.

An entity set that has a primary key is termed a strong entity set.

Weak entity set is associated with another entity set called the **identifying** or **owner entity set**. ie, weak entity set is said to be existence dependent on the identifying entity set.

Identifying entity set is said to own the weak entity set.

The relationship among the weak and identifying entity set is called the **identifying relationship**.



Discriminator in a weak entity set is a set of attributes that distinguishes the different entities among the weak entity also called as partial key.

Extended E-R Features

ER model that is supported with the additional semantic concepts is called the extended entity relationship model or EER model.

EER model deals with

- 1. Specialization
- 2. Generalization
- 3. Aggregation

1. Specialization:

The process of designating subgroupings within an entity set is called

Specialization Specialization is a top-down process.

Consider an entity set person. A person may be further classified as one of the following:

- Customer
- Employee

All person has a set of attributes in common with some additional

attributes. Specialization is depicted by a triangle component labeled ISA.

The label ISA stands for is a for example, that a customer is a person.

The ISA relationship may also be referred to as a super class-subclass relationship.

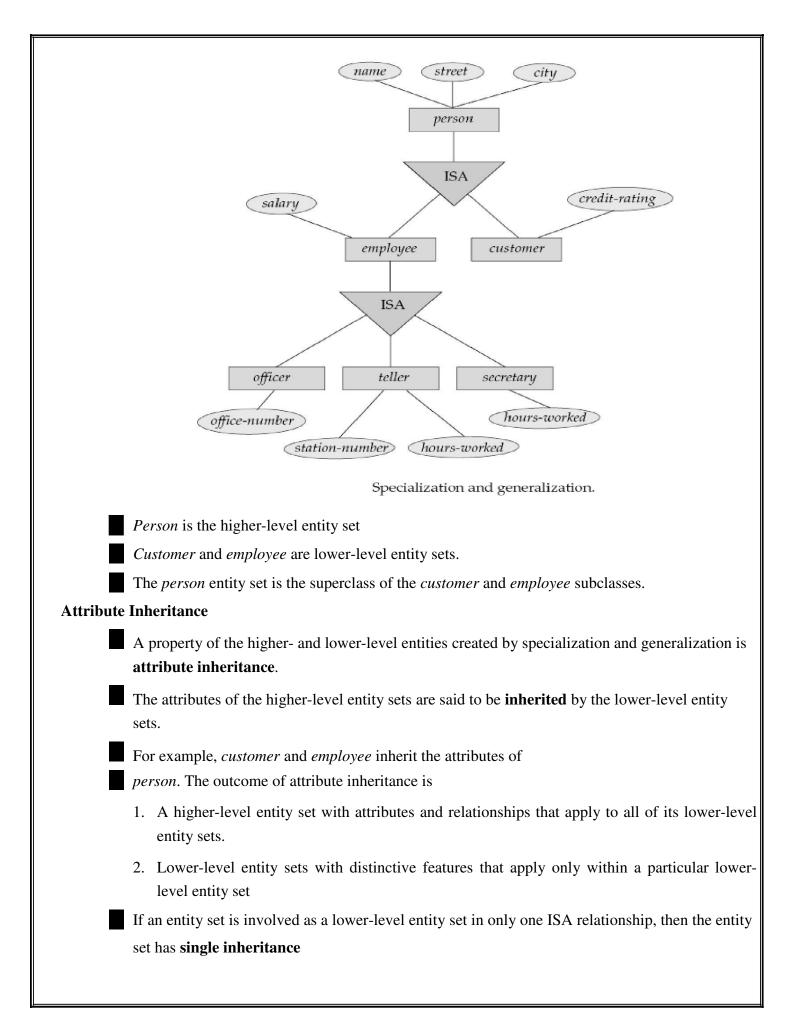
2. Generalization:

Generalization is a simple inversion of specialization.

Generalization is the process of defining a more general entity type from a set of more specialized entity types.

Generalization is a bottom-up approach.

Generalization results in the identification of a generalized super class from the original subclasses.



If an entity set is involved as a lower-level entity set in more than one ISA relationship, then the entity set has **multiple inheritance** and the resulting structure is said to be a *lattice*.

Constraints on Generalizations

1. One type of constraint determining which entities can be members of a lower-level entity set. Such membership may be one of the following:

• **Condition-defined**. In condition-defined the members of lower-level entity set is evaluated on the basis of whether or not an entity satisfies an explicit condition.

- User-defined. User defined constraints are defined by user.
- 2. A second type of constraint relates to whether or not entities may belong to more than one lowerlevel entity set within a single generalization. The lower-level entity sets may be one of the following:

• **Disjoint**. A *disjointness constraint* requires that an entity belong to no more than one lower-level entity set.

• **Overlapping**. Same entity may belong to more than one lower-level entity set within a single generalization.

3. A final constraint, the **completeness constraint** specifies whether or not an entity in the higher-level entity set must belong to at least one of the lower-level entity sets .This constraint may be one of the following:

• **Total generalization** or **specialization**. Each higher-level entity must belong to a lower-level entity set. It is represented by double line.

• Partial generalization or specialization. Some higher-level entities may not belong to any lower-level entity set.

3. Aggregation

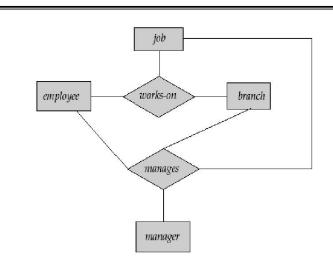
One limitation of the E-R model is that it cannot express relationships among relationships.

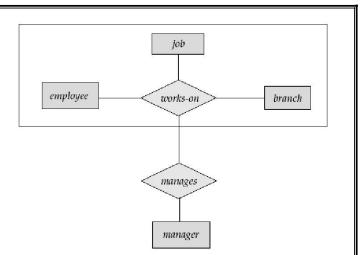
Consider the ternary relationship *works-on*, between a *employee*, *branch*, and *job*. Now, suppose we want to record managers for tasks performed by an employee at a branch. There another entity set *manager* is created.

The best way to model such a situation is to use aggregation.

Aggregation is an abstraction through which relationships are treated as higherlevel entities.

In our example works-on act as high level entity.

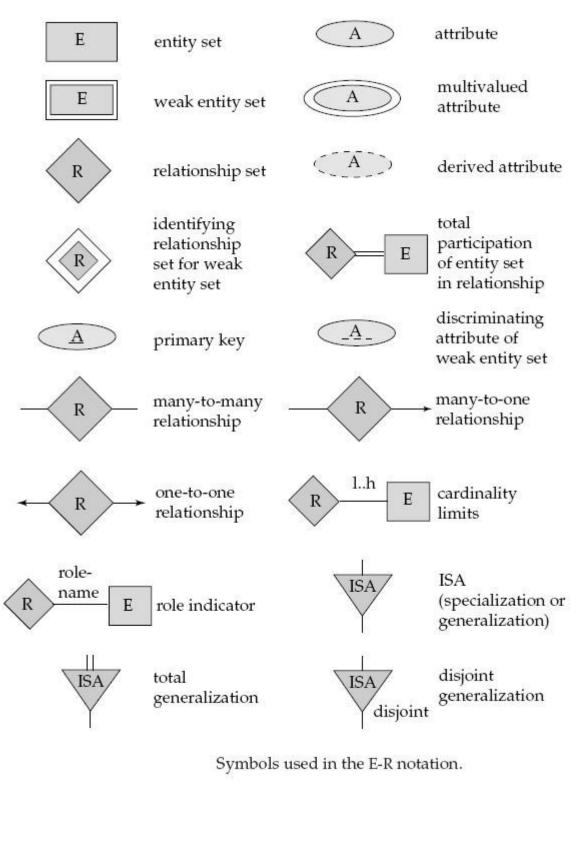




E-R diagram with redundant relationships.

E-R diagram with aggregation.

Summary of ER diagram



#. INTROD TION TO RELATIONALDATA ASES

A relational database is based on the relational model and uses a collection of tables to represent both data and the relationship among those data.

It includes DML and DDL languages.

Tables:

Each table has multiple columns and each column has unique name.

| Account number | Balance |
|----------------|---------|
| A-101 | 500 |
| A-215 | 700 |
| A-102 | 400 |

A relational model is an example of a record based model.

Record based model are structured in fixed format record of several types.

Each table contains record of particular type. Each record type defines a fixed number of fields or attributes.

The columns of the table correspond to the attribute of record type.

Data Manipulation Language (DML)

DML includes following commands

- 1. INSERT To insert one or more number of Rows.
- 2. SELECT To display one or more rows.
- 3. UPDATE Used to alter the column values in a table.
- 4. DELETE Used to delete one or more rows.

Data Definition Language (DDL)

DDL includes following commands

- 1. CREATE Command used for creating tables.
- 2. DESC Command used to view the table structure.
- 3. ALTER Command used for modifying table structure.
- 4. RENAME used to change the name of the table.
- 5. DROP Command used for removing an existing table.

2 Mark Questions

- 1. Define data, database, database management system, database system?
- 2. List any eight applications of DBMS.
- 3. What are the disadvantages of keeping organization information in a file processing system?
- 4. What are the advantages of using a DBMS (Centralized control of data)?
- 5. With the block diagram, discuss briefly the various levels of data abstraction?
- 6. Define instance and schema?
- 7. Define the terms 1) physical schema 2) logical schema 3)Subschema
- 8. What is conceptual schema?
- 9. Define data model?
- 10. What is storage manager?
- 11. What are the components of storage manager?
- 12. What is the purpose of storage manager?
- 13. List the data structures implemented by the storage manager.
- 14. What is a data dictionary?
- 15. What is an entity relationship model?
- 16. What are attributes? Give examples.
- 17. What are the types of Attributes?
- 18. What is relationship? Give examples
- 19. Define the terms
- 20. Define null values.
- 21. Define the terms
- 22. What is meant by the degree of relationship set?
- 23. Define the terms
- 24. Define weak and strong entity sets?
- 25. What does the cardinality ratio specify?
- 26. Explain the two types of participation constraint.
- 27. Define the terms
- 28. Write short notes on relational model
- 29. Define the term Domain.
- 30. Specify with suitable examples, the different types of keys used in database management systems.
- 31. Define Data model.

16 Mark Questions

- 1. Explain DBMS System Architecture.
- 2. Explain E-R Model in detail with suitable example.
- 3. Explain about various data models.
- 4 Draw an E R Diagram for Banking, University, Company, Airlines, ATM, Hospital, Library, Super market, Insurance Company.
- 5. Explain in details about the various database languages.
- 6. Discuss about various operations in Relational Databases.
- 7. Discuss about database users and administrators.

UNIT II

RELATIONAL MODEL

SQL – Data definition- Queries in SQL- Updates- Views – Integrity and Security – Relational Database design-Relational Models-Design issues – Functional dependences and Normalization for Relational Databases (up to BCNF).

1. THE RELATIONAL MODEL

STRUCTURE OF RELATIONAL DATABASES:

A relational database consists of a collection of **tables**, each of which is assigned a unique name.

A row in a table represents a *relationship* among a set of values.

BASIC STRUCTURE

Each column header is attributes. Each attribute allows a set of permitted values called domain of that attribute.

A table of n-attributes must be a subset of

 $D1 \times D2 \times \cdot \cdot \cdot \times Dn.1 \times Dn$

A relation is a cartesian product of list of domains.

Mathematically table is called as a relation and rows in a table are called as tuples.

The tuples in a relation can be either **sorted or unsorted**.

Several attributes can have **same domain**. **E.g.:** customer_name,

employee_name. Attributes can also be distinct. E.g.: balance, branch_name

Attributes can have null values incase if the value is unknown or does not exist.

Database schema begins with upper case and database relation begins with lower case. Account-schema = (account-number, branch-name, balance) account (Account-schema)

| account-number | number branch-name | |
|----------------|--------------------|-----|
| A-101 | Downtown | 500 |
| A-215 | Mianus | 700 |
| A-102 | Perryridge | 400 |
| A-305 | Round Hill | 350 |
| A-201 | Brighton | 900 |
| A-222 | Redwood | 700 |
| A-217 | Brighton | 750 |

Account Table

. THE ATALO

The catalog is a place where all the schemas and the corresponding mappings are kept.

The catalog contains detailed information also called as descriptor information or meta data. Descriptor information is essential for the system to perform its job properly.

For example the authorization subsystem uses catalog information about users and security constraints to grant or deny access to a particular user.

The catalog should be self describing.

. RELATIONAL AL E RA

The relational algebra is a *procedural* query language.

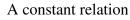
It consists of a set of operations that take one or two relations as input and produce a new relation as their result.

Formal Definition

A basic expression in the relational algebra consists of either one of the following:

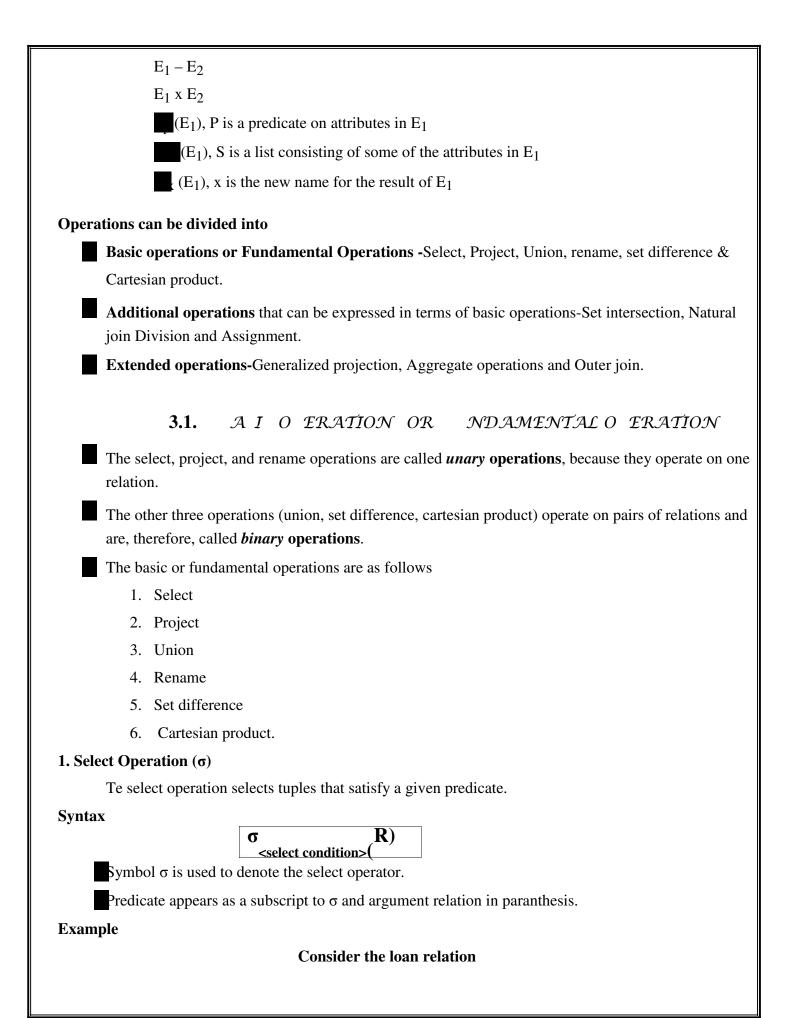


A relation in the database



Let E_1 and E_2 be relational-algebra expressions; the following are all relational-algebra

exp sions: $E_1 E_2$



| loan-number | branch-name | amount |
|-------------|-------------|--------|
| L-11 | Round Hill | 900 |
| L-14 | Downtown | 1500 |
| L-15 | Perryridge | 1500 |
| L-16 | Perryridge | 1300 |
| L-17 | Downtown | 1000 |
| L-23 | Redwood | 2000 |
| L-93 | Mianus | 500 |

Query:

(loan)

Output relation is

σ

| loan-number | branch-name | amount |
|-------------|-------------|--------|
| L-15 | Perryridge | 1500 |
| L-16 | Perryridge | 1300 |

Select operation allows all comparisons using =, _=, <, ., >,.

It allows combination of server predicates using connectives like and (A), or (V), and not (\neg).

branch-name =—Perryridge

E.g.: 1. $\sigma_{amount>1200}(loan)$

Other Examples

Consider following Book relation.

| Book_Id | Title | Author | Publisher | Year | Price |
|---------|----------|----------|-------------|------|-------|
| B001 | DBMS | Korth | McGraw_Hill | 2000 | 250 |
| B002 | Compiler | Ulman | | 2004 | 350 |
| B003 | OOMD | Rambaugh | | 2003 | 450 |
| B004 | PPL | Sabista | | 2000 | 500 |

Following are the some examples of the select operation.

Example 1: Display books published in the 2000.

Query 1: $\sigma_{year=2000}(Book)$

The output of query 1 is shown below.

| Book_Id | Title | Author | Publisher | Year | Price |
|---------|-------|---------|-------------|------|-------|
| B001 | DBMS | Korth | McGraw_Hill | 2000 | 250 |
| B004 | PPL | Sabista | | 2000 | 500 |

Example 2: Display all books having price greater than 300.

Query 2: $\sigma_{\text{price}>300}(\text{Book})$

The output of query 2 is shown below.

| Book_Id | Title | Author | Publisher | Year | Price |
|---------|----------|----------|-----------|------|-------|
| B002 | Compiler | Ulman | | 2004 | 350 |
| B003 | OOMD | Rambaugh | | 2003 | 450 |
| B004 | PPL | Sabista | | 2000 | 500 |

Example 3: Select the tuples for all books whose publishing year is 2000 or price is greater than 300.

Query 3: σ(year=2000) OR (price>300)(Book)

The output of query 3 is shown below.

| Book_Id | Title | Author | Publisher | Year | Price |
|---------|----------|----------|-------------|------|-------|
| B001 | DBMS | Korth | McGraw_Hill | 2000 | 250 |
| B002 | Compiler | Ulman | | 2004 | 350 |
| B003 | OOMD | Rambaugh | | 2003 | 450 |
| B004 | PPL | Sabista | | 2000 | 500 |

Example 4: Select the tuples for all books whose publishing year is 2000 and price is greater than 300.

Query 3: σ (year=2000) AND (price>300)(Book)

The output of query 4 is shown below.

| Book_Id | Title | Author | Publisher | Year | Price |
|---------|-------|---------|-----------|------|-------|
| B004 | PPL | Sabista | | 2000 | 500 |
| | | | | | |

2. Project operation (Π)

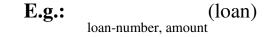
The project operation selects certain columns from a table while discarding others. It removes any duplicate tuples from the result relation.

Syntax



The symbol (pi) is used to denote the project operation

Attribute list to be projected is specified as subscript of and R denotes the relation.



| loan-number | amount |
|-------------|--------|
| L-11 | 900 |
| L-14 | 1500 |
| L-15 | 1500 |
| L-16 | 1300 |
| L-17 | 1000 |
| L-23 | 2000 |
| L-93 | 500 |

Example: The following are the examples of project operation on Book relation.

Example 1: Display all titles with author name.

Query 1: Title, Author (Book)

The output of query 1 is shown below.

| Title | Author |
|----------|----------|
| DBMS | Korth |
| Compiler | Ulman |
| OOMD | Rambaugh |
| PPL | Sabista |

Example 2: Display all book titles with authors and price.

Query 2: Title, Author, Price (Book)

The output of query 2 is shown below.

| Title | Author | Price |
|----------|----------|-------|
| DBMS | Korth | 250 |
| Compiler | Ulman | 350 |
| OOMD | Rambaugh | 450 |
| PPL | Sabista | 500 |

Composition of select and project operations

The relational operations select and project can be combined to form a complicated query.

customer-name (σ (customer)) customer-city = Harrison Input Table: customer

| customer-name | customer-street | customer-city |
|---------------|-----------------|---------------|
| Adams | Spring | Pittsfield |
| Brooks | Senator | Brooklyn |
| Curry | North | Rye |
| Glenn | Sand Hill | Woodside |
| Green | Walnut | Stamford |
| Hayes | Main | Harrison |
| Johnson | Alma | Palo Alto |
| Jones | Main | Harrison |
| Lindsay | Park | Pittsfield |
| Smith | North | Rye |
| Turner | Putnam | Stamford |
| Williams | Nassau | Princeton |

| Question | Option 1 | Option 2 | Option3 | Option 4 | Answer |
|---|---|--------------------------------------|---|--|--|
| The DBMS acts as an interface between what two components of an enterprise-class database system? | Database application and the database | Data and the database | The user and the database application | Database application and SQL | Database application and the database |
| Which of the following products was an early implementation of the relational model developed by E.F. Codd of IBM? | Nothing | DB2 | dBase | R:base | DB2 |
| The following are components of a database except | user data | metadata | reports | indexes | reports |
| An application where only one user accesses the database at a given time is an example of a(n) | single-user database application | multiuser database application | e-commerce database application | data mining database application | single-user database application |
| An on-line commercial site such as Amazon.com is an example of a(n) | single-user database application | multiuser database application | e-commerce database application | data mining database application | e- commerce database application |
| SQL stands for | Structured Query Language | Sequential Query Language | Structured Question Language | Sequential Question Language | Structured Query Language |
| Because it contains a description of its own structure, a database is considered to be | described | metadata compatible | self-describing | an application program | self- describing |
| The following are functions of a DBMS except | creating and processing forms | creating databases | processing data | administrating databases | creating and processing forms |
| Helping people keep track of things is the purpose of a(n) | database | table | instance | relationship | database |
| An Enterprise Resource Planning application is an example of a(n) | single-user database application | multiuser database application | e-commerce database application | data mining database application | multiuser database application |
| A DBMS that combines a DBMS and an application generator is | Microsoft's SQL Server | Microsoft's Access | IBM's DB2 | Oracle Corporation's Oracle | Microsoft's Access |

| You have run an SQL statement that asked the DBMS to display data in a table named USER_TABLES. The results include columns of data labeled "TableName," "NumberOfColumns" and "PrimaryKey." You are looking at | user dat | metadata | A report | indexes | metadata |
|---|---------------------------|--------------------------|------------------------|-----------------------------------|------------------------------------|
| Which of the following is not considered to be a basic element of an enterprise-class database system? | Users | Database applications | DBMS | COBOL programs | COBOL programs |
| The DBMS that is most difficult to use is | Microsoft's SQL Server | Microsoft's Access | IBM's DB2 | Oracle Corporation's Oracle | Oracle Corporation 's Oracle |
| Which of the following indicates the maximum number of entities that can be involved in a relationship? | Minimum cardinality | Maximum cardinality | ERD | Greater Entity Count (GEC) | Maximum cardinality |
| Which type of entity cannot exist in the database unless another type of entity also exists in the database, but does not require that the identifier of that other entity be included as part of its own identifier? | Weak entity | Strong entity | ID-dependent entity | ID- independent entity | Weak entity |
| In a one-to-many relationship, the entity that is on the one side of the relationship is called a(n) entity. | parent | child | instance | subtype | parent |

| Which type of entity represents an actual occurrence of an associated generalized entity? | Supertype entity | Subtype entity | Archetype entity | Instance entity | Instance entity |
|---|----------------------------|-------------------------------------|----------------------------------|-------------------------------|---------------------------------|
| A recursive relationship is a relationship between an entity and | itself | a subtype entity | an archetype entity | an instance entity | itself |
| Which of the following indicates the minimum number of entities that must be involved in a relationship? | Minimum cardinality | Maximum cardinality | ERD | Greater Entity Count (GEC) | Minimum cardinality |
| Which of the following refers to something that can be identified in the users' work environment, something that the users want to track? | Entity | Attribute | Identifier | Relationship | Entity |
| In which of the following is a single-entity instance of one type related to many entity instances of another type? | One-to-One Relationship | One-to- Many Relationshi p | Many-to- Many Relationship | Composite Relationship | One-to- Many Relationship |
| Which of the following refers to an entity in which the identifier of one entity includes the identifier of another entity? | Weak entity | Strong entity | ID-dependent entity | ID- independent entity | ID- dependent entity |
| Which type of entity is related to two or more associated entities that each contain | Supertype entity | Subtype entity | Archetype entity | Instance entity | Supertype entity |

| specialized attributes that apply to some but not all of the instances of the entity? | | | | | |
|---|----------------------------|-------------------------------------|----------------------------------|---------------------------|----------------------------------|
| An attribute that names or identifies entity instances is a(n): | entity. | attribute. | identifier. | relationship. | identifier. |
| Properties that describe the characteristics of entities are called: | entities. | attributes. | identifiers. | relationships. | attributes. |
| In which of the following can many entity instances of one type be related to many entity instances of another type? | One-to-One Relationship | One-to- Many Relationshi p | Many-to- Many Relationship | Composite Relationship | Many-to- Many Relationship |
| Entities of a given type are grouped into a(n): | database. | entity class. | attribute. | ER | entity class. |
| Which of the following is NOT a basic element of all versions of the E-R model? | Entities | Attributes | Relationships | Primary keys | Primary keys |
| In which of the following is a single-entity instance of one type of related to a single-entity instance of another type? | One-to-One Relationship | One-to- Many Relationshi p | Many-to- Many Relationship | Composite Relationship | One-to- One Relationship |
| Entities can be associated with one another in which of the following? | Entities | Attributes | Identifiers | Relationships | Relationship s |
| Which type of entity has its relationship to another entity determined by an attribute in that other entity called a discriminator? | Supertype entity | Subtype entity | Archetype entity | Instance entity | Subtype entity |
| Which type of entity represents a logical generalization whose actual occurrence is represented by a second, associated entity? | Supertype entity | Subtype entity | Archetype entity | Instance entity | Archetype entity |
| In a one-to-many relationship, the entity that is on the many side of the relationship is called a(n)entity. | parent | child | instance | subtype | child |

| Which of the following data constraints would be used to specify that the value of cells in a column must be one of a specific set of possible values? | A domain constraint | A range constraint | An intrarelation constraint | An interrelation constraint | A domain constraint |
|--|--|-----------------------------------|---|---|--|
| In a 1:N relationship, the foreign key is placed in: | either table without specifying parent and child tables. | the parent table. | the child table. | either the parent table or the child table. | the child table. |
| Which of the following column properties specifies whether or not cells in a column must contain a data value? | Null status | Data type | Default value | Data constraints | Null status |
| A primary key should be defined as: | NULL. | NOT NULL. | Either of the above can be use | None of the above are correct. | NOT NULL. |
| Which of the following column properties would be used to specify that cells in a column must contain a monetary value? | Null status | Data type | Default value | Data constraints | Data type |
| Which of the following situation requires the use of ID- dependent entities? | Association relationships only | Multivalued attributes only | Archetype/inst ance relationships only | All of the above use ID dependent entities | All of the above use ID dependent entities |
| A foreign key is: | a column containing the primary key of another table. | used to define data types. | used to define null status. | all of the above are above correct. | a column containing the primary key of another table. |
| Which of the following columns is(are) are required in a table? | A foreign key | An alternate key | A primary key | A surrogate key. | A primary key |
| In a 1:1 relationship, the foreign key is placed in: | either table without specifying parent and child tables. | the parent table. | the child table. | either the parent table or the child table. | either table without specifying parent and child tables. |
| Which of the following column properties would be used to specify that cells in a column must be immediately filled with a monetary value of \$10,000? | Null status | Data type | Default value | Data constraints | Default value |
| The identifier of an entity will become theof the new table. | foreign key | main attribute | primary key | identity key | primary key |

| Which of the following data constraints would be used to specify that the value of a cell in one column must be less than the value of a cell in another column in the same row of the same table? | A domain constraint | A range constraint | An intrarelation constraint | An interrelation constraint | An intrarelatio n constraint |
|--|---|---|--------------------------------------|-------------------------------------|--------------------------------------|
| A unique, DBMS-supplied identifier used as the primary key of a relation is called a(n): | primary key. | foreign key. | composite key. | surrogate key. | surrogate key. |
| Which is not true about surrogate keys? | They are short. | They are fixe | They have meaning to the user. | They are numeri | They have meaning to the user. |
| For every relationship, how many possible types of actions are there when enforcing minimum cardinalities? | Two | Three | Four | Six | Six |
| Which constraint requires that the binary relationship indicate all combinations that must appear in the ternary relationship? | MUST COVER | MUST NOT | Both of the above. | None of the above is correct. | MUST COVER |
| Each entity is represented as a(n): | tuple. | table. | attribute. | file. | table. |
| For every relationship, how many possible sets of minimum cardinalities are there? | Тwo | Three | Four | Six | Four |
| If a relationship has a cascade updates constraint, then if in the parent table is changed, then the same change will automatically be made to any corresponding foreign key value. | the primary key | any alternate key | a surrogate key | a foreign key | the primary key |
| Which of the following column properties would be used to specify that cells in a column must contain a monetary value that is less than another monetary value in the same row? | Null status | Data type | Default value | Data constraints | Data constraints |
| Poor data administration can lead to which of the following? | A single definition of the same data entity | Familiarity with existing data | Missing data elements | All of the above. | Missing data elements |

| A traditional data administrator performs which of the following roles? | Tune database performance | Establish backup and recovery procedures | Resolve data ownership issues | Protect the security of the database. | Resolve data ownership issues |
|--|------------------------------|---|---|--|---|
| If both data and database administration exist in an organization, the database administrator is responsible for which of the following? | Data modeling | Database design | Metadata | All of the above. | Database design |
| Which of the following is part of an administrative policy to secure a database? | Authentication policies | Limiting particular areas within a building to only authorized people | Ensure appropriate responses rates are in external maintenance agreements | All of the above. | All of the above. |
| The fact that the same operation may apply to two or more classes is called what? | Inheritance | Polymorphi sm | Encapsulation | Multiple classification | Polymorphi sm |
| The transaction log includes which of the following? | The before-image of a record | The after- image of a record | The before and after- image of a record | The essential data of the record | The essential data of the record |

DATA BASE MANAGEMENT SYSTEMS

QUESTION BANK

UNIT I

PART A

1. What is Database Management System? Why do we need a DBMS ?

DBMS is a collection of interrelated data and a set of programs to access those data. It is to provide a way to store and retrieve information that is both convenient and efficient.

2. List any two advantages of database systems.

Security, Integrity, Atomicity, Concurrent access anomalies

3. Define Data Model and its types.

Data model is a collection of tools for describing Data, Data relationships, Data semantics, Data constraints. Types: Relational model, Entity-Relationship data model Object-based data models, Semi structured data model (XML), Network model, Hierarchical model.

4. Explain the role and functions of the database administrator(DB Manager).

Coordinates all the activities of the database system, Storage structure and access method definition, Schema and physical organization modification, Granting users authority to access the database, Backing up data, Monitoring performance and responding to changes.

5. Give the limitations of E-R model? How do you overcome this?

- 2. .Limited constraint representation
- 3. .Limited relationship representation
- 4. .No data manipulation language
- 5. .Loss of information content

6. What are the limitations of file system.

Data redundancy and inconsistency, Difficulty in access data, data isolation, integrity, atomicity, concurrent access anomalies.

7. What is logical data Independency ?.

Application programs are said to exhibit physical data independence if they do not depend on the physical schema and thus need not be rewritten if the physical schema changes.

8. Define DML?.

DML is a language that enables users to access or manipulate data as organized by the appropriate data model. Data Manipulation Language Insert, Select, Delete

9. Define Data Dictionary?.

A data dictionary is a data structure which stores meta data about the structure of the database ie. the schema of the database.

10. Define Data independence.

Application programs are said to exhibit physical data independence if they do not depend on the physical schema and thus need not be rewritten if the physical schema changes.

11. With an relevant example explain Ternary Relationship.

A relationship is an association among several entities. Ternary relationship has three relationship. Eg. Parent related to child.

12. List any eight applications of DBMS.

a) Banking

- b) Airlines
- c) Universities
- d) Credit card transactions
- e) Tele communication
- f) Finance
- g) Sales
- h) Manufacturing
- i) Human resources

13. What are the advantages of DBMS?.

Controlling redundancy

- b) Restricting unauthorized access
- c) Providing multiple user interfaces
- d) Enforcing integrity constraints.
- e) Providing back up and recovery

14. Give the levels of data abstraction?

- a) Physical level
- b) logical level
- c) view level

15. Define instance and schema?

Instance: Collection of data stored in the data base at a particular moment is called an Instance of the database.

Schema: The overall design of the data base is called the data base schema.

16. Define the terms 1) physical schema 2) logical schema.

Physical schema: The physical schema describes the database design at the physical level, which is the lowest level of abstraction describing how the data are actually stored. **Logical schema:** The logical schema describes the database design at the logical level, which describes what data are stored in the database and what relationship exists among the data.

17. What is conceptual schema?

The schemas at the view level are called subschemas that describe different views of the database.

18. What is storage manager?

A storage manager is a program module that provides the interface between the

low level data stored in a database and the application programs and queries submitted to the system.

19. What are the components of storage manager?

The storage manager components include

- a) Authorization and integrity manager
- b) Transaction manager
- c) File manager
- d) Buffer manager

20. What is the purpose of storage manager?

The storage manager is responsible for the following

a) Interaction with he file manager

- b) Translation of DML commands in to low level file system commands
- c) Storing, retrieving and updating data in the database

21. List the data structures implemented by the storage manager.

The storage manager implements the following data structure

- a) Data files
- b) Data dictionary
- c) indices

PART B

- 1. Briefly explain about Database system architecture:
- 2. Explain about the Purpose of Database system:
- 3. Briefly explain about Views of data:
- 4. Explain about different kinds of data models:
- 5. Explain about Database Languages:
- 6. Briefly explain about Entity-Relationship model:

7. Explain about Relational Databases:

8. Explain about Advantages and disadvantages of DBMS:

9.i. Construct an ER diagram for a car insurance company that has a set of customers, each of whom owns one/more cars. Each car has associated with it zero to any number of recorded accidents,

ii. Construct appropriate tables for the above ER diagram.

10.i. Define data model. Explain the difference types of data models with relevant examples.

ii. Explain the role and functions of data base administrator.

11. Discuss in detail about database system architecture with neat Diagram.

12. Draw an E-R diagram for a banking enterprise with almost all components and Explain.

13.Explain the architecture of database system.

(i) With a neat Diagram, explain the structure of a DBMS.

(ii)Draw an F-R diagram for a small marketing company database, assuming your own data requirements.

14.(i) Compare the features of file system with database system.

(ii) Explain the differences between physical level, conceptual level and view level of data abstraction.

(iii) Mention any four major responsibilities of DBA.

15.(i)Discuss the various disadvantages in the file system and explain how it can be overcome by the database "system.

(ii)What are the different Data models present? Explain in detail.

16.(i)Explain the Database system structure with a neat Diagram.

(ii) Construct an ER diagram for an employee payroll system.

17.(i) Construct an ER diagram for a car insurance company that has a set of Customers, each of whom owns one/more cars. Each car has associated with it zero to any number of recorded accidents.

ii)Construct appropriate tables for the above ER diagram.

18.(i) Define data model. Explain the different types of data models with relevant Examples.

(ii) Explain the role and functions of the database administrator.

DATA BASE MANAGEMENT SYSTEMS

QUESTION BANK

UNIT II

1. Define query language. Give the classification of query language.

Query language is used to query a database. It can define the structure of the data, modify and specify security constraints. Procedural and non procedural query language

2. What are three characteristics of a Relational Database System?

Three characteristics of a Relational Database System are BCNF, Lossless join and Dependency preservation.

3. State the differences between Security and Integrity?

Integrity ensures that changes made to the database by authorized users do not result in loss of data consistency. Eg. Check constraints. Security refers to protection from malicious access. It is done in database systems, OS, Network physical, Human.

4. What are the pitfalls in relational database design?.

Repetition of Information.

PART A

Inability to represent certain information.

5. Distinguish between primary key and candidate key.

Primary key denotes a candidate key for identify entities with in entity set. Super keys for which no proper subset is a super key is called candidate key.

6. Give the reasons why Null values might be introduces into the database.

Null values are introduced into the database to indicate the absence of information about the value. Eg. Select loan number form loan where amount is null.

7. What is static SQL? How does it differ from dynamic SQL?

In static SQL, SQL queries must be present in compile time. Dynamic SQL allows programs to construct and submit SQL queries at run time.

8. What are the different types of integrity constraints used in designing a relational database?

Domain constraints, Referential Integrity, Check constraints, Assertion, Triggers.

9. With an example explain a weak entity in an ER diagram.

An entity without primary key is weak entity. Eg. Payment number in payment relation having payment number payment date and payment amount as fields.

10. What is referential integrity.

The condition that value which is appearing in one relation for a given set of attributes also appears for a certain set of attributes in another relation is called as referential integrity.

11. What is domain integrity? Give example.

There are many domain types, such as integer, character and date/time. Create domain can be used to define new domains. Value of one type can be cast to other.

Eg. Create domain Dollars numeric (12,2) Create domain Pounds numeric (12,2) Cast r.A as pounds

12. Consider the following relations. Empno(Eno,Name,Date_of_Birth, Sex,Date_Of_Joining,Basic_Pay,Dept)

Develop an SQL query that will find and display the Dept and Average Basic_Pay in eac Dept.

Select dept, Average(Basic_Pay) from empno Group by dept.

13. Define a distributed database management system.

A distributed database system consists of loosely coupled sites that share no physical component. Database systems that run on each site are independent of each other. Transactions may access data at one or more sites

14. List the SQL statements used for Transaction control.

Commit, roll back.

15. What is a nested relation ?

Attributes in relations have multiple values. Example: library information system. Each book has title, a set of authors, Publisher, and a set of keywords.

16. State the advantages of distributed systems.

18. A distributed database system consists of loosely coupled sites that share no physical component

- 19. Database systems that run on each site are independent of each other
- 20. Transactions may access data at one or more sites

17. Name the different types of joins supported in SQL.

Inner join, Natural inner join, natural left join, natural right join and full outer join.

18. Define the terms fragmentation and replication, in terms of where data is stored.

System maintains several identical replica of relation and stores each replica and a different site. Fragmentation is that the system partitions the relation into several fragments and stores each fragment at a different site.

19. What are the types of transparencies that a distributed database must support? Why?

Fragmentation, Replication and location. The user of a distributed database system should not be required to know either where the data are physically located ore how the data can be accessed at the specific local site.

20. What is an entity relationship model?

The entity relationship model is a collection of basic objects called entities and relationship among those objects. An entity is a thing or object in the real world that is distinguishable from other objects.

21. Define the terms

1.Entity set

2.Relationship set

Entity set: The set of all entities of the same type is termed as an entity set. **Relationship set**: The set of all relationships of the same type is termed as a relationship set.

22. Define single valued and multivalued attributes.

Single valued attributes: attributes with a single value for a particular entity are called single valued attributes.

Multivalued attributes: Attributes with a set of value for a particular entity are called multivalued attributes.

23. What are stored and derived attributes?

Stored attributes: The attributes stored in a data base are called stored attributes. **Derived attributes:** The attributes that are derived from the stored attributes are called derived attributes.

24. What are composite attributes?

Composite attributes can be divided in to sub parts.

25. Define null values.

In some cases a particular entity may not have an applicable value for an attribute or if we do not know the value of an attribute for a particular entity. In these cases null value is used.

26. What is a candidate key?

Minimal super keys are called *candidate keys*.

27. What is a primary key?

Primary key is chosen by the database designer as the principal means of identifying an entity in the entity set.28. What is a super key?

A *super key* is a set of one or more attributes that collectively allows us to identify uniquely an entity in the entity set.

29. Define- relational algebra.

The relational algebra is a procedural query language. It consists of a set of operations that take one or two relation as input and produce a new relation as output.

30. Write short notes on domain relational calculus

The domain relational calculus uses domain variables that take on values from an attribute domain rather than values for entire tuple.

31. What is foreign key?

A relation schema r1 derived from an ER schema may include among its

attributes the primary key of another relation schema r2.this attribute is called a *foreign key* from r1 referencing r2.

32. What is horizontal fragmentation?

Horizontal fragmentation splits the relation by assuming each tuple of r to one or more fragments.

33. What is vertical fragmentation?

Vertical fragmentation splits the relation by decomposing the scheme R of relation r.

34. Give the forms of triggers?

_ The triggering event can be insert or delete. _ For updated the trigger can specify columns. _ The referencing old row as clause

_ The referencing new row as clause

_ The triggers can be initiated before the event or after the event.

PART B

- 1. With relevant examples discuss the following in SQL.
 - i.Data Definition language
 - ii.Data Manipulation language
 - iii.Data Control language
 - iv.Views
- 2. What is normalization? Explain normalization techniques using functional dependencies with relevant examples.
- 3. Explain in detail about Relational Algebra, Domain Relational Calculus and Tuple Relational Calculus with suitable examples.
- 4. Briefly present a survey on Integrity and Security.
- 5. Explain in detail about INF, 2NF, 3NF and BCNF with suitable examples.
- 6. Describe about the Multi-Valued Dependencies and Fourth normal form with suitable example.
- 7. What is meant by relational calculus? Write query example for tuple relational and domain relational calculus
- 8. Write the overall process of designing a database application.
- 9. Write the SQL statement for the following
 - i) Union ii) Intersect iii) Except
- 10. i) Write short motes on integrity
 - ii) Compare BCNF and 3NF
- 11. (i)Consider the following relational database

employee (employee-name, street, city)

works (employee-name, camp any-narne, salary)

company (company-name, city)

manages (employee-name, manager-name)

Give an expression in SQL to express each of the following queries :

Find the names and cities of residence of all employees who work for XYZ Bank. Find the names, street address,

and cities of residence of all employees who work for XYZ Bank and earn more than Rs. 10,000 per annum.

Find the names of all employees in this database who live in the same city as the company for which they work.

Find the names of all employees who live in the same city and on

the same street as do their managers.

(ii)Define the term distributed database management system and mention the issues to be considered in the design of

the same.

- 12. (i) What are the relational algebra operations supported in SQL?Write the SQL statement for each operation.
 - (ii) What is data integrity? Explain the types of integrity constraints
- 13. (i) Explam INF, 2NF, 3OT and BCNF with suitable example.
 - (ii) Consider the universal relation E = {A, B, C, D, E, F, G, H, i}
 and the set of functional dependencies
 F = {(A, B) -> {C}, {A} ->{A E}, \B] ~> {F}, {F} -> {G, a}, {D} -> {i,
 - J}). What is the key for R? Decompose R into 2NF, then 3NF relations.
- 14. What are the pitfalls in relational database design? With a suitable example, explain the role of functional dependency in the process of normalization.
- 15. (i) Explain the use of trigger with your own example.(ii) Discuss the terms Distributed databases and client/server databases.
- 16. (i) What is a view? How can it be created? Explain with an example.(ii) Discuss in detail the operators SELECT, PROJECT, UNION with suitable Examples.
- 17. (a) Explain INF, 2NF and 3NF with an example.
- (b) Explain the Boyce-Codd normal form with an example. Also state how it Differs from that of 3NF.

- 19. (i) With relevant examples discuss the following in SQL.
 - (i)Data Definition Language.
 - (ii)Data Manipulation Language
 - (iii)Data Control Language
 - (iv)Views
- 20. What is normalization? Explain normalization techniques using functional Dependencies with relevant examples

DATA BASE MANAGEMENT SYSTEMS

QUESTION BANK

UNIT III

PART A

1. With an example explain a Functional Dependency.

Let *R* be a relation schema, a subset of .*R*. b *is also subset of R*. The functional dependency a->bholds on *R* if and only if for any legal relations *r*(R), whenever any two tuples *t*1 and *t*2 of *r* agree on the attributes a, they also agree on the attributes b. That is, t1[a] = t2[a] =>t1[b] = t2 [b] eg.

| Α | В | С |
|----|----|----|
| A1 | B1 | C1 |
| A1 | B2 | C1 |

Here A->C.

2. Define Normalization?

Normalization is to generate a set of relation schemas that allows to store information with out redundancy and to retrieve information easily

3. Justify the need for normalization.

A bad relational database have the following pit falls: inability to represent some data and redundancy. To overcome this pit falls, normalization is to be done. Good database design should be Lossless join and Dependency preservation. Decomposition of databases are done to make the database consistent and efficient.

4. Why it is necessary to decompose a relation?

A relation is to be decomposed in order to avoid repetition of information and inability to represent certain information.

5. What is decomposition and how does it address redundancy?

Decomposition is reducing a large database to set of smaller database. When splitting database, some of the fields will be repeated in databases to maintain association with records.

6. Give the comparison between BCNF and 3NF.

3NF can be achieved without sacrificing losslessness or dependency preservation. Disadvantage: Use null values to represent some of the meaningful relation and repetition of information. If it is difficult to get a dependency-preserving BCNF algorithm, it is preferable to opt for BCNF and use techniques such as materialized views.

7. Explain trivial dependency?

Functional dependency of the form a->b is trivial if b is subset of a. Trivial functional dependencies are satisfied by all the relations.

8. Define canonical cover?

A canonical cover Fc for F is a set of dependencies such that F logically implies all dependencies in FC and Fc logically implies all dependencies in F. Fc must have the following properties.

9. List the properties of canonical cover.

Fc must have the following properties.

_ No functional dependency in Fc contains an extraneous attribute. _ Each left side of a functional dependency in Fc is unique.

10. Explain the desirable properties of decomposition.

Lossless-join decomposition _
 Dependency preservation
 Repetition of information

11. What is 2NF?

A relation schema R is in 2NF if it is in 1NF and every non-prime attribute A in R is fully functionally dependent on primary key.

12. Define Boyce codd normal form

A relation schema R is in BCNF with respect to a set F of functional dependencies if, for all functional dependencies in F+ of the form a->b.

13. What is first normal form?

The domain of attribute must include only atomic (simple, indivisible) values.

14. What is meant by computing the closure of a set of functional dependency?

The closure of F denoted by F+ is the set of functional dependencies logically implied by F.

15. What are axioms?

Axioms or rules of inference provide a simpler technique for reasoning about functional dependencies.

16. What are the uses of functional dependencies?

_ To test relations to see whether they are legal under a given set of functional dependencies.

_ To specify constraints on the set of legal relations.

17. What is meant by normalization of data?

It is a process of analyzing the given relation schemas based on their Functional Dependencies (FDs) and primary key to achieve the properties

_ Minimizing redundancy

_ Minimizing insertion, deletion and updating anomalies.

18. Why certain Functional dependencies are called trivial dependency?. The functional dependency a-> β is trivial if βC a

19. What is multi valued dependency?

If A->B then, two tuples with same A value have different B value. Eg. Customer name -> customer street, customer city.

20. What is fourth normal form?.

A relation schema R is in 4NF with respect to a set D of functional and multivalued dependencies if for all multivalued dependencies in D+ of the form

 \Box . \Box , where \Box . \Box *R* and \Box . \Box *R*, at least one of the following hold: \Box is trivial (i.e., \Box . \Box . \Box or \Box . \Box . \Box = *R*)

2.□ is a superkey for schema *R* If a relation is in 4NF it is in BCNF

21. What are the closure set of functional dependency?.

-Given a set F of functional dependencies, there are certain other functional dependencies that are logically implied by F.

25. For example: If $A \square B$ and $B \square C$, then we can infer that $A \square C$

-The set of all functional dependencies logically implied by F is the *closure* of F. -We denote the *closure* of F by F+.

-F+ is a superset of F.

22. What are the goals of normalization?.

Let R be a relation scheme with a set F of functional dependencies. Decide whether a relation scheme R is in "good" form.

In the case that a relation scheme R is not in "good" form, decompose it into a set of relation scheme $\{R1, R2, ..., Rn\}$ such that

- 27. each relation scheme is in good form
- 28. the decomposition is a lossless-join decomposition
- 29. Preferably, the decomposition should be dependency preserving.

23. How would you find the extraneous attribute?.

Consider a set *F* of functional dependencies and the functional dependency $\Box . \Box . \Box$ in

F.

- Attribute A is extraneous in \Box if $A \Box \Box$

and *F* logically implies $(F - \{\Box.\Box.\Box\}) \Box \{(\Box - A) \Box.\Box\}$. -Attribute *A* is extraneous in \Box if $A \Box.\Box$ and the set of functional dependencies

 $(F - \{\Box.\Box.\Box\}) \Box \{\Box.\Box(\Box - A)\}$ logically implies *F*.

24. What do you mean by lossless join decomposition?.

A decomposition of *R* into *R*1 and *R*2 is lossless join if and only if at least one of the following dependencies is in F+:

o
$$R1 \square R2 \square R1$$
 o $R1 \square R2 \square R2$

25. Give the uses of Multivalued dependency?.

We use multivalued dependencies in two ways:

- To test relations to determine whether they are legal under a given set of functional and multivalued dependencies
- To specify constraints on the set of legal relations. We shall thus concern ourselves *only* with relations that satisfy a given set of functional and multivalued dependencies.

PART B

- Explain the following with relevant examples.

 B-Tree ii. B+Tree iii.Static and dynamic hashing.
- 2. Write a relevant example discuss the steps involved in processing a query.
- 3. Construct a B+ tree to insert the following key elements (order of the tree is 3) 5, 3, 4, 9, 7, 15, 14,21,22, 23.
- 4. Describe in detail about how records are represented in a file and how to Organize them in a file.
- 5. Write about the file organization technique and type.
- 6. Explain the followingi) B+- treeii) Dynamic hashing & static hashing
- 7. Mention the purpose of indexing. How this can be done by B+ tree? Explain.
- 8. Explain following with relevant examples:
 - (i) B tree
 - (ii) B+Tree
 - (iii) Static and dynamic hashing
- 9. With a relevant example discuss the steps involved in processing a query.

DATA BASE MANAGEMENT SYSTEMS

QUESTION BANK

UNIT IV

PART A

1. State the atomicity property of a transaction.

Atomicity: Either all operations of the transaction are properly reflected in the database or none.

Correctness: Execution of a transaction in isolation preserves the consistency **Isolation**.: Each transaction must be unaware of other concurrently executing transactions. **Durability:** After a transaction completes successfully, the changes it has made to the database persist, even if there are system failures.

2. Define deadlock.

A transaction waits for another transaction to release lock on data item. Mean while, transaction also waits for some other data item to be released by first transaction.

3. When are two schedules conflict equivalent?

If a schedule S can be transformed into a schedule S' by a series of swaps of nonconflicting instructions, we say that S and S' are conflict equivalent.

4. State the benefits of strict two-phase locking.

Cascading rollbacks can be avoided by a modification of two-phase locking. Any data written by an uncommitted transaction are locked in exclusive mode until the transaction commits, preventing any other transactions from reading the data.

5. What benefit is provided by strict-two-phase locking? What are the disadvantages results?

The transactions can be serialized in the order of their lock points. In strict two-phase locking, transaction must hold all its exclusive locks till it commits/aborts.

6. What is concurrency control?

Concurrency control is maintaining consistency

7. What is transaction?

Collections of operations that form a single logical unit of work are called transactions.

8. What are the two statements regarding transaction?

The two statements regarding transaction of the form: _ Begin transaction End transaction 9. When is a transaction rolled back?

Any changes that the aborted transaction made to the database must be undone. Once the changes caused by an aborted transaction have been undone, then the transaction has been rolled back.

10. What are the states of transaction?

The states of transaction are

- Active
- Partially committed
- Failed
- Aborted
- Commited
- Terminated

11. What is a shadow copy scheme?

It is simple, but efficient, scheme called the shadow copy schemes. It is based on making copies of the database called shadow copies that one transaction is active at a time. The scheme also assumes that the database is simply a file on disk.

12. Give the reasons for allowing concurrency?

The reasons for allowing concurrency is if the transactions run serially, a short transaction may have to wait for a preceding long transaction to complete, which can lead to unpredictable delays in running a transaction.

So concurrent execution reduces the unpredictable delays in running transactions.

13. What is average response time?

The average response time is that the average time for a transaction to be completed after it has been submitted.

14. What are the two types of serializability?

The two types of serializability is _

Conflict serializability

_ View serializability

15. Define lock?

Lock is the most common used to implement the requirement is to allow a transaction to access a data item only if it is currently holding a lock on that item.

16. What are the different modes of lock?

The modes of lock are: _ Shared

_ Exclusive

17. Differentiate strict two phase locking protocol and rigorous two phase locking protocol.

In **strict two phase locking protocol** all exclusive mode locks taken by a transaction is held until that transaction commits.

Rigorous two phase locking protocol requires that all locks be held until the transaction commits.

18. How the time stamps are implemented

- Use the value of the system clock as the time stamp. That is a transaction's time stamp is equal to the value of the clock when the transaction enters the system.

-Use a logical counter that is incremented after a new timestamp has been assigned; that is the time stamp is equal to the value of the counter.

19. What are the time stamps associated with each data item?

-W-timestamp (Q) denotes the largest time stamp if any transaction that executed WRITE (Q) successfully.

- R-timestamp (Q) denotes the largest time stamp if any transaction that executed READ (Q) successfully.

20. What is recovery management component?

Ensuring durability is the responsibility of a software component of the base system called the recovery management component.

21. Define the phases of two phase locking protocol

_ Growing phase: a transaction may obtain locks but not release any lock.

_ Shrinking phase: a transaction may release locks but may not obtain any new locks.

22. What is meant by log-based recovery?

The most widely used structures for recording database modifications is the log.

The log is a sequence of log records, recording all the update activities in the database. There are several types of log records.

23. What are the two methods for dealing deadlock problem?

The two methods for dealing deadlock problem is deadlock detection and deadlock recovery.

24. What are the two types of errors?

The two types of errors are: _ Logical error

_ System error

25. What are the storage types?

The storage types are: _ Volatile storage _ Nonvolatile storage

PART B

1. Explain testing for serializability with respect to concurrency control schemes. How will you determine, whether a schedule is serialiazble or not.

2. Explain the following protocols for concurrency control.

i.Lock based protocols

ii.Time stamp based protocols.

3. Discuss in detail about Transaction Recovery, System Recovery and Media Recovery

4. Write down in detail about Deadlock and Serializability.

5. Explain about

i)Concurrent execution

ii)Recoverability

6. Explain various recovery techniques during transaction

7.(i) Explain about immediate update and deferred update recovery techniques.

- (ii) Explain the concepts of serializability.
- 8.(i) Explain Two-phase locking protocol.
- (ii) Describe about the deadlock prevention schemes
- 9.(i)How can you implement atomicity in transactions? Explain.
- (ii)Describe the concept of serilalizability with suitable example.
- 10. How concurrency is performed? Explain the protocol that is used to maintain the concurrency concept.
- 11. Explain testing for serializability with respect to concurrency control schemes. How will you determine, whether a schedule is serializable or not.
- 12. Explain the following concurrency control:
 - (i) Lock based protocol
 - (ii) Time stamp based protocol
- 13.Explain briefly about the Deadlock.
- 14. Explain in details about Database Recovery techniques
- 15. What is Transaction Processing? Explain in detail about concept, properties and State?
- 16. Explain ACID in detail.
- 17.Explain serializability.
- 18.Explain in details about Log-Based Recovery.

DATA BASE MANAGEMENT SYSTEMS

QUESTION BANK

UNIT V

PART A

1) What are ordered Indices?

In an **ordered index**, index entries are stored sorted on the search key value. E.g., author catalog in library.

Two types: Sparse Index and Dense index.

2) Distinguish between sparse index and dense index.

An index record appears for every search key value in the file is dense index. An index record appears for only some of the search key values is sparse index. Each index record contains a search key value and a pointer to the first data record with that search value.

3. What is a heap file? How are pages organized in a heap file?

Heap file:- Any record can be placed anywhere in the file where there is space for the record. There is no ordering of records.

4. How does a B-tree differ from B+ - tree? Why is a B+ - tree usually preferred as an access structure to a data file?

B-tree allows search-key values to appear only once; eliminates redundant

storage of search keys. Search keys in nonleaf nodes appear nowhere else in the B-tree; All paths from root to leaf are of the same length. Each node that is not a root or a

leaf has between $\Box \widetilde{\Box}$.and *n* children. A leaf node has between (n-1)/2.and *n*-1 values.

5. Define Static hashing. Static Hashing:

Hash function h is a function from the set of all search-key values K to the set of all bucket addresses B. Hash function is used to locate records for access, insertion as well as deletion. In static hashing, function h maps search-key values to a fixed set of B of bucket addresses.

6. Define Dynamic hashing.

Dynamic Hashing: Good for database that grows and shrinks in size, allows the hash function to be modified dynamically. The number of buckets also changes dynamically due to

coalescing and splitting of buckets.

7. How does a DBMS represent a relational query plan?

Parsing and translation: Translate the query into its internal form. This is then translated into relational algebra. Parser checks syntax, verifies relations **Evaluation**: The query-execution engine takes a query-evaluation plan, executes that plan, and returns the answers to the query.

8. Define RAID:

Redundant Arrays of Independent Disks.

9. Compare sequential access devices versus random access devices with an example.

Records are stored in sequential order according to the value of a search key. If n th record has to be accessed then all n-1 records has to be passed eg. Tape. In random access device, n th record can be pointed directly eg. Disk.

10. What can be done to reduce the occurrences of bucket overflows in a hash file organization?

The bucket overflows can be reduced by using overflow buckets. Overflow handling can be done using linked list. This is called overflow chaining.

11. Give the measures of quality of a disk.

_ Capacity

_ Access time _ Seek time

_ Data transfer rate _ Reliability

_ Rotational latency time.

12. Compare sequential access devices versus random access devices with an example sequential access devices random access devices

Must be accessed from the beginning It is possible to read data from any location Eg:tape storage Eg:-disk storage

Access to data is much slower Access to data is faster Cheaper than disk Expensive when compared with disk

13. What are the types of storage devices?

_ Primary storage

_ Secondary storage _

Tertiary storage

14. Define access time.

Access time is the time from when a read or write request is issued to when data transfer begins.

15. Define seek time.

The time for repositioning the arm is called the seek time and it increases with the distance that the arm is called the seek time.

16. Define average seek time.

The average seek time is the average of the seek times, measured over a sequence of random requests.

17. Define rotational latency time.

The time spent waiting for the sector to be accessed to appear under the head is called the rotational latency time.

18. Define average latency time.

The average latency time of the disk is one-half the time for a full rotation of the disk.

19. What is meant by data-transfer rate?

The data-transfer rate is the rate at which data can be retrieved from or stored to the disk.

20. What is meant by mean time to failure?

The mean time to failure is the amount of time that the system could run continuously without failure.

21. What are a block and a block number?

A block is a contiguous sequence of sectors from a single track of one platter.

Each request specifies the address on the disk to be referenced. That address is in the form of a block number.

22. What is the use of RAID?

A variety of disk-organization techniques, collectively called redundant arrays of independent disks are used to improve the performance and reliability.

23. Explain how reliability can be improved through redundancy?

The simplest approach to introducing redundancy is to duplicate every disk. This technique is called mirroring or shadowing. A logical disk then consists of two physical disks, and write is carried out on both the disk. If one of the disks fails the data can be read from the other. Data will be lost if the second disk fails before the first fail ed disk is repaired.

24. What is called bit-level striping?

Data striping consists of splitting the bits of each byte across multiple disks. This is called bit-level striping.

25. What is called block-level striping?

Block level striping stripes blocks across multiple disks. It treats the array of disks as a large disk, and gives blocks logical numbers.

26. Distinguish between fixed length records and variable length records? Fixed length records

Every record has the same fields and field lengths are fixed.

Variable length records

File records are of same type but one or more of the fields are of varying size.

27. What are the ways in which the variable-length records arise in database systems?

_ Storage of multiple record types in a file.

_ Record types that allow variable lengths for one or more fields. _ Record types that allow repeating fields.

28. Explain the use of variable length records.

_ They are used for Storing of multiple record types in a file.

_ Used for storing records that has varying lengths for one or more fields. _ Used for storing records that allow repeating fields

29. What is known as heap file organization?

In the heap file organization, any record can be placed anywhere in the file where there is space for the record. There is no ordering of records. There is a single file for each relation.

30. What is known as sequential file organization?

In the sequential file organization, the records are stored in sequential order, according to the value of a "search key" of each record.

31. What is hashing file organization?

In the hashing file organization, a hash function is computed on some attribute of

each record. The result of the hash function specifies in which block of the file the record should

be placed.

32. What is known as clustering file organization?

In the clustering file organization, records of several different relations are stored in the same file.

33. What is an index?

An index is a structure that helps to locate desired records of a relation quickly, without examining all records.

34. What are the two types of ordered indices?

_ Primary index

_ Secondary index

35. What are the types of indices?

_ Ordered indices _ Hash indices

36. What are the techniques to be evaluated for both ordered indexing and hashing?

_ Access types _ Access time _ Insertion time _ Deletion time

_ Space overhead

37. What is linear probing?

Linear probing is a type of open hashing. If a bucket is full the system inserts records in to the next bucket that has space. This is known as linear probing.

38. What is called query processing?

Query processing refers to the range of activities involved in extracting data from a database.

39. What are the steps involved in query processing?

The basic steps are:

_ parsing and translation _optimization_ evaluation40. What is called an evaluation primitive?

A relational algebra operation annotated with instructions on how to evaluate is called an evaluation primitive.

41. What is called a query evaluation plan?

A sequence of primitive operations that can be used to evaluate ba query is a query evaluation plan or a query execution plan.

42. What is called a query -execution engine?

The query execution engine takes a query evaluation plan, executes that plan, and returns the answers to the query.

43. How do you measure the cost of query evaluation?

The cost of a query evaluation is measured in terms of a number of different resources including disk accesses, CPU time to execute a query, and in a distributed database system the cost of communication

44. List out the operations involved in query processing

Selection operation Join operations. Sorting.

Projection Set operations Aggregation

45. Define query optimization.

Query optimization refers to the process of finding the lowest –cost method of evaluating a given query.

46. Which level of RAID is best? Why?

RAID level 1 is the RAID level of choice for many applications with moderate

storage requirements and high I/O requirements. RAID 1 follows mirroring and provides best write performance.

PART B

1. State and explain the features of object oriented data model. Use banking application as an

example.

- 2. Write detailed note on the following:
 - i. Distributed data base.

ii.Data mining.

- 3. Write detailed on the following
 - i) Distributed database

ii)Data mining

- 4. Discuss in detail about the object relational database and its advantage
- 5. State and explain the object oriented data model. Use banking application as an example. (16)
- 6. Write detail notes on following:
 - (i) Distributed Databases
 - (ii) Data Mining
- 7. Explain briefly about Inheritance and Multiple Inheritance.
- 8.Discuss in detail about Steps for Data Mining.
- 9. Discuss in detail about Data Warehousing.
- 10. Explain briefly about the Querying and Transformation.