



KARPAGAM ACADEMY OF HIGHER EDUCATION
(Deemed to be University Established Under Section 3 of UGC Act 1956)
Pollachi Main Road, Eachanari (Po),
Coimbatore –641 021
DEPARTMENT OF MATHEMATICS

SUBJECT :REAL ANALYSIS

SUBJECT CODE :17MMP102

L T P C

4 0 0 4

PO: After the completion of this course, the learner will be enriched with the concept of analysis which is the motivating tool in other area such as applied mathematics.

PLO: To understand the Riemann – Stieltjes Integral, Infinite series, infinite products, Sequences of functions, the Lebesgue integral, Implicit functions and extremum problems to have a sound knowledge in Measure Theory.

UNIT I

The Riemann – Stieltjes Integral:

Introduction – Basic Definitions – Linear Properties – Integration by parts – Change of variable in a Riemann – Stieltjes Integral – Reduction to a Riemann Integral – Step functions as integrators – Reduction of a Riemann – Stieltjes Integral to a finite sum – Monotonically increasing – Additive and linear properties – Riemann condition – Comparison theorems – Integrators of bounded variation – Sufficient condition for Riemann Stieltjes integral.

UNIT II

Infinite series and infinite products:

Introduction – Basic definitions – Ratio test and root test – Dirichlet test and Able's test Rearrangement of series – Riemann's theorem on conditionally convergent series – Sub series - Double sequences – Double series – Multiplication of series – Cesaro summability.

UNIT III

Sequences of functions:

Basic definitions – Uniform convergence and continuity - Uniform convergence of infinite series of functions – Uniform convergence and Riemann – Stieltjes integration – Non uniformly convergent sequence – Uniform convergence and differentiation – Sufficient condition for uniform convergence of a series.

UNIT IV

The Lebesgue integral:

Introduction- The class of Lebesgue – integrable functions on a general interval- Basic properties of the Lebesgue integral- Lebesgue integration and sets of measure zero- The Levi monotone convergence theorem- The Lebesgue dominated convergence theorem- Applications of Lebesgue dominated convergence theorem- Lebesgue integrals on unbounded intervals as limit of integrals on bounded intervals- Improper Riemann integrals- Measurable functions.

UNIT V

Implicit functions and extremum problems:

Introduction – Functions with non zero Jacobian determinant – Inverse function theorem – Implicit function theorem – Extrema of real valued functions of one variable and several variables.

SUGGESTED READINGS**TEXT BOOK**

1. Rudin. W., (1976) .Principles of Mathematical Analysis, Mcgraw Hill, New york .

REFERENCES

1. Tom .M. Apostol., (2002). Mathematical Analysis, Second edition, Narosa Publishing House, New Delhi.
2. Balli. N.P., (1981). Real Analysis, Laxmi Publication Pvt Ltd, New Delhi.
3. Gupta.S.L. and Gupta.N.R.,(2003).Principles of Real Analysis, Second edition, Pearson Education Pvt.Ltd, Singapore.
4. Royden .H.L., (2002). Real Analysis, Third edition, Prentice hall of India,New Delhi.
5. Sterling. K. Berberian., (2015).A First Course in Real Analysis, Springer Pvt Ltd, New Delhi.



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DEPARTMENT OF MATHEMATICS

Lecture Plan

Subject: Real Analysis
Sub. Code: 17MMP102

Class: I M.Sc Mathematics
Name of the faculty: S.KOHILA

S. No.	Lecture Duration (Hours)	Topics to be covered	Support Materials
UNIT – I			
1.	1	Introduction on Riemann-Stieltjes integral	T1: Chap: 6, Pg. No : 120-124
2.	1	Some basic definitions on Riemann-Stieltjes integral	T1: Chap: 6, Pg. No : 124-128
3.	1	Linear properties of Riemann-Stieltjes integral	T1: Chap: 6, Pg. No : 128-129
4.	1	Continuation on linear properties of Riemann-Stieltjes integral	T1: Chap: 6, Pg. No : 129-131
5.	1	Integration by parts	T1: Chap: 6, Pg. No : 134
6.	1	Change of variable in Riemann-Stieltjes integral	T1: Chap: 6, Pg. No : 132-133
7.	1	Reduction to Riemann integral	R1: Chap: 7, Pg. No : 145-146
8.	1	Step function as integrators	R1: Chap: 7, Pg. No : 147-148
9.	1	Reduction of a Riemann-Stieltjes integral to a finite sum	R1: Chap: 7, Pg. No : 148-149
10.	1	Monotonically increasing on Riemann integral	R1: Chap: 7, Pg. No : 150-152
11.	1	Additive and linear properties on Riemann	R1: Chap: 7, Pg. No : 153

		integral	
12.	1	Riemann condition	R1: Chap: 7, Pg. No : 154
13.	1	Comparison theorems & Integrators of bounded variation	R1: Chap: 7, Pg. No : 155-156
14.	1	Sufficient condition for Riemann stieltjes integral	R1: Chap: 7, Pg. No : 156-158
15.	1	Recapitulation and discussion of important questions on unit I	R1: Chap: 7, Pg. No : 159
Total	15 Hours		

T1. Rudin. W.,1976 .Principles of mathematical Analysis, Mcgraw hill, Newyork .

R1. Tom .M. Apostol .,2002. Mathematical Analysis, Second edition, Narosa Publishing House,New Delhi.

UNIT – II

1.	1	Introduction on infinite series & infinite products	R3: Chap: 5, Pg. No : 5.1-5.3
2.	1	Some basic definitions on infinite series and infinite products	R1: Chap: 8, Pg. No : 183-191
3.	1	Ratio test and root test	R1: Chap: 8, Pg. No : 193-194
4.	1	Drichlet test and	R3: Chap: 6, Pg. No : 6.1-6.4
5.	1	Able's test	R1: Chap: 8, Pg. No : 196
6.	1	Rearrangement of Series	R1: Chap: 8, Pg. No : 197
7.	1	Riemann's Theorem on conditionally convergent series	R1: Chap: 8, Pg. No : 197-199
8.	1	Sub series on conditionally convergent series Double sequences on	R1: Chap: 8, Pg. No : 199-200

9.	1	conditionally convergent series	R1: Chap: 8, Pg. No :200-202
10.	1	Double series on conditionally convergent series	R1: Chap: 8, Pg. No :203-205
11.	1	Problems on double series on conditionally convergent series	R1: Chap: 8, Pg. No :205-206
12.	1	Multiplication of series	R1: Chap: 8, Pg. No :206-208
13	1	conditionally convergent series	R1: Chap: 8, Pg. No :209-210
14	1	Cesaro Summability	R1: Chap: 8, Pg. No :11-212
15	1	Recapitulation and discussion of important questions on unit II	
Total	15 Hours		

R1. Tom .M. Apostol .,2002. Mathematical Analysis, Second edition, Narosa Publishing House,New Delhi.

R3. Gupta . S.L ., and N.R. Gupta ., 2003.Principles of Real Analysis, Second edition, Pearson Education Pvt.Ltd,Singapore.

UNIT –III

1.	1	Basic definitions on sequences	R5: Chap:3, Pg. No :33-36
2.	1	Basic definitions on of functions	R5: Chap:3, Pg. No :39-41
3.	1	Uniform convergence on sequences of functions	T1: Chap:7, Pg. No :143-146
4.	1	Continuity of uniform convergence on sequences of functions	T1: Chap:7, Pg. No :147-148
5.	1	Uniform convergence of infinite series	R2: Chap:9, Pg. No :533-534

6.	1	Problems on infinite series of functions	T1: Chap:7, Pg. No :152-153
7.	1	Riemann-Stieltjes integration	T1: Chap:7, Pg. No :152-154
8.	1	Non-uniformly convergent sequence	R1: Chap:9, Pg. No :228-229
9.	1	Non-uniformly convergent sequence.	R1: Chap:9, Pg. No :230-231
10.	1	Problems on non-uniformly convergent sequence	
11	1	Problems on non-uniformly convergent sequence	R1: Chap:9, Pg. No :231
12	1	Problems on non-uniformly convergent sequence .	R1: Chap:9, Pg. No :232
13	1	Necessary condition for uniform convergence of a series	R1: Chap:9, Pg. No :235
14	1	Sufficient condition for uniform convergence of a series	R1: Chap:9, Pg. No :235
15	1	Recapitulation and discussion of important questions on unit III	
Total	15 Hours		

T1. Rudin. W.,1976 .Principles of mathematical Analysis, Mcgraw hill, Newyork .

R1. Tom .M. Apostol .,2002. Mathematical Analysis, Second edition, Narosa Publishing House,New Delhi.

R5. Sterling. K. Berberian ., 2004.A First Course in Real Analysis, Springer Pvt Ltd, New Delhi.

UNIT-IV

1.	1	Introduction on Lebesgue integral	R4:Chap:4,Pg.No:75-77
2.	1	The class of Lebesgue – integrable functions	R1:Chap:10,Pg.No:254-

		on a general interval	256
3.	1	Basic properties of a Lebesgue integral	R4:Chap:4,Pg.No:85-88
4.	1	Lebesgue integration and sets of measure zero	R1:Chap:10,Pg.No:264-265
5.	1	The Levi monotone convergence theorem	R1:Chap:10,Pg.No:266-268
6.	1	The Lebesgue dominated convergence theorem	R1:Chap:10,Pg.No:268-270
7.	1	Applications of Lebesgue dominated convergence theorem	R1:Chap:10,Pg.No:270-272
8.	1	Continuation of applications of Lebesgue dominated convergence theorem	R1:Chap:10,Pg.No:272-274
9.	1	Lebesgue integrals on unbounded interval as limit of integrals on bounded intervals	R1:Chap:10,Pg.No:274-275
10.	1	Improper Riemann integrals	R1:Chap:10,Pg.No:276-279
11.	1	Measurable functions	R1:Chap:10,Pg.No:279-280
12.	1	Problems on Improper Riemann integrals	R1:Chap:10,Pg.No:278
13	1	Problems on Improper Riemann integrals	R1:Chap:10,Pg.No:282
14	1	Problems on Measurable functions	R1:Chap:10,Pg.No:279-280
15		Recapitulation and discussion of important questions on unit IV	
Total	15 Hours		
R1. Tom .M. Apostol .,2002. Mathematical Analysis, Second edition, Narosa Publishing House,New Delhi.			
UNIT – V			

1.	1	Introduction on implicit functions	R1:Chap:13,Pg.No:367-368
2.	1	Functions with non-zero Jacobian determinant	R1:Chap:13,Pg.No:368-371
3.	1	Inverse function theorem	R1:Chap:13,Pg.No:372-373
4.	1	Implicit function theorem	R1:Chap:13,Pg.No:373-375
5.	1	Extrema of real valued functions of one variable	R1:Chap:13,Pg.No:375-376
6.	1	Extrema of real valued functions of several variables-Theorems	R1:Chap:13,Pg.No:376-379
7.	1	Extrema of real valued functions of several variables -Theorems	R1:Chap:13.Pg No : 380-381
8.	1	Problems on Extrema of real valued functions of several variables	R1:Chap:13.Pg No : 380
9.	1	Problems on Inverse function theorem	R1:Chap:13.Pg No : 380
10.	1	Problems on Inverse function theorem	R1:Chap:13.Pg No : 376
11	1	Problems on Inverse function theorem	R1:Chap:13.Pg No : 378
12	1	Recapitulation and discussion of important questions on unit V	
13	1	Discussion of previous ESE question papers.	
14	1	Discussion of previous ESE question papers.	
15	1	Discussion of previous ESE question papers.	
Total	15 Hours		
R1. Tom .M. Apostol .,2002. Mathematical Analysis, Second edition, Narosa Publishing House,New Delhi.			

Text book and References:

T1. Rudin. W.,1976 .Principles of mathematical Analysis, Mcgraw hill, Newyork .

R1. Tom .M. Apostol .,2002. Mathematical Analysis, Second edition, Narosa Publishing House,New Delhi.

R2. Balli. N.P., 1981. Real Analysis, Laxmi Publication Pvt Ltd, New Delhi.

R3. Gupta . S.L ., and N.R. Gupta ., 2003.Principles of Real Analysis, Second edition, Pearson Education Pvt.Ltd,Singapore.

R4. Royden .H.L ., 2002. Real Analysis, Third edition, Prentice hall of India,New Delhi.

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Subject : REAL ANALYSIS
SUBJECT CODE: 17MMP102

SEMESTER: I
CLASS : I M.Sc

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UNIT I

The Riemann – Stieltjes Integral:

Introduction – Basic Definitions – Linear Properties – Integration by parts – Change of variable in a Riemann – Stieltjes Integral – Reduction to a Riemann Integral – Step functions as integrators – Reduction of a Riemann – Stieltjes Integral to a finite sum – Monotonically increasing – Additive and linear properties – Riemann condition – Comparison theorems – Integrators of bounded variation – Sufficient condition for Riemann Stieltjes integral.

TEXT BOOK

1. Rudin. W., (1976) .Principles of Mathematical Analysis, Mcgraw Hill, New york .

REFERENCES

1. Tom .M. Apostol., (2002). Mathematical Analysis, Second edition, Narosa Publishing House, New Delhi.
2. Balli. N.P., (1981). Real Analysis, Laxmi Publication Pvt Ltd, New Delhi.
3. Royden .H.L., (2002). Real Analysis, Third edition, Prentice hall of India, New Delhi.

1.1. The Riemann-Stieltjes Integral.

Definitions:

Let $[a, b]$ be a given interval. Then a set $P = \{x_0, x_1, \dots, x_{n-1}, x_n\}$ of $[a, b]$ such that $a = x_0 \leq x_1 \leq \dots \leq x_{n-1} \leq x_n = b$ is said to be a Partition of $[a, b]$. The set of all partitions of $[a, b]$ is denoted by $P([a, b])$. The intervals $[x_0, x_1], [x_1, x_2], \dots, [x_{n-1}, x_n]$ are called the subintervals of $[a, b]$. Write $\Delta x_i = x_i - x_{i-1}$ is called the length of the interval $[x_{i-1}, x_i]$ ($i = 1, \dots, n$) and $\max |\Delta x_i|$ is called the norm of the partition P and is denoted by $\|P\|$ or Q is called the refinement or finer of the partition $P \subset (P)$. A partition Q of $[a, b]$ such that $P \subset Q$ is called the refinement or finer of the partition $P \subset (P)$. A partition Q of $[a, b]$ such that $P \subset Q$ is called the refinement or finer of the partition $P \subset (P)$. Suppose f is a bounded real valued function defined on $[a, b]$ and $P \subset ([a, b])$. Then $M_i = \sup_{x_{i-1} \leq x \leq x_i} f(x)$, $m_i = \inf_{x_{i-1} \leq x \leq x_i} f(x)$ for each P

Suppose f is a bounded real valued function defined on $[a, b]$ and $P \subset ([a, b])$. Then

$M_i = \sup_{x_{i-1} \leq x \leq x_i} f(x)$, $m_i = \inf_{x_{i-1} \leq x \leq x_i} f(x)$ for each P $\sum_{i=1}^n M_i \Delta x_i$ and $L(P, f) = \sum_{i=1}^n m_i \Delta x_i$ are called the Upper and Lower Riemann sums $\sum_{i=1}^n M_i \Delta x_i$ and $L(P, f) = \sum_{i=1}^n m_i \Delta x_i$ or Upper and Lower Darboux sums of f on $[a, b]$ with respect to the partition P .

Further write $\int_a^b f(x) dx = \inf U(P, f)$ and $\int_a^b f(x) dx = \sup L(P, f)$ where the inf and the sup are taken over all partitions P of $[a, b]$ are called the Upper and Lower Riemann integrals of f over $[a, b]$, respectively.

If the upper and lower Riemann integrals are $R[a, b]$ and we denote $\int_a^b f(x) dx$ equal, we say that f is Riemann-integrable on $[a, b]$ and we write f the common value of these integrals by $\int_a^b f(x) dx$, i.e., $\int_a^b f(x) dx = \int_a^b f(x) dx = \int_a^b f(x) dx$.

R is bounded function then the upper and lower Riemann \rightarrow 1.1.1. Lemma . If $f : [a, b]$ integrals of f are bounded. Since f is bounded, there exist two numbers m and M such that $m \leq f(x) \leq M$ ($a \leq x \leq b$). Hence, for every partition P of $[a, b]$ we have $M \leq M_i \leq m_i \leq m$ $\sum_{i=1}^n M_i \Delta x_i \leq \sum_{i=1}^n m_i \Delta x_i \leq \sum_{i=1}^n m \Delta x_i \leq \sum_{i=1}^n M \Delta x_i \leq M(b-a)$, $i = 1, 2, 3, \dots, n$. $m(b-a) \leq L(P, f) \leq U(P, f) \leq M(b-a)$, \Rightarrow so that the numbers $L(P, f)$ and $U(P, f)$ form a bounded set. Therefore by the definition of lower and upper Riemann integrals this shows that the upper and lower integrals are defined for every bounded function f are bounded also. The question of their equality, and hence the question of the integrability of f ,

R is bounded function, P is any partition of $[a, b]$ and P^* is the \rightarrow

1.1.2. Lemma. If $f : [a, b]$ refinement of P , then $L(P, f) \leq L(P^*, f)$ and $U(P^*, f) \leq U(P, f)$. R is bounded function and P_1, P_2 are any two partitions of $[a, b] \rightarrow$

1.1.3. Lemma. f is bounded function and P_1, P_2 are any two partitions of $[a, b]$ → If $f : [a, b]$

$$L(P_1, f) \leq U(P_2, f) \text{ and } L(P_2, f) \leq U(P_1, f).$$

f, g are bounded functions and P is any partition of $[a, b]$ then → 1.1.4.

Lemma. If $f, g : [a, b]$ (i) $L(P, f + g) \geq L(P, f) + L(P, g)$ (ii) $U(P, f + g) \leq U(P, f) + U(P, g)$. f, g are bounded functions.

Theorem. If $f : [a, b] \rightarrow \mathbb{R}$ is bounded function then for $\epsilon > 0$ there exists $\delta > 0$ such that $U(P, f) - L(P, f) < \epsilon$ if $\|P\| < \delta$. f is Riemann Integrable if the oscillatory $\omega(f, \delta) \rightarrow 0$ as $\delta \rightarrow 0$.

1.1.3. Theorem. If $f : [a, b] \rightarrow \mathbb{R}$ is bounded function and $\omega(f, \delta) \rightarrow 0$ as $\delta \rightarrow 0$, then f is Riemann Integrable.

1.1.4. Theorem. Every continuous function $f : [a, b] \rightarrow \mathbb{R}$ is Riemann Integrable. → 1.1.5. Theorem. Every monotone function $f : [a, b] \rightarrow \mathbb{R}$ is Riemann Integrable. Students you studied the properties given above and other properties of Riemann Integrals in previous classes therefore we are not interested to investigate these here. However we shall immediately consider a more general situation. Let $f : [a, b] \rightarrow \mathbb{R}$ be a monotonically increasing α -Riemann bounded function and →

1.1.2 Definition. Let $f : [a, b] \rightarrow \mathbb{R}$ be a function on $[a, b]$. Let $P = \{x_0, x_1, \dots, x_{n-1}, x_n\}$ such that $a = x_0 \leq x_1 \leq \dots \leq x_{n-1} \leq x_n = b$ be any Partition of $[a, b]$. We write (x_{i-1}, x_i) , $i = 1, 2, 3, \dots, n$. $\alpha(x_i) - \alpha(x_{i-1}) = \Delta\alpha_i$ is bounded on $[a, b]$, $\Delta\alpha_i$ are finite therefore $\alpha(a)$ and $\alpha(b)$ are finite. By the definition of monotone function $\Delta\alpha_i \geq 0$, $i = 1, 2, 3, \dots, n$. $\Delta\alpha$ is monotonically increasing function then clearly α is also since $P \subset [a, b]$. We define $M_i = \sup_{x \in [x_{i-1}, x_i]} f(x)$, $m_i = \inf_{x \in [x_{i-1}, x_i]} f(x)$ for each P . $U(P, f, \Delta\alpha) = \sum_{i=1}^n M_i \Delta\alpha_i$ and $L(P, f, \Delta\alpha) = \sum_{i=1}^n m_i \Delta\alpha_i$ are called the Upper and Lower Riemann Stieltjes sums respectively. Further we define $U(f, \alpha) = \sup_P U(P, f, \Delta\alpha)$ and $L(f, \alpha) = \sup_P L(P, f, \Delta\alpha)$. $\int_a^b f d\alpha = U(f, \alpha) = L(f, \alpha)$ where the inf and the sup are taken over all partitions P of $[a, b]$, are called the Upper and Lower Riemann Stieltjes integrals of f over $[a, b]$, respectively.

If the upper and lower Riemann Stieltjes integrals are equal, we say that f is Riemann Stieltjes integrable on $[a, b]$

Lower Riemann Stieltjes integrals of f over $[a, b]$, respectively. If the upper and lower Riemann Stieltjes integrals are equal, we say that f is Riemann Stieltjes integrable on $[a, b]$

$$\int_a^b f d\alpha \text{ or } \int_a^b f(x) d\alpha$$

E.2. PROPERTIES

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as $n \rightarrow \infty$, we have

$$\int_0^{10} f(x) d\alpha(x) = 50 + 55 = 105.$$

E.2. Properties

Theorem E.4. Let c_1, c_2 be two constants in \mathbb{R} .

(1) If $f, g \in R(\alpha)$ on $[a, b]$, then $c_1f + c_2g \in R(\alpha)$ on $[a, b]$, and

$$\int_a^b (c_1f + c_2g) d\alpha = c_1 \int_a^b f d\alpha + c_2 \int_a^b g d\alpha.$$

(2) If $f \in R(\alpha)$ and $f \in R(\beta)$ on $[a, b]$, then $f \in R(c_1\alpha + c_2\beta)$ on $[a, b]$, and

$$\int_a^b f d(c_1\alpha + c_2\beta) = c_1 \int_a^b f d\alpha + c_2 \int_a^b f d\beta.$$

(3) If $c \in [a, b]$, then

$$\int_a^b f d\alpha = \int_a^c f d\alpha + \int_c^b f d\alpha.$$

Definition E.5. If $a < b$, we define

$$\int_b^a f d\alpha = - \int_a^b f d\alpha.$$

Theorem E.6. If $f \in R(\alpha)$ and α has a continuous derivative on $[a, b]$, then the Riemann integral $\int_a^b f(x)\alpha'(x) dx$ exists and

$$\int_a^b f(x) d\alpha(x) = \int_a^b f(x)\alpha'(x) dx.$$

E.3.1. Integration by parts.

Theorem E.7 (Integration by parts). *If $f \in R(\alpha)$ on $[a, b]$, then $\alpha \in R(f)$ on $[a, b]$,*

and

$$\int_a^b f(x) d\alpha(x) = f(b)\alpha(b) - f(a)\alpha(a) - \int_a^b \alpha(x) df(x).$$

Example E.8. As in Example E.3, $f(x) = x$, and $\alpha(x) = x + [x]$. Then

$$\begin{aligned} \int_0^{10} f(x) d\alpha(x) &= f(10)\alpha(10) - f(0)\alpha(0) - \int_0^{10} \alpha(x) df(x) \\ &= 10 \times 20 - 0 \times 0 - \int_0^{10} (x + [x]) dx \\ &= 200 - 50 - \int_0^{10} [x] dx = 150 - 45 = 105 \end{aligned}$$

E.3.2. Change of variables.

Theorem E.9 (Change of variables). *Suppose that $f \in R(\alpha)$ on $[a, b]$ and g is a strictly increasing continuous function on $[c, d]$ with $a = g(c)$, $b = g(d)$. Let $h = f \circ g$, $\beta = \alpha \circ g$. Then $h \in R(\beta)$ on $[c, d]$ and*

$$\int_a^b f(x) d\alpha(x) = \int_c^d f(g(t)) d\alpha(g(t)) = \int_c^d h(t) d\beta(t).$$

Example E.10. Let $y = \sqrt{x}$, we have

$$\begin{aligned} \int_0^4 ([\sqrt{x}] + x^2) d\sqrt{x} &= \int_0^2 ([y] + y^4) dy = \int_0^2 [y] dy + \int_0^2 y^4 dy \\ &= 1 + \frac{1}{5} y^5 \Big|_{y=0}^2 = \frac{37}{5} \end{aligned}$$

E.3. TECHNIQUE OF INTEGRATIONS

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E.3.3. Step functions as α . By Remark C.6 and Theorem E.4(2), we have

$$\int_a^b f(x) dF(x) = \int_a^b f(x) dF_{ac}(x) + \int_a^b f(x) dF_{sc}(x) + \int_a^b f(x) dF_d(x) \quad (\text{E.2})$$

Remark E.11. If $\alpha \equiv \text{constant}$ on $[a, b]$, then $S(P, f, \alpha) = 0$ for all partition P , and

$$\int_a^b f(x) d\alpha(x) = 0.$$

Example E.15. (1) Consider

$$f(x) = 1 \quad \text{for } x \in [-1, 1], \quad \text{and} \quad \alpha(x) = -I_{\{0\}},$$

then

$$\int_{-1}^1 f(x) d\alpha(x) = f(0)(\alpha(0+) - \alpha(0-)) = 0$$

(2) Consider

$$f(x) = 2I_{\{0\}} + I_{[-1,0) \cup (0,1]} \quad \text{and} \quad \alpha(x) = -I_{[0,1]}.$$

Then both of α and f are discontinuous from the left at $x = 0$. This implies that the Riemann-Stieltjes integral $\int_{-1}^1 f d\alpha$ does not exist.

E.3. TECHNIQUE OF INTEGRATION

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Theorem E.16 (Reduction of a Riemann-Stieltjes Integral to a finite sum). *Let α be a step function on $[a, b]$ with jump*

$$c_k = \alpha(x_k+) - \alpha(x_k-) \quad \text{at } x = x_k.$$

Let f be defined on $[a, b]$ in such a way that not both of f and α are discontinuous from the left or from the right at x_k . Then $\int_a^b f(x) d\alpha(x)$ exists and

$$\int_a^b f(x) d\alpha(x) = \sum_{k=1}^n f(x_k) c_k.$$

Theorem E.16 (Reduction of a Riemann-Stieltjes Integral to a finite sum). *Let α be a step function on $[a, b]$ with jump*

$$c_k = \alpha(x_k+) - \alpha(x_k-) \quad \text{at } x = x_k.$$

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Theorem E.16 (Reduction of a Riemann-Stieltjes Integral to a finite sum). *Let α be a step function on $[a, b]$ with jump*

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Let f be defined on $[a, b]$ in such a way that not both of f and α are discontinuous from the left or from the right at x_k . Then $\int_a^b f(x) d\alpha(x)$ exists and

$$\int_a^b f(x) d\alpha(x) = \sum_{k=1}^n f(x_k) c_k.$$

Theorem E.12. Given $c \in (a, b)$. Define

$$\alpha(x) = pI_{[a,c)} + rI_{\{c\}} + qI_{(c,b]}$$

(as given in Figure E.1). Suppose at least one of the functions f or α is continuous from the left at c , and at least one is continuous from the right at c . Then $f \in R(\alpha)$ and

$$\int_a^b f(x) d\alpha(x) = f(c)(\alpha(c+) - \alpha(c-)) = f(c)(q - p).^1$$

Remark E.13. The integral $\int_a^b f d\alpha$ does not exist if both of f and α are discontinuous from the left or from the right at c .

Remark E.14. (1) If $\alpha(x) = pI_{\{a\}} + qI_{(a,b]}$, then

$$\int_a^b f(x) d\alpha(x) = f(a)(\alpha(a+) - \alpha(a))$$

Example E.15. (1) Consider

$$f(x) = 1 \quad \text{for } x \in [-1, 1], \quad \text{and} \quad \alpha(x) = -I_{\{0\}},$$

then

$$\int_{-1}^1 f(x) d\alpha(x) = f(0)(\alpha(0+) - \alpha(0-)) = 0$$

(2) Consider

$$f(x) = 2I_{\{0\}} + I_{[-1,0) \cup (0,1]} \quad \text{and} \quad \alpha(x) = -I_{[0,1]}.$$

Then both of α and f are discontinuous from the left at $x = 0$. This implies that the Riemann-Stieltjes integral $\int_{-1}^1 f d\alpha$ does not exist.

Example E.17. (1) Let

$$f(x) = \begin{cases} 3 & \text{if } x \leq 0 \\ 3 - 4x & \text{if } 0 < x < 1 \\ -1 & \text{if } x \geq 1 \end{cases}$$

and

$$\alpha(x) = \begin{cases} 0 & \text{if } x \leq 0 \\ 2 & \text{if } 0 < x < 1 \\ 0 & \text{if } x \geq 1 \end{cases}$$

Since f is continuous, $\int_{-3}^3 f(x) d\alpha(x)$ exists and

$$\begin{aligned} \int_{-3}^3 f(x) d\alpha(x) &= f(0)(\alpha(0+) - \alpha(0-)) + f(1)(\alpha(1+) - \alpha(1-)) \\ &= 3(2 - 0) + (-1)(0 - 2) = 8. \end{aligned}$$

(2) Let $\alpha(x) = 2I_{[0,2)} + 5I_{[2,3)} + 6I_{[3,\infty)}$

$$\begin{aligned} \int_{-5}^{10} e^{-3x} d\alpha(x) &= e^{-3 \cdot 0}(2 - 0) + e^{-3 \cdot 2}(5 - 2) + e^{-3 \cdot 3}(6 - 5) \\ &= 2 + 3e^{-6} + e^{-9}. \end{aligned}$$

Example E.18. Suppose F is the Cantor function (see Figure C.1). By integration by parts, we have

$$\int_0^1 x dF(x) = xF(x)|_{x=0}^1 - \int_0^1 F(x) dx = 1 - \int_0^1 F(x) dx.$$

Since $\int_0^1 F(x) dx$ is the area of the Cantor function on $[0, 1]$, we get

$$\int_0^1 F(x) dx = \frac{1}{2}.$$

Hence,

$$\int_0^1 x dF(x) = \frac{1}{2}.$$

E.3.4. Comparison theorem.

Theorem E.19. Assume that α is an increasing function on $[a, b]$. If $f, g \in R(\alpha)$ on $[a, b]$, and if $f(x) \leq g(x)$ for $x \in [a, b]$, then

$$\int_a^b f(x) d\alpha(x) \leq \int_a^b g(x) d\alpha(x).$$

Corollary E.20. If $g(x) \geq 0$ and α is an increasing function on $[a, b]$, then

$$\int_a^b f(x) d\alpha(x) \geq 0.$$

Theorem E.21. Assume that α is an increasing function on $[a, b]$. If $f \in R(\alpha)$ on $[a, b]$, then

(1) $|f| \in R(\alpha)$ on $[a, b]$, and

$$\left| \int_a^b f(x) d\alpha(x) \right| \leq \int_a^b |f(x)| d\alpha(x).$$

(2) $f^2 \in R(\alpha)$ on $[a, b]$.

Definition E.23. A function $\alpha : [a, b] \rightarrow \mathbb{R}$ is said to be of bounded variation if there exists a constant M such that

$$\sum_{k=1}^n |\alpha(x_k) - \alpha(x_{k-1})| \leq M$$

for every partition $\{x_0, x_1, \dots, x_n\}$ of $[a, b]$.

Theorem E.24. Let α be defined on $[a, b]$, then α is of bounded variation on $[a, b]$, if and only if there exist two increasing functions α_1 and α_2 , such that $\alpha = \alpha_1 - \alpha_2$

Theorem E.25. If f is continuous on $[a, b]$, and if α is of bounded variation on $[a, b]$, then $f \in R(\alpha)$. Moreover, the function

$$F(t) = \int_0^t f(x) d\alpha(x)$$

has the following properties :

- (1) F is of bounded variation on $[a, b]$.
- (2) Every continuous point of α is also a continuous point of F .

POSSIBLE QUESTIONS**UNIT-I****PART-B (5 × 6 = 30)**

1. Show that Reimann Steiljes integral can be reduced to a finite sum.
2. Prove: If $f \in R(\alpha)$ and $f \in R(\beta)$, then $f \in R(c_1 \alpha + c_2 \beta)$.
3. If $f \in R(\alpha)$ on $[a, b]$ then $\alpha \in R(f)$ on $[a, b]$ we have
$$\int_a^b f(x) d\alpha(x) + \int_a^b \alpha(x) df(x) = f(b)\alpha(b) - f(a)\alpha(a).$$
4. Explain change of variables in Riemann – Stieltjes integral ?
5. State and prove linearity property in Reimann Steiltjes integrals.
6. Let $f(x) = \alpha(x) = x$ on $[0, 1]$. Then show that $f \in R(\alpha)$ and find the integral.

PART- C (1 × 10 =10)

1. Prove that the modulus of Reimann Steiltjes integrable functions are also Reimann Steiltjes integrable.
2. State and Prove relation between Riemann integral and Reimann Steiltjes integral.
3. Show that every bounded function on a closed interval is Reimann integrable.
4. Prove that a function f is continuous iff its bounded variation function is also continuous.
5. State and prove necessary and sufficient condition for Reiamann Steiltjes integrability



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DEPARTMENT OF MATHEMATICS
Multiple Choice Questions (Each Question Carries One Mark)
Subject Name: REAL ANALYSIS Subject Code: 17MMP102

UNIT-I	Option 1	option 2	Option 3	Option 4	ANSWERS
$x(y+z) = xy + xz$ is ----- law	commutative	associative	distributive	closure	distributive
If $x < y$, then for every z we have -----	$(x+z) < (y+z)$	$(x+z) > (y+z)$	$(x+z) = (y+z)$	$x+z = 0$	$(x+z) < (y+z)$
If $x > 0$ and $y > 0$, then -----	xy less than equal to 0	$xy > 0$	xy greater than equal to 0	$xy < 0$	$xy > 0$
If $x > y$ and $y > z$, then -----	$x < y$	$x = z$	$x > z$	$x < z$	$x > z$
If a less than equal to $b + \delta$ for every $\delta > 0$, then -----	$a < b$	$a > b$	a greater than equal to b	a less than equal to b	a less than equal to b
The set of all points between a and b is called -----	integer	interval	elements	set	interval
The set $\{x: a < x < b\}$ is -----	(a, b)	$[a, b]$	$[a, b)$	$[a, b)$	(a, b)
A real number is called a positive integer if it belongs to -----	interval	open interval	closed interval	inductive set	inductive set
If d is a divisor of n , then -----	$n = c$	$n < cd$	$n > cd$	$n = cd$	$n = cd$
If $a bc$ and $(a, b) = 1$, then -----	$a c$	$a b$	$b a$	$c a$	$a c$
If $a bc$ and $(a, b) = 1$, then $a c$ is -----	Unique factorisation theorem	additive property	approximation property	Euclid's lemma	Euclid's lemma
Rational numbers is of the form -----	p/q	$p + q$	p/q	$p - q$	p/q
e is -----	rational	irrational	prime	composite	irrational
An integer n is called ----- if the only possible divisors of n are 1 and n	rational	irrational	prime	composite	prime
If $d a$ and $d b$, then d is called -----	LCM	common divisor	prime	function	common divisor
If $(a, b) = 1$, then a and b are called -----	twin prime	common factor	LCM	relatively prime	relatively prime
If an upper bound 'b' of a set S is also a member of S then 'b' is called -----	rational	irrational	maximum element	minimum element	maximum element
If an lower bound 'b' of a set S is also a member of S then 'b' is called -----	rational	irrational	maximum element	minimum element	minimum element
A set with no upper bound is called -----	bounded above	bounded below	prime	function	bounded above
A set with no lower bound is called -----	bounded above	bounded below	prime	function	bounded below
The least upper bound is called -----	bounded above	bounded below	supremum	infimum	supremum
The greatest lower bound is called -----	bounded above	bounded below	supremum	infimum	infimum
The supremum of $\{3, 4\}$ is -----	3	4	$\{3, 4\}$	$\{3, 4\}$	4
Every finite set of numbers is -----	bounded	unbounded	prime	bounded above	bounded
A set S of real numbers which is bounded above and bounded below is called -----	bounded set	inductive set	super set	subset	bounded set
The set N of natural numbers is -----	bounded	not bounded	irrational	rational	not bounded
The completeness axiom is -----	$b = \sup S$	$S = \sup b$	$b = \inf S$	$S = \inf b$	$b = \sup S$
The infimum of $\{3, 4\}$ is -----	3	4	$\{3, 4\}$	$\{3, 4\}$	3
$\sup C = \sup A + \sup B$ is called ----- property	approximation	additive	archimedean	comparison	additive
For any real x , there is a positive integer n such that -----	$n < x$	$n < x$	$n = x$	$n = 0$	$n > x$
If $x > 0$ and if y is an arbitrary real number, there is a positive number n such that $nx > y$ is ----- property	approximation	additive	archimedean	comparison	archimedean
The set of positive integers is -----	bounded above	bounded below	unbounded above	unbounded below	unbounded above
The absolute value of x is denoted by -----	$ x $	$ x $	$x < 0$	$x > 0$	$ x $
If $x < 0$ then -----	$ x = x$	$ x = x $	$ x = -x$	$ x = -x$	$ x = -x$
If $S = [0, 1)$ then $\sup S =$ -----	0	1	$(0, 1)$	$[0, 1]$	1
Triangle inequality is -----	$ a + b $ greater equal to $ a + b $	$ a > a + b $	$ b > a + b $	$ a + b $ less than equal to $ a + b $	$ a + b $ less than equal to $ a + b $
$ x + y $ greater than equal to -----	$ x + y $	$ x y $	$ x - y $	$ x - y $	$ x - y $
Set of real numbers S is bounded above implies S has a -----	supremum	infimum	additive property	comparison property	supremum
In $\{(3n+2)/(2n+1) \text{ such that } n \text{ is in } \mathbb{N}\}$, the greatest lower bound is -----	5 divided by 3	8 divided by 5	11 divided by 47	3 divided by 2	3 divided by 2
In Cauchy-Schwarz inequality, the equality holds iff -----	$ax = 0$	$ax + bx = 0$	$ax + bx = 0$	$bx = 0$	$ax + bx = 0$
If a set consists of a finite number of elements is called -----	infinite set	finite set	cantor set	null set	finite set
A sequence $\langle S_n \rangle$ is said to be monotonically increasing if -----	$S_n \geq S_{n+1}$	$S_n \leq S_{n+1}$	$S_n \geq S_{n+1}$	$S_n \leq S_{n+1}$	$S_n \geq S_{n+1}$
If A, B, C are three sets then what is $A - (B - C) =$ -----	$(A - B) \cup (A \cap C)$	$A - (B \cap C)$	$(A - B) \cup C$	$(A - B) \cup (A - C)$	$(A - B) \cup (A \cap C)$
If S is a non-empty set of real numbers and S is unbounded below, then -----	$\inf S = -\infty$	$\inf S = \infty$	$\sup S = -\infty$	$\sup S = \infty$	$\inf S = -\infty$
If $P(A)$ denotes the power set of A and A is the void set then $P(P(P(A))) =$ -----	0	1	4	16	16
If $X \subseteq \mathbb{R}$ then -----	$X/\infty = \infty$	$X/\infty = 0$	$X/\infty = X$	$X/\infty = -\infty$	$X/\infty = 0$
If $x < 0$ then -----	$X(-\infty) = -\infty$	$X(-\infty) = \infty$	$X(-\infty) = 0$	$X(-\infty) = X$	$X(-\infty) = -\infty$
If \mathbb{R}^* is an extended real number system then the least upper bound is -----	∞	negative infinity	0	no least upper bound	∞
Let $f: \mathbb{R} \rightarrow \mathbb{R}$ be a function defined as $f(x) = x x $ then -----	f is 1-1 but not onto	neither f is 1-1 nor onto	f is 1-1 both onto	f is onto but not one-one	f is onto but not one-one
The value of $(0, \infty)$ is -----	∞	0	not defined	can not be determined	0



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UNIT II

Infinite series and infinite products:

Introduction – Basic definitions – Ratio test and root test – Dirichlet test and Able's test
Rearrangement of series – Riemann's theorem on conditionally convergent series – Sub series -
Double sequences – Double series – Multiplication of series – Cesaro summability.

TEXT BOOK

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Definition :

Given an infinite sequence $\{a_n\}$, the n th partial sum S_n is the sum of the first n terms of the sequence, that is,

A series is **convergent** if the sequence of its partial sums S_n tends to a limit; that means that the partial sums become closer and closer to a given number when the number of their terms increases. More precisely, a series converges, if there exists a number L such that for any arbitrarily small positive number ϵ , there is a (sufficiently large) integer N such that for all $n > N$,

If the series is convergent, the number L (necessarily unique) is called the **sum of the series**.

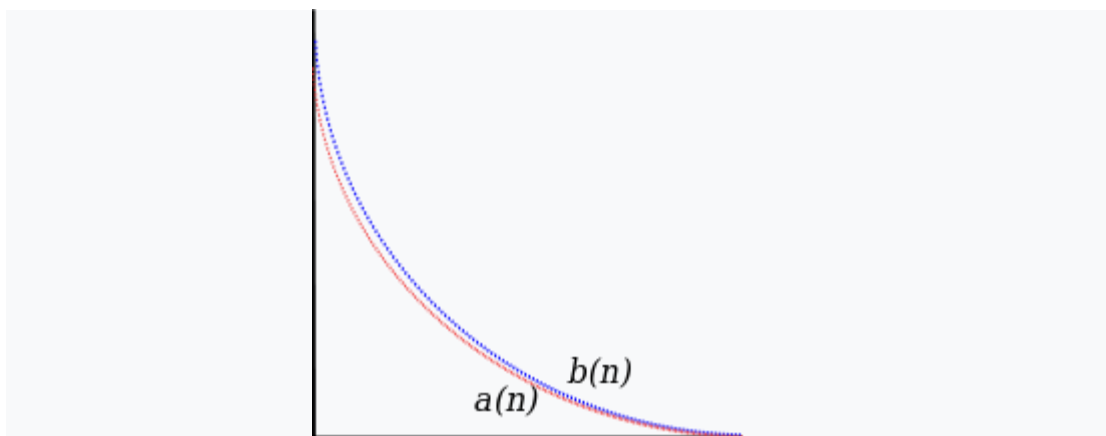
Any series that is not convergent is said to be divergent.

- The reciprocals of the positive integers produce a divergent series (harmonic series):
- Alternating the signs of the reciprocals of positive integers produces a convergent series:
- The reciprocals of prime numbers produce a divergent series (so the set of primes is "large"):
- The reciprocals of triangular numbers produce a convergent series:
- The reciprocals of factorials produce a convergent series (see e):
- The reciprocals of square numbers produce a convergent series (the Basel problem):
- The reciprocals of powers of 2 produce a convergent series (so the set of powers of 2 is "small"):

The reciprocals of powers of any n produce a convergent series:
 Alternating the signs of reciprocals of powers of 2 also produces a convergent series:

- Alternating the signs of reciprocals of powers of any n produces a convergent series:
- The reciprocals of Fibonacci numbers produce a convergent series (see ψ):

There are a number of methods of determining whether a series converges or diverges.



If the blue series, $\sum b(n)$, can be proven to converge, then the smaller series, $\sum a(n)$, must converge. By contraposition, if the red series, $\sum a(n)$, is proven to diverge, then $\sum b(n)$ must also diverge.

Comparison test. The terms of the sequence $a(n)$ are compared to those of another

sequence $b(n)$. If,

for all n , $a(n) \leq b(n)$, and $\sum b(n)$ converges, then so does

However, if,

for all n , $a(n) \leq b(n)$, and $\sum b(n)$ diverges, then so does

Ratio test. Assume that for all n , $a(n) > 0$. Suppose that there exists L such that

If $r < 1$, then the series converges. If $r > 1$, then the series diverges. If $r = 1$, the ratio test is inconclusive, and the series may converge or diverge.

Root test or n th root test. Suppose that the terms of the sequence in question are non-negative. Define r as follows:

where "lim sup" denotes the limit superior (possibly ∞ ; if the limit exists it is the same value).

If $r < 1$, then the series converges. If $r > 1$, then the series diverges. If $r = 1$, the root test is inconclusive, and the series may converge or diverge.

The ratio test and the root test are both based on comparison with a geometric series, and as such they work in similar situations. In fact, if the ratio test works (meaning that the limit exists and is not equal to 1) then so does the root test; the converse, however, is not true. The root test is therefore more generally applicable, but as a practical matter the limit is often difficult to compute for commonly seen types of series.

Integral test. The series can be compared to an integral to establish

convergence or divergence. Let $f(x)$ be a positive and monotone decreasing function. If

$\int_1^\infty f(x) dx$ converges, then the series converges. But if the integral diverges, then the series does so as well.

Limit comparison test. If $\sum a_n$ and $\sum b_n$ are series of positive terms, and the limit $\lim_{n \rightarrow \infty} \frac{a_n}{b_n}$ exists and is not zero, then $\sum a_n$ converges if and only if $\sum b_n$ converges.

Alternating series test. Also known as the *Leibniz criterion*, the alternating series test states that for an alternating series of the

form $\sum (-1)^n a_n$, if a_n is monotone decreasing, and has a limit of 0 at infinity, then the series converges.

Cauchy condensation test. If $\{a_n\}$ is a positive monotone decreasing sequence, then $\sum a_n$ converges if and only if $\sum 2^n a_{2^n}$ converges.

Dirichlet's test

Abel's test

Raabe's test

Conditional and absolute convergence[edit]

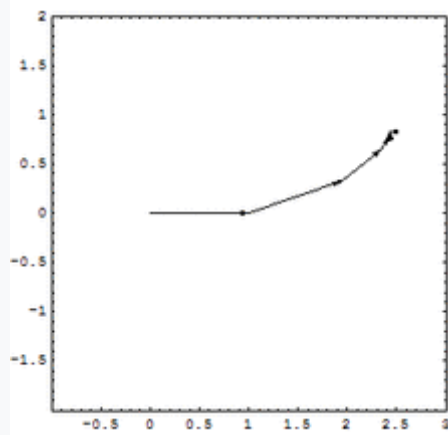


Illustration of the absolute convergence of the power series of $\text{Exp}[z]$ around 0 evaluated at $z = \text{Exp}[1/3]$. The length of the line is finite.

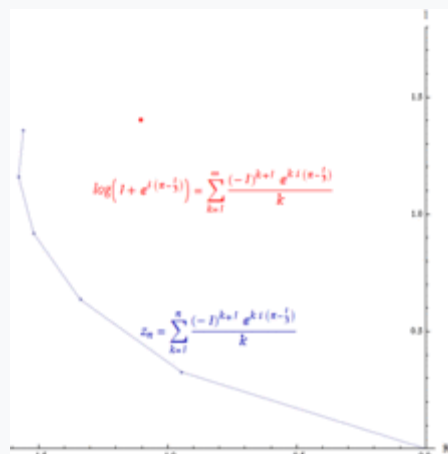


Illustration of the conditional convergence of the power series of $\log(z+1)$ around 0 evaluated at $z = \exp((\pi-1/3)i)$. The length of the line is infinite.

For any sequence $\{a_n\}$, $\sum_{n=1}^{\infty} a_n$ for all n . Therefore,

This means that if $\sum_{n=1}^{\infty} a_n$ converges, then $\sum_{n=1}^{\infty} |a_n|$ also converges (but not vice versa).

If the series $\sum_{n=1}^{\infty} a_n$ converges, then the series $\sum_{n=1}^{\infty} |a_n|$ is absolutely convergent. An absolutely convergent sequence is one in which the length of the line created by joining together all of the increments to the partial sum is finitely long. The power series of the exponential function is absolutely convergent everywhere.

If the series $\sum a_n$ converges but the series $\sum b_n$ diverges, then the

series $\sum (a_n + b_n)$ is conditionally convergent. The path formed by connecting the partial sums of a conditionally convergent series is infinitely long. The power series of the logarithm is conditionally convergent.

The Riemann series theorem states that if a series converges conditionally, it is possible to rearrange the terms of the series in such a way that the series converges to any value, or even diverges.

Uniform convergence[edit]

Main article: uniform convergence

Let $\{f_n\}$ be a sequence of functions. The series $\sum f_n$ is said to

converge uniformly to f if the sequence $\{S_n\}$ of partial sums defined by

$S_n(x) = \sum_{k=0}^n f_k(x)$ converges uniformly to f .

There is an analogue of the comparison test for infinite series of functions called the Weierstrass M-test.

Cauchy convergence criterion[edit]

The **Cauchy convergence criterion** states that a series

$\sum a_n$ converges if and only if the sequence of partial sums is

a Cauchy sequence. This means that for every $\epsilon > 0$ there is a

positive integer N such that for all $m, n > N$ we have

1. SEQUENCES

1.1. Sequences. An infinite sequence of real numbers is an ordered unending list of real numbers. E.g.:

$$1, 2, 3, 4, \dots$$

We represent a generic sequence as a_1, a_2, a_3, \dots , and its n -th as a_n .

In order to define a sequence we must give enough information to find its n -th term. Two ways of doing this are:

1. With a formula. E.g.:

1

$$a_n = \frac{1}{n}$$

$$a_n = \frac{1}{10^n}$$

$$a_n = \sqrt{3n - 7}$$

1.2. Limit of a Sequence. We say that a sequence a_n *converges* to a limit L if the difference $|a_n - L|$ can be made as small as we wish by taking n large enough. We write $a_n \rightarrow L$, or more formally:

$$\lim_{n \rightarrow \infty} a_n = L.$$

E.g.:

$$\lim_{n \rightarrow \infty} \frac{1}{n} = 0.$$

If a sequence does not converge we say that it *diverges*. E.g., the following sequences diverge:

$$n = 1, 2, 3, 4, \dots \rightarrow \quad \text{diverges (to } +\infty)$$

$$(-1)^n = -1, 1, -1, 1, \dots \rightarrow \quad \text{diverges}$$

1.3. Limit Laws for Sequences. If $\lim_{n \rightarrow \infty} a_n = A$ and $\lim_{n \rightarrow \infty} b_n = B$, then:

$$\lim_{n \rightarrow \infty} (a_n + b_n) = A + B$$

$$\lim_{n \rightarrow \infty} (a_n - b_n) = A - B$$

$$\lim_{n \rightarrow \infty} (a_n b_n) = AB$$

$$\lim_{n \rightarrow \infty} (a_n/b_n) = A/B \quad (\text{provided } B \neq 0)$$

So, a “complicated” limit such as $L = \lim_{n \rightarrow \infty} \frac{1 + \frac{1}{n}}{3 + \frac{1}{10^n}}$ can be computed by replacing smaller parts of it with their limits $1/n \rightarrow 0$, $1/10^n \rightarrow 0$:

$$L = \frac{1 + 0}{3 + 0} = \frac{1}{3}.$$

1.4. Squeeze Law. If $a_n \leq c_n \leq b_n$, and a_n and b_n have the same limit: $a_n \rightarrow L$, $b_n \rightarrow L$, then c_n has also the same limit: $c_n \rightarrow L$. This can be used to compute limits such as the following one:

$$\lim_{n \rightarrow \infty} \frac{\sin n}{n}.$$

In this case we have:

$$-\frac{1}{n} \leq \frac{\sin n}{n} \leq \frac{1}{n}.$$

Since $-1/n \rightarrow 0$ and $1/n \rightarrow 0$ then $\frac{\sin n}{n} \rightarrow 0$ also.

1.5. Limits of Functions of Sequences. If $a_n = f(n)$ for some function f and $\lim_{x \rightarrow \infty} f(x) = L$, then $\lim_{n \rightarrow \infty} a_n = L$. This basically allows us to replace limits of sequences with limits of functions. In particular this is useful for using L'Hôpital's rule in computing limits of sequences. E.g:

$$\lim_{n \rightarrow \infty} \frac{e^n}{n} = \lim_{x \rightarrow \infty} \frac{e^x}{x} = (\text{L'Hôpital's rule}) = \lim_{x \rightarrow \infty} \frac{e^x}{1} = \infty.$$

2.1. Series. An infinite series is an expression of the form

$$\sum_{n=1}^{\infty} a_n = a_1 + a_2 + a_3 + \cdots ,$$

¹Proof: We use induction. First note that $0 < a_1 = \sqrt{6} < 3$. By adding 6 and taking square roots we get $\sqrt{6} < \sqrt{6 + a_1} < \sqrt{6 + 3} = 3$, i.e.: $a_1 < a_2 < 3$. Now assume $a_n < a_{n+1} < 3$ for a given $n \geq 1$ (induction hypothesis). Again, by adding 6 and taking square roots we get $\sqrt{6 + a_n} < \sqrt{6 + a_{n+1}} < 3$, i.e. $a_{n+1} < a_{n+2} < 3$ (induction step). From here we get that $a_n < a_{n+1} < 3$ for every $n \geq 1$, which proves both, a_n is increasing and is bounded by 3.

²C.H. Edwards, Jr. & David E. Penney: *Calculus with Analytic Geometry*, 5th edition, Prentice Hall.

where $\{a_n\}$ is a sequence of numbers—sometimes the series starts at $n = 0$ or some other term instead of $n = 1$. Its N th partial sum is

$$S_N = \sum_{n=1}^N a_n = a_1 + a_2 + a_3 + \cdots + a_N.$$

2.2. Sum of a Series. The sum

$$S = \sum_{n=1}^{\infty} a_n$$

of a series is defined as the limit of its partial sums

$$S = \lim_{N \rightarrow \infty} S_N = \lim_{N \rightarrow \infty} \sum_{n=1}^N a_n$$

if it exists—it this case we say that the series converges. For instance, consider the following series:

$$\sum_{n=1}^{\infty} \frac{1}{2^n} = \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \cdots$$

Its partial sum is

$$S_N = \sum_{n=1}^N \frac{1}{2^n} = \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \cdots + \frac{1}{2^N} = 1 - \frac{1}{2^N}.$$

Hence, its sum is

$$S = \lim_{N \rightarrow \infty} S_N = \lim_{N \rightarrow \infty} \left(1 - \frac{1}{2^N}\right) = 1,$$

i.e.:

$$\sum_{n=1}^{\infty} \frac{1}{2^n} = 1.$$

A series may or may not have a sum. For instance, in the following series:

$$\sum_{n=0}^{\infty} (-1)^n = 1 - 1 + 1 - 1 + 1 - 1 + \cdots$$

the sequence of partial sums $S_N = 1, 0, 1, 0, 1, 0, \dots$ diverges, and the series has no sum.

2.3. Telescopic Series. Telescopic series are series for which all terms of its partial sum can be canceled except the first and last ones. For instance, consider the following series:

$$\sum_{n=1}^{\infty} \frac{1}{n(n+1)} = \frac{1}{2} + \frac{1}{6} + \frac{1}{12} + \cdots$$

Its n th term can be rewritten in the following way:

$$a_n = \frac{1}{n(n+1)} = \frac{1}{n} - \frac{1}{n+1}.$$

Hence, its N th partial sum becomes:

$$\begin{aligned} S_N &= \sum_{n=1}^N \frac{1}{n(n+1)} = \sum_{n=1}^N \left(\frac{1}{n} - \frac{1}{n+1} \right) \\ &= \left(1 - \frac{1}{2} \right) + \left(\frac{1}{2} - \frac{1}{3} \right) + \left(\frac{1}{3} - \frac{1}{4} \right) + \cdots + \left(\frac{1}{N} - \frac{1}{N+1} \right) \\ &= 1 - \frac{1}{N+1}. \end{aligned}$$

Hence:

$$\sum_{n=1}^{\infty} \frac{1}{n(n+1)} = \lim_{N \rightarrow \infty} \left(1 - \frac{1}{N+1} \right) = 1.$$

2.4. Geometric series. A geometric series $\sum_{n=0}^{\infty} a_n$ is a series in which each term is a fixed multiple of the previous one: $a_{n+1} = r a_n$, where r is called the *ratio*. A geometric series can be rewritten in this way:

$$\sum_{n=0}^{\infty} a r^n = a + a r + a r^2 + a r^3 + \cdots.$$

If $|r| < 1$ its sum is

$$\sum_{n=0}^{\infty} a r^n = \frac{a}{1-r}.$$

Note that a is the first term of the series. If $a \neq 0$ and $|r| \geq 1$, the series diverges.

Examples:

$$\sum_{n=0}^{\infty} \frac{1}{2^n} = 1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \cdots = \frac{1}{1 - \frac{1}{2}} = 2.$$

$$\sum_{n=0}^{\infty} \frac{(-1)^n}{2^n} = 1 - \frac{1}{2} + \frac{1}{4} - \frac{1}{8} + \cdots = \frac{1}{1 - (-\frac{1}{2})} = \frac{2}{3}.$$

Note that in the last example $r = a_{n+1}/a_n = \frac{(-1)^{n+1}/2^{n+1}}{(-1)^n/2^n} = -1/2$.

6.1. Alternating Series Test. If an alternating series verifies:

1. a_n it is decreasing: $a_n \geq a_{n+1} > 0$ for every n , and
2. the n th term tends to zero: $\lim_{n \rightarrow \infty} a_n = 0$,

then the series converges.

So, in this particular case the “reciprocal” of the n th term test holds.

E.g.:

$$\sum_{n=1}^{\infty} \frac{(-1)^{n+1}}{n} = 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \cdots = \ln 2$$

$$\sum_{n=0}^{\infty} \frac{(-1)^{n+1}}{2n+1} = 1 - \frac{1}{3} + \frac{1}{5} - \frac{1}{7} + \cdots = \frac{\pi}{4}.$$

6.3. Absolute Convergence. A general series $\sum a_n$ is said to be *absolutely convergent* if the series of absolute values of its terms $\sum |a_n|$ is convergent.

We have that a series can be:

1. Convergent and absolutely convergent, e.g:

$$\sum_{n=0}^{\infty} \left(-\frac{1}{2}\right)^n = 1 - \frac{1}{2} + \frac{1}{4} - \frac{1}{8} + \cdots$$

2. Convergent but not absolutely convergent—in this case the series is called **conditionally convergent**—, e.g:

$$\sum_{n=1}^{\infty} \frac{(-1)^{n+1}}{n} = 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \cdots$$

3. Not convergent nor absolutely convergent, e.g:

$$\sum_{n=1}^{\infty} n = 1 + 2 + 3 + 4 + \cdots$$

However, a series cannot be absolutely convergent and not convergent, because absolute convergence implies convergence:

$$\text{absolute convergent} \implies \text{convergent}$$

Example: Does the series $\sum_{n=1}^{\infty} \frac{\cos n}{n^2}$ converge? Answer: Look at the series of absolute values:

$$\sum_{n=1}^{\infty} \frac{|\cos n|}{n^2} \leq \sum_{n=1}^{\infty} \frac{1}{n^2}.$$

By comparison test, it converges (the right hand side is a p -series with $p > 1$), hence the given series is absolutely convergent, which implies that it is indeed convergent.

6.4. Ratio Test. Suppose that the limit $\rho = \lim_{n \rightarrow \infty} \left| \frac{a_{n+1}}{a_n} \right|$ exists or is infinity. Then

1. If $\rho < 1 \implies \sum a_n$ converges absolutely.
2. If $\rho > 1 \implies \sum a_n$ diverges.
3. If $\rho = 1 \implies$ the ratio test is inconclusive.

As a rule of thumb, for geometric series $\rho = |r|$ (the ratio), and the conclusion of the ratio test is analogous to the one for geometric series, i.e., the series converges for $|r| < 1$ and diverges for $|r| > 1$.

Example: For the series $\sum_{n=1}^{\infty} \frac{n^2}{2^n}$ we have

$$\rho = \lim_{n \rightarrow \infty} \left| \frac{a_{n+1}}{a_n} \right| = \lim_{n \rightarrow \infty} \frac{(n+1)^2/2^{n+1}}{n^2/2^n} = \lim_{n \rightarrow \infty} \frac{(n+1)^2}{2n^2} = \frac{1}{2} < 1,$$

hence, it converges absolutely.³

6.5. Root Test. In some cases in which the ratio test is unable to provide an answer, the root test may help. It says the following: Suppose that the limit $\rho = \lim_{n \rightarrow \infty} \sqrt[n]{|a_n|}$ exists or is infinity. Then

1. If $\rho < 1 \implies \sum a_n$ converges absolutely.
2. If $\rho > 1 \implies \sum a_n$ diverges.
3. If $\rho = 1 \implies$ the root test is inconclusive.

Example: Consider the following series $\sum_{n=1}^{\infty} \frac{1}{2^{n+\sin n}}$. For this series the ratio test cannot be used, because

$$\frac{a_{n+1}}{a_n} = 2^{-1+\sin n - \sin(n+1)} = 2^{-1-2\sin \frac{1}{2} \cos(n+\frac{1}{2})}$$

which has no limit. However, the root test shows that the series is absolutely convergent:

$$\lim_{n \rightarrow \infty} \sqrt[n]{\frac{1}{2^{n+\sin n}}} = \lim_{n \rightarrow \infty} \frac{1}{2^{1+\sin n/n}} = \frac{1}{2} < 1.$$

POSSIBLE QUESTIONS**UNIT-II****PART-B ($5 \times 6 = 30$)**

1. Prove: If $\sum a_n$ converges absolutely then the subseries $\sum b_n$ is also converges

absolutely. Also $|\sum_{n=1}^{\infty} b_n| \leq \sum_{n=1}^{\infty} |a_n|$.

2. Suppose $\sum_{n=1}^{\infty} a_n$ converges absolutely and $\sum_{n=1}^{\infty} a_n = A$ and $\sum_{n=1}^{\infty} b_n = B$; then show that the product of these two series converges to AB .

3. State and prove Ratio test for convergence of series.

4. State and Prove Riemann theorem on conditionally convergent series .

5. Using Root test, show that the series

$$\frac{1}{2} + \frac{1}{3} + \frac{1}{2^2} + \frac{1}{3^2} + \frac{1}{2^3} + \frac{1}{3^3} + \cdots \dots \dots \dots \dots \dots \dots \text{converges}$$

PART- C ($1 \times 10 = 10$)

1. State and prove uniform convergence in double sequences

2. Let $\sum a_n$ be an absolutely convergent series having sum S then every rearrangement of $\sum a_n$ also converges absolutely has sum S .



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DEPARTMENT OF MATHEMATICS

Multiple Choice Questions (Each Question Carries One Mark)

Subject Name: REAL ANALYSIS

Subject Code: 17MMP102

UNIT II	OPTION1	OPTION2	OPTION3	OPTION4	ANSWERS
The coordinates (x,y) of a point represent an ----- of numbers	function	relation	ordered pair	set	ordered pair
$(a, b) = \{ \}$ -----	$\{(a), (b), (a, b)\}$	$\{(a), (b)\}$	$\{(a), (a, b)\}$	$\{(a), (b), \{\}\}$	$\{(a), (a, b)\}$
$(a, b) = (c, d)$ if and only if -----	$a = c$ & $b = d$	$a = b$ & $c = d$	$a = d$ & $c = b$	$ab = cd$	$a = c$ & $b = d$
$A \times B$ denotes the ----- of the sets A & B	product	cartesian product	polar form	complement	cartesian product
Any set of ordered pairs is called -----	function	relation	ordered pair	set	relation
If S is a relation, the set of all elements that occur as first members in S is called the -----	function	codomain	domain	range	domain
If S is a relation, the set of all elements that occur as second members in S is called the -----	function	codomain	domain	range	range
If (x, y) belongs to F and (x, z) belongs to F, then -----	$x = z$	$x = y$	$xy = z$	$y = z$	$y = z$
A mapping S into itself is called -----	function	relation	domain	transformation	transformation
If $F(x) = F(y)$ implies $x = y$ is a ----- function	one-one	onto	into	inverse	one-one
One-one function is also called -----	injective	bijective	transformation	codomain	injective
$S = \{(a, b) : (b, a) \text{ is in } S\}$ is called -----	inverse	domain	codomain	converse	converse
The composite functions are denoted by -----	$G \circ F$	$G \circ F$	$G \circ F$	$G + F$	$G \circ F$
$G \circ F(x) = \{ \}$ -----	$G[F(x)]$	$F[G(x)]$	$G(x)$	$F(x)$	$G[F(x)]$
In general the composite function $G \circ F$ is -----	$G \circ F = F \circ G$	$G \circ F$ is not equal to $F \circ G$	$G \circ F < F \circ G$	$G \circ F > F \circ G$	$G \circ F$ is not equal to $F \circ G$
If $m < n$, then $K(m) < K(n)$ implies that K is -----	sequence	subsequence	order preserving	equinumerous	order preserving
Similar sets are also called as ----- set	denumerable	uncountable	finite	equinumerous	equinumerous
If A and B are two sets and if there exists a one-one correspondence between them, then it is called ----- set	denumerable	uncountable	finite	equinumerous	equinumerous
A set which is equinumerous with the set of all positive integers is called ----- set	finite	infinite	countably infinite	countably finite	countably infinite
A set which is either finite or countably infinite is called ----- set	countable	uncountable	similar	equal	countable
Uncountable sets are also called ----- set	denumerable	non-denumerable	similar	equal	non-denumerable
Countable sets are also called ----- set	denumerable	non-denumerable	similar	equal	denumerable
Every subset of a countable set is -----	countable	uncountable	rational	irrational	countable
The set of all real numbers is -----	countable	uncountable	rational	irrational	uncountable
The cartesian product of the set of all positive integers is -----	countable	uncountable	rational	irrational	countable
The set of those elements which belong either to A or to B or to both is called -----	complement	intersection	union	disjoint	union
The set of those elements which belong to both A and B is called -----	complement	intersection	union	disjoint	intersection
Union of sets is -----	commutative	not commutative	not associative	disjoint	commutative
The complement of A relative to B is denoted by -----	$B - A$	B	A	$A - B$	$B - A$
If A intersection B is the empty set, then A and B are called -----	commutative	not commutative	not associative	disjoint	disjoint
$B - (\text{union } A) = \{ \}$ -----	union (B - A)	$B - (\text{intersection } A)$	intersection (B - A)	$\{ \}$	intersection (B - A)
$B - (\text{intersection } A) = \{ \}$ -----	union (B - A)	$B - (\text{union } A)$	intersection (B - A)	$\{ \}$	union (B - A)
Union of countable sets is -----	uncountable	infinite	countable	disjoint	countable
The set of all rational numbers is -----	uncountable	infinite	countable	disjoint	countable
The set of intervals with rational end points is ----- set	uncountable	infinite	countable	disjoint	countable
A relation which is symmetric, reflexive and transitive is called ----- relation	equivalence	component	composite	countable	equivalence
Any collection of disjoint intervals of positive length is -----	equivalence relation	countable set	composite function	uncountable set	countable set
If A similar to B and B similar to C, then -----	C similar to A	A similar to C	$A < C$	$A = C$	A similar to C
If the root of an algebraic equation $f(x) = 0$, then the real number is called -----	prime	positive	algebraic	composite	algebraic
For all subsets A and B of S with B contained in A, we have -----	$f(A + B) = f(A)$	$f(A + B) = f(B)$	$f(A - B) = f(A) - f(B)$	$f(A - B) = f(A)$	$f(A - B) = f(A) - f(B)$
If $f(A \cup B) = f(A) + f(B)$, then the function f is called -----	additive	multiplicative	disjoint	equinumerous	additive
$f(A \cup B) = \{ \}$ -----	$f(A) + f(B)$	$f(A) - f(B)$	$f(A) + f(B) - f(B - A)$	$f(A) + f(B - A)$	$f(A) + f(B - A)$
The sequence $\langle (-1)^n \rangle$ is -----	monotonically increasing	monotonically decreasing	either increasing or decreasing	neither monotonically increasing nor monotonically decreasing	neither monotonically increasing nor monotonically decreasing
An unbounded sequence -----	a limit point	does not have a limit point	may or may not have a limit point	unique limit point	may or may not have a limit point
The series $\sum_{n=1}^{\infty} (-1)^{n-1} \frac{1}{n}$ -----	convergent	conditionally convergent	absolutely divergent	need not convergent	conditionally convergent
Every absolutely convergent series is -----	convergent	conditionally convergent	absolutely divergent	need not convergent	convergent
The sequence $\langle 1/n \rangle$ is -----	convergent & bounded	divergent & unbounded	divergent & bounded	convergent & Unbounded	convergent & bounded
If a sequence $\{a_n\}$ converges to a real number then the given sequence is -----	unbounded sequence	convergent	divergent & bounded	bounded	unbounded sequence
Every subsequence has a -----	limit point	convergent	monotonic subsequence	non monotonic sequence	monotonic subsequence
The series $1 + r + r^2 + r^3 + \dots$ is oscillatory if -----	$r = 1$	$r = -1$	$r > 1$	$r < 1$	$r = -1$



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DEPARTMENT OF MATHEMATICS

Subject : REAL ANALYSIS
SUBJECT CODE: 17MMP102

SEMESTER: I
CLASS : I M.Sc

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UNIT III

Sequences of functions:

Basic definitions – Uniform convergence and continuity - Uniform convergence of infinite series of functions – Uniform convergence and Riemann – Stieltjes integration – Non uniformly convergent sequence – Uniform convergence and differentiation – Sufficient condition for uniform convergence of a series

TEXT BOOK

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Uniform Convergence.

1 Introduction.

In this course we study amongst other things **Fourier series**. The **Fourier series** for a periodic function $f(x)$ with period 2π is defined as the series

$$\frac{a_0}{2} + \sum_{k=1}^{\infty} (a_k \cos kx + b_k \sin kx),$$

where the coefficients a_k, b_k are defined as

$$a_k = \frac{1}{\pi} \int_{-\pi}^{\pi} f(x) \cos kx \, dx, \quad b_k = \frac{1}{\pi} \int_{-\pi}^{\pi} f(x) \sin kx \, dx,$$

with $k = 0, 1, \dots$ (note that this means that $b_0 = 0$).

This is an example of a **functional series**, which is a series whose terms are functions:

$$\sum_{k=0}^{\infty} u_k(x).$$

As usual with series, we define the above infinite sum as a limit:

$$\sum_{k=0}^{\infty} u_k(x) = \lim_{N \rightarrow \infty} \sum_{k=0}^N u_k(x),$$

providing the limit exists. Note that different values of x will, in general, give different limits, if they exist.

In this lecture we shall look at functional series, and functional sequences, and we shall consider first the question of convergence. To deal with this, we consider two types of convergence: **pointwise convergence** and **uniform convergence**. There are three main results: the first one is that **uniform convergence** of a sequence of continuous functions gives us a continuous function as a limit. The second main result is **Weierstrass' Majorant Theorem**, which gives a condition that guarantees that a functional series converges to a continuous function. The third result is that integrals of a sequence of functions which converges uniformly to a limit function $f(x)$ also converge with the limit being the integral of $f(x)$. These results are not only good for your mental health, they are also important tools in our later discussion of Fourier series, and that is the reason for looking at them.

Definition 2.1 (Pointwise convergence.) Suppose $\{f_n(x) : n = 0, 1, 2, \dots\}$ is a sequence of functions defined on an interval I . We say that $f_n(x)$ **converges pointwise to the function $f(x)$ on the interval I** if

$$f_n(x) \rightarrow f(x), \quad \text{as } n \rightarrow \infty, \quad \text{for each } x \in I.$$

We call the function $f(x)$ **the limit function**.

Uniform continuity

Definition. A function $f: D \rightarrow \mathbb{R}$ is *continuous at a point $d \in D$* if for any $\epsilon > 0$ there exists $\delta > 0$ such that for all $x \in D$,

$$|x - d| < \delta \Rightarrow |f(x) - f(d)| < \epsilon.$$

The following is an equivalent ‘sequential’ definition of the same concept.

Definition. A function $f: D \rightarrow \mathbb{R}$ is *continuous at a point $d \in D$* if for any sequence $(x_n)_{n \in \mathbb{N}}$ in D with $d = \lim_{n \rightarrow \infty} x_n$, the convergence $f(d) = \lim_{n \rightarrow \infty} f(x_n)$ holds as well.

Definition. A function $f: D \rightarrow \mathbb{R}$ is *continuous* if it is continuous at every point of D .

Definition. A function $f: D \rightarrow \mathbb{R}$ is *uniformly continuous* if for all $\epsilon > 0$ there exists $\delta > 0$ such that for all $x, y \in D$,

$$|x - y| < \delta \Rightarrow |f(x) - f(y)| < \epsilon.$$

The crucial point to note here is that for uniform continuity, the number δ may only depend on ϵ . For continuity (at x), it may depend on x as well. Using quantifiers, we express continuity of f as follows:

$$\forall x \in D \forall \epsilon > 0 \exists \delta > 0 \forall y \in D: |x - y| < \delta \Rightarrow |f(x) - f(y)| < \epsilon.$$

Uniform continuity, on the other hand, means that

$$\forall \epsilon > 0 \exists \delta > 0 \forall x, y \in D: |x - y| < \delta \Rightarrow |f(x) - f(y)| < \epsilon.$$

It is clear that uniform continuity implies continuity. But continuity does not imply uniform continuity.

Example. Define $f: \mathbb{R} \rightarrow \mathbb{R}, x \mapsto \sin(x^2)$. Then f is continuous but not uniformly continuous (see Exercise 1.1).

Example 2.1 $f_n(x) = x - \frac{1}{n}$. Then $f_n(x)$ converges pointwise to x for each $x \in \mathbb{R}$:

$$|f_n(x) - x| = \frac{1}{n} \rightarrow 0 \quad \text{as } n \rightarrow \infty.$$

Example 2.2 $f_n(x) = e^{-nx}$ on $[1, 3]$. For each $x \in [1, 3]$ we have $nx \rightarrow \infty$ as $n \rightarrow \infty$ and therefore $f_n(x) \rightarrow 0$ as $n \rightarrow \infty$ for each $x \in [1, 3]$. Thus $f_n(x)$ converges pointwise to $f(x) = 0$ for each $x \in [1, 3]$.

Example 2.3 $f_n(x) = e^{-nx}$ on $[0, 3]$. For each $0 < x \leq 3$ we have $nx \rightarrow \infty$ as $n \rightarrow \infty$ and therefore $f_n(x) \rightarrow 0$ as $n \rightarrow \infty$ for each $0 < x \leq 3$. However, at $x = 0$ we have $f_n(0) = 1$ for all n . Thus $f_n(x)$ converges pointwise to the function $f(x)$ defined by $f(0) = 1$, $f(x) = 0$ for each $0 < x \leq 3$. This is not a continuous function, despite the fact that each function $f_n(x)$ is continuous.

Example 2.4 Let the sequence f_n be defined as

$$f_n(x) = \frac{nx}{(nx+1)^3} \quad x \in [0, \infty[.$$

Then $f_n(0) = 0$ and for each fixed $x > 0$

$$\begin{aligned} f_n(x) &= \frac{n^2x}{(nx+1)^3} \\ &= \frac{n^2x}{n^3(x+\frac{1}{n})^3} \\ &= \frac{1}{n} \frac{x}{(x+\frac{1}{n})^3} \rightarrow 0 \quad \text{as } n \rightarrow \infty. \end{aligned}$$

So that $f_n(x) \rightarrow f(x) = 0$ pointwise on $[0, \infty[$.

Then for $x > -\frac{1}{n}$ we have

$$f'_n(x) = \frac{n^2(1-2nx)}{(nx+1)^4}$$

and we see that for $x > 0$ we have $f'_n(x) \rightarrow 0$ as $n \rightarrow \infty$ whereas $f'_n(0) = n^2 \rightarrow \infty$. Here we see that $f'_n \rightarrow f'$ only on for $x > 0$. This shows that differentiability is not always respected by pointwise convergence.

The last two examples then lead us to pose the question: what extra condition (other than just pointwise convergence) can guarantee that the limit function is also continuous or differentiable? The answer to this is given by the concept of **uniform convergence**.

3 Uniform convergence

We define for a real-valued (or complex-valued) function f on a non-empty set I the **supremum norm** of f on the set I :

$$\|f\|_I = \sup_{x \in I} |f(x)|.$$

Note that if f is a bounded function on I then

$$\sup_{x \in I} |f(x)| = \sup\{|f(x)| : x \in I\}$$

exists, by the so-called **supremum axiom**. Observe that

$$|f(x)| \leq \|f\|_I \quad \text{for all } x \in I,$$

and that $|f(x)|$ takes on values which are arbitrarily near $\|f\|_I$. In particular $\|f\|_I =$ the largest value of $|f(x)|$ whenever such a value exists (such as when I is a closed, bounded interval and $f(x)$ is a continuous function on I).

Definition. Consider a sequence $(f_n)_{n \in \mathbb{N}}$ of functions $f_n: D \rightarrow \mathbb{R}$ and another function $f: D \rightarrow \mathbb{R}$.

- (i) We say that f_n converges to f *pointwise* if for all $x \in D$ and all $\epsilon > 0$, there exists $N \in \mathbb{N}$ such that for all $n \geq N$,

$$|f_n(x) - f(x)| < \epsilon.$$

- (ii) We say that f_n converges to f *uniformly* if for all $\epsilon > 0$ there exists $N \in \mathbb{N}$ such that for all $x \in D$ and all $n \geq N$,

$$|f_n(x) - f(x)| < \epsilon.$$

The above definition is equivalent to the following.

Definition. Consider a sequence $(f_n)_{n \in \mathbb{N}}$ of functions $f_n: D \rightarrow \mathbb{R}$ and another function $f: D \rightarrow \mathbb{R}$.

- (i) We say that f_n converges to f *pointwise* if

$$f(x) = \lim_{n \rightarrow \infty} f_n(x)$$

for every fixed $x \in D$.

- (ii) We say that f_n converges to f *uniformly* if

$$\sup_{x \in D} |f_n(x) - f(x)| \rightarrow 0$$

as $n \rightarrow \infty$.

Clearly uniform convergence implies pointwise convergence, but pointwise convergence does not imply uniform convergence.

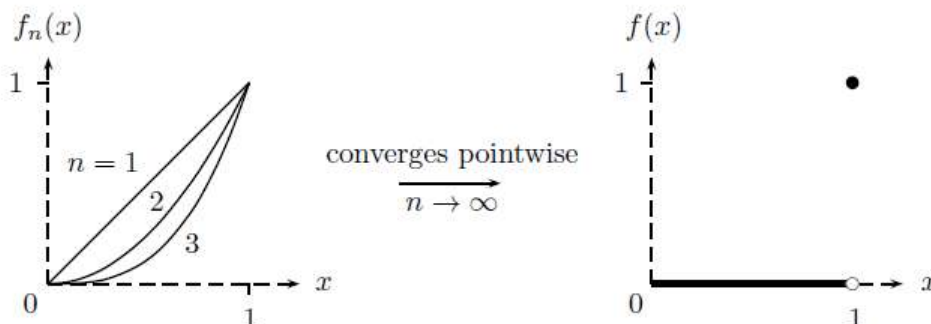
Example. Let $D = [0, 1]$. Consider the sequence $(f_n)_{n \in \mathbb{N}}$ of functions $f_n: D \rightarrow \mathbb{R}$, $x \mapsto x^n$. The pointwise limit is the discontinuous function $f: D \rightarrow \mathbb{R}$ given by

$$f(x) = \begin{cases} 0 & \text{if } x \in [0, 1), \\ 1 & \text{if } x = 1. \end{cases}$$

The convergence is not uniform, however, since

$$\sup_{x \in D} |f_n(x) - f(x)| = 1$$

for all $x \in \mathbb{N}$.



5.4. Properties of uniform convergence

In this section we prove that, unlike pointwise convergence, uniform convergence preserves boundedness and continuity. Uniform convergence does not preserve differentiability any better than pointwise convergence. Nevertheless, we give a result that allows us to differentiate a convergent sequence; the key assumption is that the derivatives converge uniformly.

5.4.1. Boundedness. First, we consider the uniform convergence of bounded functions.

Theorem 5.14. Suppose that $f_n : A \rightarrow \mathbb{R}$ is bounded on A for every $n \in \mathbb{N}$ and $f_n \rightarrow f$ uniformly on A . Then $f : A \rightarrow \mathbb{R}$ is bounded on A .

Proof. Taking $\epsilon = 1$ in the definition of the uniform convergence, we find that there exists $N \in \mathbb{N}$ such that

$$|f_n(x) - f(x)| < 1 \quad \text{for all } x \in A \text{ if } n > N.$$

Choose some $n > N$. Then, since f_n is bounded, there is a constant $M_n \geq 0$ such that

$$|f_n(x)| \leq M_n \quad \text{for all } x \in A.$$

It follows that

$$|f(x)| \leq |f(x) - f_n(x)| + |f_n(x)| < 1 + M_n \quad \text{for all } x \in A,$$

meaning that f is bounded on A (by $1 + M_n$). \square

We do not assume here that all the functions in the sequence are bounded by the same constant. (If they were, the pointwise limit would also be bounded by that constant.) In particular, it follows that if a sequence of bounded functions converges pointwise to an unbounded function, then the convergence is not uniform.

Example 5.15. The sequence of functions $f_n : (0, 1) \rightarrow \mathbb{R}$ in Example 5.2, defined by

$$f_n(x) = \frac{n}{nx + 1},$$

cannot converge uniformly on $(0, 1)$, since each f_n is bounded on $(0, 1)$, but their pointwise limit $f(x) = 1/x$ is not. The sequence (f_n) does, however, converge uniformly to f on every interval $[a, 1)$ with $0 < a < 1$. To prove this, we estimate for $a \leq x < 1$ that

$$|f_n(x) - f(x)| = \left| \frac{n}{nx + 1} - \frac{1}{x} \right| = \frac{1}{x(nx + 1)} < \frac{1}{nx^2} \leq \frac{1}{na^2}.$$

Thus, given $\epsilon > 0$ choose $N = 1/(a^2\epsilon)$, and then

$$|f_n(x) - f(x)| < \epsilon \quad \text{for all } x \in [a, 1) \text{ if } n > N,$$

5.4.2. Continuity. One of the most important property of uniform convergence is that it preserves continuity. We use an “ $\epsilon/3$ ” argument to get the continuity of the uniform limit f from the continuity of the f_n .

Theorem 5.16. If a sequence (f_n) of continuous functions $f_n : A \rightarrow \mathbb{R}$ converges uniformly on $A \subset \mathbb{R}$ to $f : A \rightarrow \mathbb{R}$, then f is continuous on A .

Proof. Suppose that $c \in A$ and $\epsilon > 0$ is given. Then, for every $n \in \mathbb{N}$,

$$|f(x) - f(c)| \leq |f(x) - f_n(x)| + |f_n(x) - f_n(c)| + |f_n(c) - f(c)|.$$

By the uniform convergence of (f_n) , we can choose $n \in \mathbb{N}$ such that

$$|f_n(x) - f(x)| < \frac{\epsilon}{3} \quad \text{for all } x \in A,$$

and for such an n it follows that

$$|f(x) - f(c)| < |f_n(x) - f_n(c)| + \frac{2\epsilon}{3}.$$

(Here we use the fact that f_n is close to f at both x and c , where x is an arbitrary point in a neighborhood of c ; this is where we use the uniform convergence in a crucial way.)

Since f_n is continuous on A , there exists $\delta > 0$ such that

$$|f_n(x) - f_n(c)| < \frac{\epsilon}{3} \quad \text{if } |x - c| < \delta \text{ and } x \in A,$$

which implies that

$$|f(x) - f(c)| < \epsilon \quad \text{if } |x - c| < \delta \text{ and } x \in A.$$

This proves that f is continuous. □

This result can be interpreted as justifying an “exchange in the order of limits”

$$\lim_{n \rightarrow \infty} \lim_{x \rightarrow c} f_n(x) = \lim_{x \rightarrow c} \lim_{n \rightarrow \infty} f_n(x).$$

Such exchanges of limits always require some sort of condition for their validity — in this case, the uniform convergence of f_n to f is sufficient, but pointwise convergence is not.

It follows from Theorem 5.16 that if a sequence of continuous functions converges pointwise to a discontinuous function, as in Example 5.3, then the convergence is not uniform. The converse is not true, however, and the pointwise limit of a sequence of continuous functions may be continuous even if the convergence is not uniform, as in Example 5.4.

Definition 3.1 A sequence of functions $f_n(x)$ defined on an set I is said to **converge uniformly to $f(x)$ on I** if

$$\|f_n - f\|_I \rightarrow 0 \text{ as } n \rightarrow \infty.$$

We write this as

$$\lim_{n \rightarrow \infty} f_n = f \text{ uniformly on } I$$

or as

$$f_n \rightarrow f \text{ uniformly on } I \text{ as } n \rightarrow \infty.$$

Uniform convergence implies **pointwise convergence**, however there are sequences which converge pointwise but not uniformly. Indeed we have

$$|f_n(x) - f(x)| \leq \sup_{x \in I} |f_n(x) - f(x)| = \|f_n - f\|_I,$$

so that

$$f_n \rightarrow f \text{ uniformly on } I \text{ as } n \rightarrow \infty$$

$$\implies |f_n(x) - f(x)| \rightarrow 0 \text{ for each } x \in I$$

$$\implies f_n \rightarrow f \text{ pointwise on } I.$$

Example 3.1 $f_n(x) = e^{-nx}$ on $[1, 3]$. We have seen above that $f_n(x)$ converges pointwise to $f(x) = 0$ for each $x \in [1, 3]$. Then we have $|f_n(x) - f(x)| = |f_n(x)|$ and we then have

$$\begin{aligned} \|f_n - f\| &= \sup_{x \in [1, 3]} |f_n(x)| \\ &= \sup_{x \in [1, 3]} |e^{-nx}| \\ &= \sup_{x \in [1, 3]} e^{-nx} \\ &= e^{-n} \rightarrow 0 \text{ as } n \rightarrow \infty. \end{aligned}$$

Thus we have uniform convergence in this case. Note that the last step follows from the observation that e^{-nx} is strictly decreasing for $x \geq 0$ with $n \geq 0$, so that $e^{-n} \geq e^{-nx}$ for all $x \geq 1$.

Theorem 4.3 Suppose that $\{f_n(x); n = 0, 1, 2, \dots\}$ is a sequence of functions on an interval I and satisfying the following conditions:

- (i) $f_n(x)$ is **differentiable** on I for each $n = 0, 1, 2, \dots$
- (ii) $f_n(x)$ converges **pointwise** to $f(x)$ on I
- (iii) $f'_n(x)$ is continuous for each n and $f'_n \rightarrow g$ converges **uniformly** on I where $g(x)$ is a continuous function on I .

Then the limit function $f(x)$ is differentiable and $f'(x) = g(x)$.

Proof: First note that

$$f_n(x) - f_n(a) = \int_a^x f'_n(t) dt$$

for each $f_n(x)$ and for each choice of $x, a \in I$. Because $f_n(x)$ converges pointwise to $f(x)$ for all $x \in I$, the left-hand side converges to $f(x) - f(a)$ as $n \rightarrow \infty$. Also, $f'_n \rightarrow g$ uniformly on I so by Theorem 4.2 we have that

$$\int_a^x f'_n(t) dt \rightarrow \int_a^x g(t) dt,$$

and we then find that

$$f(x) - f(a) = \int_a^x g(t) dt.$$

Now, $g(t)$ is continuous, so that, by the Fundamental Theorem of Calculus,

$$\frac{d}{dx} \int_a^x g(t) dt = g(x)$$

so that $f(x)$ must be differentiable and $f'(x) = g(x)$.

This result is very useful, as we shall see, in examining the differentiability of functional series.

5 Applications to functional series.

Definition 5.1 A functional series is a series

$$\sum_{k=0}^{\infty} u_k(x)$$

where each term of the series $u_k(x)$ is a function on an interval I .

We can also define **pointwise convergence** for functional series:

Definition 5.2 *The functional series*

$$\sum_{k=0}^{\infty} u_k(x)$$

is **pointwise convergent** for each $x \in I$ if the limit

$$\sum_{k=0}^{\infty} u_k(x) = \lim_{N \rightarrow \infty} \sum_{k=0}^N u_k(x)$$

exists for each $x \in I$.

Thus, we always define a **sequence of partial sums** $S_N(x)$ given as

$$S_N(x) = \sum_{k=0}^N u_k(x)$$

so that

$$S_0(x) = u_0(x), \quad S_1(x) = u_0(x) + u_1(x), \quad S_2(x) = u_0(x) + u_1(x) + u_2(x)$$

and if

$$\lim_{N \rightarrow \infty} S_N(x)$$

exists for x then we say that the series

$$\sum_{k=0}^{\infty} u_k(x) = \lim_{N \rightarrow \infty} S_N(x)$$

converges at x . It converges pointwise on the interval I if

$$\lim_{N \rightarrow \infty} S_N(x)$$

exists for each $x \in I$.

Theorem 5.1 *Suppose that the functional series*

$$\sum_{k=0}^{\infty} u_k(x)$$

is defined on an interval I and that there is a sequence of positive constants M_k so that

$$|u_k(x)| \leq M_k, \quad k = 0, 1, 2, \dots$$

for all $x \in I$. If

$$\sum_{k=0}^{\infty} M_k$$

converges, then

$$\sum_{k=0}^{\infty} u_k(x)$$

converges uniformly on I .

Proof: If the conditions are fulfilled then we immediately have, from the **Comparison Theorems for Positive Series**, that, for each $x \in I$, the series

$$\sum_{k=0}^{\infty} |u_k(x)|$$

is convergent, so that

$$\sum_{k=0}^{\infty} u_k(x)$$

is absolutely convergent, and therefore convergent. This means that

$$\sum_{k=0}^{\infty} u_k(x)$$

is **pointwise convergent** on I , and we denote the limit by $S(x)$. We now show that the partial sums

$$S_N(x) = \sum_{k=0}^N u_k(x)$$

converges uniformly to $S(x)$ on I under the conditions of the theorem. We have

$$S(x) - S_N(x) = \sum_{k=N+1}^{\infty} u_k(x)$$

(all we do is subtract the first N terms from the series). Then it follows that

$$|S(x) - S_N(x)| \leq \sum_{k=N+1}^{\infty} |u_k(x)| \leq \sum_{k=N+1}^{\infty} M_k$$

for each $x \in I$, since $|u_k(x)| \leq M_k$ for each $x \in I$ according to our assumption. Then

$$\|S - S_N\|_I \leq \sum_{k=N+1}^{\infty} M_k.$$

We also know (by assumption) that $\sum_{k=0}^{\infty} M_k$ converges, so we must have that $\sum_{k=N+1}^{\infty} M_k \rightarrow 0$ as $N \rightarrow \infty$. Consequently,

$$\|S - S_N\|_I \rightarrow 0 \quad \text{as } N \rightarrow \infty,$$

and our result is proved.

Corollary 5.1 *If*

(i) *the functional series*

$$S(x) = \sum_{k=0}^{\infty} u_k(x) \quad \text{converges uniformly on interval } I,$$

(ii) *$u_k(x)$ is a continuous function on I for each $k = 0, 1, 2, \dots$,*

then $S(x)$ is continuous on I .

Proof: Because a finite sum of continuous functions is again a continuous function, it follows that the partial sums

$$S_N(x) = \sum_{k=0}^N u_k(x)$$

are continuous functions for $N = 0, 1, 2, \dots$. Then by Theorem 4.1, we have that $S(x) = \lim_{N \rightarrow \infty} S_N(x)$ is a continuous function.

Example 5.1 *Take the functional series*

$$\sum_{k=1}^{\infty} \frac{\sin kx}{k^2}.$$

We have

$$|u_k(x)| = \left| \frac{\sin kx}{k^2} \right| = \frac{|\sin kx|}{k^2} \leq \frac{1}{k^2}$$

since $|\sin t| \leq 1$ for all real t . We know (standard positive series) that

$$\sum_{k=1}^{\infty} \frac{1}{k^2}$$

converges (series of the form $\sum 1/k^\alpha$ converge for $\alpha > 1$ and diverge for $\alpha \leq 1$). Hence, by Weierstrass' Majorant Theorem,

$$\sum_{k=1}^{\infty} \frac{\sin kx}{k^2}$$

converges uniformly for all x , and by Corollary 5.1 this series is a continuous function of x for all $x \in \mathbb{R}$.

Theorem 5.2 *If*

(i) *the functional series*

$$\sum_{k=0}^{\infty} u_k(x) \quad \text{converges uniformly on the interval } I$$

(ii) $u_k(x)$ *is continuous on* I *for each* $k = 0, 1, 2, \dots$,

then

$$\int_a^x \left(\sum_{k=0}^{\infty} u_k(t) \right) dt = \sum_{k=0}^{\infty} \left(\int_a^x u_k(t) dt \right)$$

for all $a, x \in I$. In other words, if the series of continuous functions converges uniformly on I , then the integral of the sum is the sum of the integrals of the functions, just as in the case of a finite sum.

Proof: Put $S(t) = \sum_{k=0}^{\infty} u_k(t)$ and $S_N(t) = \sum_{k=0}^N u_k(t)$, then we have $S_N \rightarrow S$ uniformly on I so that

$$\lim_{N \rightarrow \infty} \int_a^x S_N(t) dt = \int_a^x \lim_{N \rightarrow \infty} S_N(t) dt = \int_a^x S(t) dt,$$

according to Theorem 4.2. Note that since $S_N(t)$ is a finite sum of functions, we see that

$$\begin{aligned}\int_a^x S_N(t)dt &= \int_a^x \left(\sum_{k=0}^N u_k(t) \right) dt \\ &= \sum_{k=0}^N \left(\int_a^x u_k(t)dt \right),\end{aligned}$$

and we then find that

$$\int_a^x \left(\sum_{k=0}^{\infty} u_k(t) \right) dt = \sum_{k=0}^{\infty} \left(\int_a^x u_k(t)dt \right).$$

We can also say something about the differentiability of the series $\sum u_k(x)$, using Theorem 4.3 In this case, as in the previous two theorems, we replace $f_n(x)$ by $S_N(t)$ and $f(x)$ by $S(t)$. Thus, we want the following:

- $S_N(x) \rightarrow S(x)$ pointwise on I
- $S'_N(x) \rightarrow G(x)$ uniformly on I
- $S_N(x)$ is continuously differentiable for each N

and then we may conclude that $S(x)$ is continuously differentiable with $S'(x) = G(x)$. All we need is to formulate these requirements and result as follows:

Theorem 5.3 Suppose that $\sum_{k=0}^{\infty} u_k(x)$ satisfies the following conditions:

- $\sum_{k=0}^{\infty} u_k(x)$ converges pointwise on I
- $\sum_{k=0}^{\infty} u'_k(x)$ converges uniformly on I
- $u_k(x)$ is continuously differentiable for each k

Then $\sum_{k=0}^{\infty} u_k(x)$ is continuously differentiable and

$$\frac{d}{dx} \left(\sum_{k=0}^{\infty} u_k(x) \right) = \sum_{k=0}^{\infty} u'_k(x).$$

Proof: We have that each $S'_N(x)$ is continuous on I and

$$S_N(x) \longrightarrow S(x) \quad \text{pointwise on } I$$

$$S'_N \longrightarrow G \quad \text{uniformly on } I.$$

Then by Theorem 4.3, $S(x)$ is differentiable and $S'(x) = G(x)$ on I . In other words:

$$\frac{d}{dx} \left(\sum_{k=0}^{\infty} u_k(x) \right) = \sum_{k=0}^{\infty} u'_k(x)$$

Definition 6.1 *We say that the integral*

$$F(x) = \int_a^{\infty} f(x, y) dy$$

converges uniformly on I if:

(i) $F(x) = \int_a^\infty f(x, y)dy$ converges pointwise for each $x \in I$;

(ii) the family of functions F_R defined as

$$F_R(x) = \int_a^R f(x, y)dy$$

converges uniformly to F on I . That is, if

$$\|F_R - F\|_I \longrightarrow 0 \quad \text{as } R \rightarrow \infty.$$

A test for uniform convergence of integrals is an analogy of the Weierstrass functional series:

Theorem 6.1 (M-test) Suppose

(i) $f(x, y)$ is continuous on $I \times [a, \infty[$

(ii) $|f(x, y)| \leq M(y)$ for all $x \in I$ and $y \in [a, \infty[$

(iii) $\int_a^\infty M(y)dy$ converges.

Then

$$F(x) = \int_a^\infty f(x, y)dy$$

converges uniformly on I .

Proof: We have

$$F_R(x) - F(x) = \int_R^\infty f(x, y)dy$$

from which we obtain

$$\begin{aligned} |F_R(x) - F(x)| &\leq \int_R^\infty |f(x, y)|dy \\ &\leq \int_R^\infty M(y)dy, \end{aligned}$$

by assumption. Consequently,

$$\|F_R - F\|_I \leq \int_R^\infty M(y)dy \rightarrow 0 \quad \text{as } R \rightarrow \infty,$$

because

$$\int_a^\infty M(y)dy \text{ converges}$$

implies that

$$\int_R^\infty M(y)dy \longrightarrow 0 \text{ as } R \rightarrow \infty.$$

Thus, $F_R \rightarrow F$ uniformly on I .

Now we come to proving that if $F_R \rightarrow F$ uniformly on an interval I , then F is continuous if each F_R is continuous. here, the problem is to show that $F_R(x)$ is continuous. We have the following result:

Lemma 7.1 Suppose that $f(x)$ is a real-valued continuous function on the closed, bounded interval $[a, b]$. Then

$$\sup_{x \in [a, b]} f(x) = \max\{f(x) : x \in [a, b]\}, \quad \inf_{x \in [a, b]} f(x) = \min\{f(x) : x \in [a, b]\}.$$

That is, the supremum of a continuous function over a closed, bounded interval is equal to its largest value over that interval, and the infimum is the least value of the function over the interval.

Proof: Since $f(x)$ is continuous and the interval is closed, then $f(x)$ has a largest value and a least value on the interval: there exist $x_1, x_2 \in [a, b]$ so that $f(x_1) \leq f(x) \leq f(x_2)$ for all $x \in [a, b]$, and we now see that

$$\sup_{x \in [a, b]} f(x) = f(x_2), \quad \inf_{x \in [a, b]} f(x) = f(x_1),$$

and the result is proved.

We sketch the proof and refer to any good book on analysis for further details.

Our task is to prove that for each $x \in [c, d]$ we have

$$F_R(x+h) \longrightarrow F_R(x) \quad \text{as } h \rightarrow 0.$$

Then we have

$$|F_R(x+h) - F_R(x)| \leq \int_a^R |f(x+h, y) - f(x, y)| dy.$$

Now if f is continuous on $[c, d] \times [a, R]$ it can be shown that for any $\epsilon > 0$ there is a $\delta > 0$ so that

$$|f(x_0, y_0) - f(x_1, y_1)| < \epsilon$$

whenever $\sqrt{(x_0 - x_1)^2 + (y_0 - y_1)^2} < \delta$. That is, whenever the distance between the points (x_0, y_0) and (x_1, y_1) is less than δ . This is called **uniform continuity**. Using this fact, we choose $\epsilon > 0$ (arbitrarily small) and then for each given R we find a $\delta > 0$ so that

$$|f(x+h, y) - f(x, y)| < \frac{\epsilon}{R-a}$$

whenever $|h| < \delta$. From this it follows that

$$|F_R(x+h) - F_R(x)| < \epsilon \quad \text{whenever } |h| < \delta.$$

Note that we have $\epsilon > 0$ arbitrarily small, and for each such choice there is a corresponding δ . From this it follows that

$$|F_R(x+h) - F_R(x)| \longrightarrow 0 \quad \text{as } h \rightarrow 0.$$

Hence $F_R(x)$ is continuous on $[c, d]$, for each choice of $R > 0$.

7.3 Cauchy's condition for uniform convergence of series

In this section we record, without proof, a result which is of some interest: **Cauchy's criterion for uniform convergence of functional series**, and we make some comments on some aspects of uniform convergence.

Theorem 7.1 Suppose that $\{u_k(x)\}$ is a sequence of continuous functions defined on an interval I . Then the series

$$\sum_{k=0}^{\infty} u_k(x)$$

is uniformly convergent on I if and only if for each choice of $\epsilon > 0$, however small, there exists a (corresponding) integer $N > 0$ so that

$$|u_{k+1}(x) + u_{k+2}(x) + \cdots + u_m(x)| < \epsilon$$

for all $m > k \geq N$ and for all $x \in I$. In particular, we then have, on putting $m = k + 1$,

$$|u_{k+1}(x)| < \epsilon$$

for all $k \geq N$ and all $x \in I$.

The proof of this result requires more mathematical machinery than we have at hand, and can be found in any good textbook on Mathematical Analysis.

As an illustration of the usefulness of this result we look at the series expansion of e^x . We have

$$e^x = 1 + x + \frac{x^2}{2!} + \cdots + \frac{x^n}{n!} + \cdots = \sum_{k=0}^{\infty} \frac{x^k}{k!}.$$

This expansion is true for each $x \in \mathbb{R}$, so we have pointwise convergence of the series. However, we do not have uniform convergence on \mathbb{R} . To see this, we apply the last comment in the theorem: we need to find, for a given $\epsilon > 0$, an $N > 0$ so that **for all** $x \in \mathbb{R}$ we have

$$|u_{k+1}(x)| = \frac{|x|^{k+1}}{k!} < \epsilon$$

whenever $k \geq N$. However, if we choose any k we may choose x so that

$$\frac{|x|^{k+1}}{k!}$$

is as large as we like, contradicting the requirement for uniform convergence. Hence we do not have uniform convergence on the whole of \mathbb{R} . However, if we only consider $x \in [-a, a]$ for some $a > 0$ then we can prove uniform convergence of the series on this **closed, bounded interval**. This can be done using Weierstrass' Majorant Theorem. This phenomenon occurs often, and then we say that the series **converges uniformly on closed, bounded intervals** or **converges on compact sets**. Another example of this phenomenon occurs in **power series** (which are the first kind of functional series taught in elementary calculus courses). For instance, for the geometric series

$$\sum_{k=0}^{\infty} x^k$$

we have absolute convergence for $|x| < 1$ and divergence for $|x| \geq 1$. For $|x| < 1$ we have

$$\sum_{k=0}^{\infty} x^k = \frac{1}{1-x} = S(x).$$

The corresponding partial sums are

$$S_N(x) = \sum_{k=0}^N x^k = \frac{1-x^{N+1}}{1-x}.$$

The sequence $S_N(x)$ does not converge uniformly to $S(x)$ on the interval $] - 1, 1[$: we have

$$|S_N(x) - S(x)| = \frac{|x|^{N+1}}{1-x}$$

and as $x \rightarrow 1_-$ we see that $|x|^{N+1} \rightarrow 1$ and $1/(1-x) \rightarrow \infty$, so that we may make $|S_N(x) - S(x)|$ as large as we like, and it then follows that $\|S_N - S\|$ does not exist, so it is impossible for $\|S_N - S\| \rightarrow 0$ as $N \rightarrow \infty$. However, if we consider the series on the closed, bounded interval $[-a, a]$ with a fixed $0 < a < 1$, we have

$$|S_N(x) - S(x)| \leq \frac{a^{N+1}}{1-a}$$

for all $x \in [-a, a]$, and therefore

$$\|S_N - S\| = \sup_{x \in [-a, a]} |S_N(x) - S(x)| \leq \frac{a^{N+1}}{1-a} \rightarrow 0 \text{ as } N \rightarrow \infty,$$

because $0 < a < 1$ gives $a^{N+1} \rightarrow 0$ when $N \rightarrow \infty$. So we have uniform convergence of the series **on compact subsets of** $] - 1, 1[$.

POSSIBLE QUESTIONS**UNIT-III****PART-B (5 × 6 = 30)**

1. Show that the sequence of functions $\{f_n\}$ converges uniformly on E iff $\forall \epsilon > 0, \exists$ an integer N such that for every $m, n \geq N$; $|f_n(x) - f_m(x)| < \epsilon$.
2. Prove that limit of sequences of functions in $R(\alpha)$ which converge uniformly on $[a, b]$ is also in $R(\alpha)$.
3. State and prove Weistrass M-test for Uniform convergence of series of functions.
4. Let series $\sum f_n(x) = f(x)$ converges uniformly on S be such that each f_n is continuous at a point x_0 of S then f is also continuous at x_0 .
5. State and prove Cauchy criterion for convergence of sequences of functions.
6. If $\lim_{n \rightarrow \infty} f_n = f$ and $\lim_{n \rightarrow \infty} g_n = g$ on $[a, b]$; define $h(x) = \int_a^x f(t)g(t)dt$ and $h_n(x) = \int_a^x f_n(t)g_n(t)dt$ for each $x \in [a, b]$ then prove that $h_n \rightarrow h$ uniformly on $[a, b]$.
7. Let the sequence $\{f_n\}$ is a real valued function such that $f_n \in R(\alpha)$ on $[a, b]$ for each $n \in \mathbb{N}$. If $f_n \rightarrow f$ uniformly on $[a, b]$, then $f \in R(\alpha)$ on $[a, b]$ ”.

PART- C(1 × 10 =10)

1. Let $\{f_n\}$ be sequence of differentiable functions on $[a, b]$. If $\{f_n\}$ converges to f uniformly, then show that $f'(x) = \lim_{n \rightarrow \infty} f'_n(x)$.
2. State and prove Arzela Theorem.
3. Let the series $\sum f_n(x)$ be such that $\{f_n\}$ is uniformly convergent on S . Suppose $\{g_n\}$ be a sequence of real valued function such that $g_{(n+1)} \leq g_n(x)$ and uniformly bounded on S . Then show that the series $\sum f_n(x)g_n(x)$ converges uniformly on S .
4. Let α be of bounded variation on $[a, b]$ and each term of the series $\sum f_n$ is Riemann Stieltjes integrable with respect to α on $[a, b]$ and if $\sum_{n=1}^{\infty} f_n(x) = f(x)$ ($a \leq x \leq b$) series converges . uniformly on $[a, b]$, then show that $f \in R(\alpha)$ and

$$\int_a^b f d\alpha = \sum_{n=1}^{\infty} \int_a^b f_n d\alpha$$



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DEPARTMENT OF MATHEMATICS
Multiple Choice Questions (Each Question Carries One Mark)

Subject Name: REAL ANALYSIS Subject Code: 17MMP102

UNIT III	OPTION1	OPTION2	OPTION3	OPTION4	ANSWERS
A Sequence of functions is said to boundedly convergent on T is seq	point wise convergent	it contains no limit points	it contains some limit points	it contains infinite limit points	it contains all of its limit points
A set F is closed if	it contains all of its limit points	uncountable collection of disjoint	countable collection of disjoint	uncountable collection of disjoint	countable collection of disjoint
Every open set of real numbers is the union of	countable collection of disjoint	open intervals	closed intervals	closed intervals	open intervals
Composite number n is	a prime number and $n > 1$	Uncountable	infinite	finite	
The union of a finite or collection of countable sets is	countable	a does not belongs to S	a is not lower bound of S	a is not upper bound of S	
An element a is an minimal element of set S, then	a belongs to S	$n \times x$	$n \times x$	$n \times x$	
For every real number x , there is a positive integer n such that	$n > x$	uncountable subset	proper subset	improper subset	countable subset
Every infinite set has a	countable subset	finite	countable	uncountable	finite
Set of real numbers is bounded above is Sup S	infinite	minimal element only	maximal and minimal	no maximal no minimal	minimal element only
The half interval $(0, 1)$ have	maximal element only	finite	countable	uncountable	infinite
Set of real numbers is unbounded above is Sup S	infinite	closed	semi open	semi closed	closed
The arbitrary intersection of closed set is	open	a singleton set	a finite set	not a well defined set	not a well defined set
The set of intelligent student in a class is	a null set	sum of prime numbers	product of prime numbers	prime numbers or a product of prime numbers	prime numbers or a product of prime numbers
Every integer $n > 1$ is	prime numbers	non ordered set	set of irrational numbers	does not satisfies principle induction	ordered set
The set of integer is	ordered set	unbounded below	unbounded above	no maximal element	bounded above
The closed interval $S = [0, 1]$ is	bounded above		1	2 empty	1
If $S = [0, 1]$ the least upper bound for S is	a point of closure to S	0 prime number	not a point of closure	non prime number	a point of closure to S
If S is a set of real numbers which is bounded below then inf S is	closed set	closed set	uncountable set	countable set	closed set
A finite set is	open set	inf $E > \sup E$	inf $E = \sup E$	inf $E \neq \sup E$	
If E is a nonempty set then	inf $E < \sup E$	∞	∞ (negative)	no infimum	∞ (negative)
If R is a extended real number system then inf R is	0	0	1	0	2
The set of negative integers having least upper bound is	-1	closure of E contains non empty opensets	closure of E contains empty opensets	closure of E contains non empty closedsets	closure of E contains non empty opensets
The set E is nowhere dense if	closure of E contains no non empty opensets	upper bound	maximal element	minimal element	lower bound
The set of natural numbers has	lower bound		0	1 no maximal element	a
Let $S = [0, 1]$ the maximal element of S is	a	its complement is closed set	its complement is null set	its complement is semiclosed set	its complement is closed set
If A is a non-empty open set then	closed set	closed set	empty set	unbounded set	open set
the intersection finite collection of open set is	open set	finite set	unbounded set	unbounded set	open set
The set of real number R is	open set	bounded from below	bounded from above	bounded	unbounded
The set of real numbers is	unbounded	closed set	empty set	non empty set	closed set
The intersection of any collection of closed set is	open set	limit point	infinite limit point	finite limit point	limit point
An infinite set must possess a	does not have a limit point	open intervals	open	closed intervals	open
The empty set is	imperfect	open set	uncountable set	countable set	closed
single ton set $\{x\}$ is	closed	closed	$\{0\}$	a	open set
The union of any or collection of open sets is	a	open	semi open	either open or closed	closed
The derived set of a set is	closed	limit point	largest limit point	no limit point	smallest limit point
Every bounded infinite set has	smallest limit point	countable	finite	infinite	countable
The set of all integers is	uncountable	countable	finite	infinite	countable
The cartesian product of two countable set is	uncountable	E^* is null	E^* is open	E^* is closed interval	E^* is closed
Let E^* is the set of point of closure of E	E^* is closed	semi open	closed intervals	open intervals	open set
Null set	open set	bounded above by 0 & minimal element is 0	bounded below by 1 & no maximal element	bounded above by 1 & maximal element is 1	bounded above by 1 & maximal element is 1
$S = [0, 1]$ is	bounded above by 1 & maximal element is 1	A - B is non empty set	A - B is closed set may not always be a closed set	A - B is empty set	A - B is open set
If A is open set and B is closed set then	A - B is open set	may be closed set	set	open set	may not always be a closed set
The union of an arbitrary family of closed set	closed set	equal sequence	range set of a sequence	null sequence	range set of a sequence
The set of all distinct element of a sequence is called	constant sequence	bounded below	bounded	neither bounded above nor bounded below	neither bounded above nor bounded below
$< (-1)^n n >$	bounded above	oscillates finitely	diverges	converges or oscillates finitely	converges or oscillates finitely
A bounded sequence	converges	one limit	many limit	no limit point	more than one limit
A sequence can not converge to	more than one limit	need not be a member of the sequence	need not be a member of the sequence	not a member of the sequence	need not be a member of the sequence
limit point of a sequence	member of the sequence	no limit point	a limit point	more than two limit point	a limit point
Every bounded real sequence has	many limit point				



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DEPARTMENT OF MATHEMATICS

Subject : QUANTITATIVE METHODS FOR MANAGEMENT
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Unit IV

Transportation problem - Mathematical formulation of Transportation problem - Initial Basic Feasible solution - Optimum solution for non degeneracy and degeneracy models - Unbalanced Transportation problems and Maximization case in Transportation problem. The Assignment problem - Mathematical formulation of Assignment problem – Hungarian method – Unbalanced Assignment problem - Maximization case in Assignment problem.

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TRANSPORTATION AND ASSIGNMENT PROBLEMS

Transportation Problems:

Introduction

Transportation deals with the transportation of a commodity (single product) from 'm' sources (origins or supply or capacity centers) to 'n' destinations (sinks or demand or requirement centers). It is assumed that

- (i) Level of supply at each source and the amount of demand at each destination and
- (ii) The unit transportation cost of transportation is linear.

It is also assumed that the cost of transportation is linear.

The objective is to determine the amount to be shifted from each sources to each destination such that the total transportation cost is minimum.

Note: The transportation model also can be modified to Account for multiple commodities.

1. Mathematical Formulation of a Transportation problem:

Let us assume that there are m sources and n destinations.

Let a_i be the supply (capacity) at source i , b_j be the demand at destination j , c_{ij} be the unit transportation cost from source i to destination j and x_{ij} be the number of units shifted from sources i to destination j .

Then the transportation problems can be expressed mathematically as

$$\text{Minimize } Z = \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij}$$

Subject to the constraints

$$\begin{aligned} \sum_{j=1}^n x_{ij} &= a_i, & i &= 1, 2, 3, \dots, m \\ \sum_{i=1}^m x_{ij} &= b_j, & j &= 1, 2, 3, \dots, N. \end{aligned}$$

And $x_{ij} \geq 0$, for all i and j .

Note 1: The two sets of constraints will be consistent if

$$\sum_{i=1}^m a_i = \sum_{j=1}^n b_j$$

(total supply) (total demand)

Which is the necessary and sufficient condition for a transportation problems to have a feasible solution. Problems satisfying this condition are **balanced transportation problems**.

Note 2: If $\sum a_i \neq \sum b_j$

Note 3: For any transportation problems, the coefficient of all x_{ij} in the constraints are unity.

Note 4: The objective function and the constraints being all linear, the transportation problems is a special class of linear programming problem. Therefore it can be solved by simplex method. But the number of variables being large, there will be too many calculations. So we can look for some other technique which would be simpler than the usual simplex method.

Standard transportation table:

Transportation problem is explicitly represented by the following transportation table.

		Destination							
		D_1	D_1	D_1	D_1	D_1	supply
Source	S_1	C_{11}	C_{12}	C_{13}		C_{1j}		C_{1n}	a_1
	S_1	C_{21}	C_{22}	C_{23}		C_{2j}		C_{2n}	a_2
									.
									.
	S_1	C_{i1}	C_{i2}			C_{ij}		C_{in}	.
	S_1	C_{m1}	C_{m2}			C_{mj}		C_{mn}	a_n
Demand		b_1	b_2	b_3	b_n	$\sum a_i = \sum b_j$

Definition 4: A basic feasible solution that contains less than $m + n - 1$ non-negative allocations is said to be a degenerate basic feasible solution.

Definition 5: A feasible solution (not necessarily basic) is said to be an **optimal solution** if it minimize is atmost $m + n - 1$.

Note: The number of non-basic variables in an $m \times n$ balanced transportation problem is almost $m + n - 1$.

Note: The number of non-basic variables in an $m \times n$ balanced transportation problem is atleast $mn - (m + n - 1)$.

II. Methods for finding initial basic feasible solution

The transportation problems has a solution is and only if the problem is balanced. Therefore before starting to find the initial basic feasible solution, check whether the given transportation problem is balanced. If not once has to balance the transportation problems first. The way to doing this is discussed in section 7.4 page 7.40. In this section all the given transportation problems are balanced.

Method I: North west corner rule:

Step I: The first assignment is made in the cell occupying the upper left-hand (north-west) corner of the transportation table. The maximum possible amount is allocated there. That is $x_{11} = \min \{a_1, b_1\}$.

Case (i): If $\min \{a_1, b_1\} = a_1$, then put $x_{11} = a_1$, decrease b_1 by a_1 and move vertically to the 2nd row (i.e.,) to the cell (2, 1) cross out the first row.

Case (ii): If $\min \{a_1, b_1\} = b_1$, then put $x_{11} = b_1$, decrease a_1 by b_1 and move horizontally right (i.e.,) to the cell (2, 1) cross out the first column.

Case (iii): If $\min \{a_1, b_1\} = a_1 = b_1$, the put $x_{11} = a_1 = b_1$ and move diagonally to the cell (2, 2) cross out the first row and the first column.

Step 2: Repeat the procedure until all the rim requirements are satisfied.

Method 2: Lest cost method (or) Matrix minima method (or) Lowest cost entry

method:

Step 1: Identify the cell with smallest cost and allocate $x_{ij} = \min \{a_i, b_j\}$

Case (i): If $\min \{a_i, b_j\} = a_i$, then put $x_{ij} = a_i$, cross out the i th row and decrease b_j by a_i , go to step(2).

Case (ii): If $\min \{a_i, b_j\} = b_j$, then put $x_{ij} = b_j$, cross out the j th column and decrease a_j by b_j , go to step(2).

Case (iii): If $\min \{a_i, b_j\} = a_i = b_j$, then put $x_{ij} = a_i = b_j$, cross out either i th row and j th column but not both, go to step(2).

Step 2: Repeat step (1) for the resulting reduced transportation table until all the rim requirements are satisfied.

Method 3: Vogel's approximation method (VAM) (or) Unit cost penalty method:

Step 1: Find the difference (penalty) between the smallest and next smallest costs in each row (column) and write them in brackets against the corresponding row (column).

Step 2: Identify the row (or) column with large penalty. If a tie occurs, break the tie arbitrarily. Choose the cell with smallest cost in that selected row or column and allocate as much as possible to this cell and cross out the satisfied row or column and go to step (3).

Step 3: Again compute the column and row penalties for the reduced transportation table and then go to step (2). Repeat the procedure until all the rim requirements are satisfied.

Example 1: Determine basic feasible solution to the following transportation problems using North West Corner Rule:

		Sink				
		A	B	C	D	E
Origin	P	2	11	10	3	7
	Q	1	4	7	2	1
	R	3	9	4	8	12
Demand		3	3	4	5	6
		Supply				
		4	8	9		

[MU. BE. Apr 94]

Solution:

Since $a_i = b_j = 21$, the given problem is balanced. \therefore There exists a feasible solution to the transportation problem.

2	11	10	3	7	4
3					
1	4	7	2	1	8
3	9	4	8	12	9

3	3	4	5	6
---	---	---	---	---

Following North West Corner rule, the first allocation is made in the cell(1,1)

Here $x_{11} = \min \{a_1, b_1\} = \min \{4, 3\} = 3$

Allocate 3 to the cell(1,1) and decrease 4 by 3 i.e., $4 - 3 = 1$

As the first column is satisfied, we cross out the first column and the resulting reduced Transportation table is

11 1	10	3	7	1
4	7	2	1	8
9	4	8	12	9
3	4	5	6	

Here the North West Corner cell is (1,2).

So allocate $x_{11} = \min \{1, 3\} = 1$ to the cell (1,2) and move vertically to cell (2, 2). The resulting transportation table is

4 2	7	2	1	8
9	4	8	12	9
2	4	5	6	

Allocate $x_{22} = \min \{8, 2\} = 2$ to the cell (2, 2) and move horizontally to cell (2, 3). The resulting transportation table is

7 4	2	1	6
4	8	12	9

4	5	6
---	---	---

Allocate $x_{23} = \min \{6, 4\} = 4$ and move horizontally to cell (2, 4). The resulting reduced transportation table is

2	1	2
2		
8	12	9
5	6	

Allocate $x_{24} = \min \{2, 5\} = 2$ and move vertically to cell (3, 4). The resulting reduced transportation table is

8	12	9
3		
3	6	

Allocate $x_{34} = \min \{9, 3\} = 3$ and move horizontally to cell (3, 5).which is

12		6
6		
6		

Allocate $x_{35} = \min \{6, 6\} = 6$

Finally the initial basic feasible solution is as shown in the following table.

2	11	10	3	7
3	1			
1	4	7	2	1
	2	4	2	
3	9	4	8	12
			3	6

From this table we see that the number of positive independent allocations is equal to $m + n - 1 = 3 + 5 - 1 = 7$. This ensures that the solution is non degenerate basic feasible.

$$\begin{aligned}
 \therefore \text{The initial transportation cost} &= \text{Rs. } 2 \times 3 + 11 \times 1 + 4 \times 2 + 7 \times 4 + 2 \times 2 + 8 \times 3 \\
 &\quad + 12 \times 6 \\
 &= \text{Rs. } 153/-
 \end{aligned}$$

Example 2:

Find the initial basic feasible solution for the following transportation problem by Least Cost Method.

		To				Supply
		1	2	1	4	
From		3	3	2	1	50
		4	2	5	9	20
		Demand	20	40	30	10

[MU. BE. Apr 95, BE. Nov 96]

Solution:

Since $\sum a_i = \sum b_j = 100$, the given TPP is balanced. There exists a feasible solution to the transportation problem.

1	2	1	4	30
20				
3	3	2	1	50
4	2	5	9	20
	20	40	30	10

By least cost method, $\min c_{ij} = c_{11} = c_{13} = c_{24} = 1$

Since more than one cell having the same minimum c_{ij} , break the tie.

Let us choose the cell (1,1) and allocate $x_{11} = \min \{a_1, b_1\} = \min \{30, 20\} = 20$ and cross out the satisfied column and decrease 30 by 20.

The resulting reduced transportation table is

2	1	4	10
	10		
3	2	1	50

2	5	9	20
40	30	10	

Here $\min c_{ij} = c_{13} = c_{24} = 1$. Choose the cell (1,3) and allocate $x_{13} = \min \{a_1, b_3\} = \min \{10, 30\} = 10$ and cross out the satisfied row.

The resulting reduced transportation table is

3	2	1	50
		10	
2	5	9	20
40	20	10	

Here $\min c_{ij} = c_{24} = 1$

∴ Allocate $x_{24} = \min \{a_2, b_4\} = \min (50, 10) = 10$ and cross out the satisfied column.

The resulting transportation is

3	2	40
	20	
2	5	20
40	20	

Here $c_{ij} = c_{23} = c_{32} = 2$. Choose the cell (2,3) and allocate $x_{23} = \min \{a_2, b_3\} = \min (40, 20) = 10$ and cross out the satisfied column.

The resulting reduced transportation table is

3	20
2	20
40	

Here $\min c_{ij} = c_{32} = 2$. Choose the cell (3,2) and allocate $x_{32} = \min \{a_3, b_2\} = \min (20, 40) = 20$ and cross out the satisfied row.

The resulting reduced transportation table is

3	20
20	
20	

Finally the initial basic feasible solution is as shown in the following table.

1 20	2	1 10	4
3	3 20	2 20	1 10
4	2 20	5	9

From this table we see that the number of positive independent allocations is equal to

$m + n - 1 = 3 + 4 - 1 = 6$. This ensures that the solution is non degenerate basic feasible.

\therefore The initial transportation = Rs. $1 \times 20 + 1 \times 10 + 3 \times 20 + 2 \times 20$

Cost $1 \times 10 + 2 \times 20$

= Rs. $20 + 10 + 60 + 40 + 10 + 40$

= Rs. 180/-

Example 3:

Find the initial basic feasible solution for the following transportation problem by VAM.

Distribution centres

		D_1	D_1	D_1	D_1	Availability
Origin	S_1	11	13	17	14	250
	S_2	16	18	14	10	300
	S_3	21	24	13	10	400
Requirements		200	225	275	250	

Solution:

Since $\sum a_i = \sum b_j = 100$, the given is balanced. \therefore There exists a feasible solution to this problem.

11 200	13	17	14	250 (2)
16	18	14	10	300 (4)
21	24	13	10	400 (3)
200	225	275	250	

(5) (5) (1) (0)

First let us find the difference (penalty) between the smallest and next smallest costs in each row and column and write them in brackets against the respective rows and columns.

The largest of these differences is (5) and is associated with the first two columns of the transportation table. We choose the first column arbitrarily.

In this selected column, the cell (1,1) has the minimum unit transportation cost $c_{11} = 11$.

∴ Allocate $x_{11} = \min(250, 200) = 200$ to this cell (1,1) and decrease 250 by 200 and cross out the satisfied column.

The resulting reduced transportation table is

13 50	17	14	50	(1)
18	14	10	300	(4)
24	13	10	400	(3)
225	275	250		
(5)	(1)	(0)		

The row and column differences are now computed for this reduced transportation table. The largest of these is (5) which is associated with the second column. Since $c_{12} = 13$ is the minimum cost, we allocate $x_{12} = \min(50, 225) = 50$ to the cell (1,2) and decrease 225 by 50 and cross out the satisfied row.

Continuing in this manner, the subsequent reduced transportation tables and the differences for the surviving rows and columns are shown below:

18 175	14	10	300	(4)
24	13	10	400	(3)
175	275	250		
(6)	(1)	(0)		
(i)				

14	10	125	(4)
	125		
13	10	400	(3)
		250	
(1)	(0)		

(ii)

14	10	400
	125	
275	125	

(iii)

13	275
275	
275	

(iv)

Finally the initial basic feasible solution is as shown in the following table.

11	13	17	14
200	50		
16	18	14	10
	175		125
21	24	13	10
		275	125

From this table we see that the number of positive independent allocation is equal to

$m + n - 1 = 3 + 4 - 1 = 6$. This ensures that the solution is non degenerate basic feasible.

\therefore The initial transportation = Rs. $11 \times 200 + 13 \times 50 + 18 \times 175 +$

$$\begin{aligned}\text{Cost} &= + 10 \times 125 + 13 \times 275 + 10 \times 125 \\ &= \text{Rs. } 12075/-\end{aligned}$$

Example 4:

Find the starting solution of the following transportation model

1	2	6	7
0	4	2	12
3	1	5	11
10	10	10	

Using (i). *North West Corner rule*

(ii). *Least Cost method*

(iii). *Vogel's approximation method.*

Solution:

Since $\sum a_i = \sum b_j = 100$, the given Transportation problem is balanced. \therefore There exists a basic feasible solution to this problem.

(i). **North West Corner rule:** Using this method, the allocation are shown in the tables below:

1	2	6	7
7			
0	4	2	12
3	1	5	11
10	10	10	

(i)

0	4	2	12
3			
3	1	5	11
3	10	10	

(ii)

4	2	9
8		
1	5	11
10	10	

(iii)

1	5	11
1		
1	10	

5	10	10
10		

(iv)

(v)

The starting solution is as shown in the following table

1	2	6
7		
0	4	2
3	9	
3	1	5
	1	10

$$\therefore \text{The initial transportation cost} = \text{Rs. } 1 \times 7 + 0 \times 3 + 4 \times 9 + 1 \times 1 + 5 \times 10$$

$$= \text{Rs. } 94/-$$

(ii). **Least Cost method:** Using this method, the allocation are as shown in the table below:

1	2	6	7
		7	
0	4	2	12
10			
3	1	5	11
10	10	10	

(i)

2	6	7
4	2	2
1	5	11
10		10

(ii)

6	7
2	2
5	1
10	

(iii)

6	7
5	1
8	

(iv)

6	7
7	

(v)

The starting solution is as shown in the following table:

1	2	6	7
0	4	2	2
3	1	5	1
10		10	1

$$\therefore \text{The initial transportation cost} = \text{Rs. } 6 \times 7 + 0 \times 10 + 2 \times 2 + 1 \times 10 + 5 \times 1$$

$$= \text{Rs. } 61/-$$

(iii). **Vogel's approximation Method:** Using this method, the allocations are shown in the table below:

1	2	6	7 (1)
0	4	2	12 (2)
3	1	5	11 (2)
10	10	10	
(1)	(1)	(3)	

(i)

1	2	7 (1)
0	4	2 (4)
3	1	11 (2)
10	10	
(1)	(1)	

(ii)

1	2	7 (1)
3	1	11 (2)
8	10	
(2)	(1)	

(iii)

1	7	7
3		1
8		

(iv)

3	1	1
1		

(v)

The starting solution is as shown in the following table:

1 7	2	6
0 2	4	2 10
3 1	1 10	5

$$\therefore \text{The initial transportation cost} = \text{Rs. } 1 \times 7 + 0 \times 2 + 2 \times 10 + 3 \times 1 + 1 \times 10$$

$$= \text{Rs. } 40/-$$

Note: For the above problem, the number of positive allocation in independent positions is $(m + n - 1)$ (i.e., $m + n - 1 = 3 + 3 - 1 = 5$). This ensures that the given problem has a non-degenerate basic feasible solution by using all the three methods. This need not be the case in all the problems.

Transportation Algorithm (or) MODI Method (modified distribution method) (Test for optimal solution).

Step 1: Find the initial basic feasible solution of the given problems by Northwest Corner rule (or) Least Cost method or VAM.

Step 2: Check the number of occupied cells. If these are less than $m + n - 1$, there exists degeneracy and we introduce a very small positive assignment of $\epsilon (\approx 0)$ in suitable independent positions, so that the number of occupied cells is exactly equal to $m + n - 1$.

Step 3: Find the set of values u_i, v_j ($i = 1, 2, 3, \dots, m; j = 1, 2, 3, \dots, n$) from the relation $c_{ij} = u_i + v_j$ for each occupied cell (i, j) , by starting initially with $u_i = 0$ or $v_j = 0$ preferably for which the corresponding row or column has maximum number of individual allocations.

Step 4: Find $u_i + v_j$ for each unoccupied cell (i, j) and enter at the upper right corner of the corresponding cell (i, j) .

Step 5: Find the cell evaluations $d_{ij} = c_{ij} - (u_i + v_j)$ (d_{ij} = upper left – upper right) for each unoccupied cell (i, j) and enter at the lower right corner of the corresponding cell (i, j) .

Step 6: Examine the cell evaluations d_{ij} for all unoccupied cells (i, j) and conclude that

- (i) If all $d_{ij} > 0$, then the solution under the test is optimal and unique.
- (ii) If all $d_{ij} > 0$, with atleast one $d_{ij} = 0$, then the solution under the test is optimal and an alternative optimal solution exists.
- (iii) If atleast one $d_{ij} < 0$, then the solution is not optimal. Go to the next step.

Step 7: Form a new $B > F > S$ by giving maximum allocation to the cell for which d_{ij} is most negative by making an occupied cell empty. For that draw a closed path consisting of horizontal and vertical lines beginning and ending at the cell for which d_{ij} is most negative and having its **other corners at some allocated cells**. Along this closed loop indicate $+\theta$ and $-\theta$ alternatively at the corners. Choose minimum of the allocations from the cells having $-\theta$. Add this minimum allocation to the cells with $+\theta$ and subtract this minimum allocation from the allocation to the cells with $-\theta$.

Step 8: Repeat steps (2) to (6) to test the optimality of this new basic feasible solution.

Step 9: Continue the above procedure till an optimum solution is attained.

Note: The Vogels approximation method (VAM) takes into account not only the least cost c_{ij} but also the costs that just exceed the least cost c_{ij} and therefore yields better initial solution than obtained from other methods in general. This can be justified by the above example (4). So to find the initial solution, give preference to VAM unless otherwise specified.

Example 1: Solve the transportation problem:

	1	2	3	4	Supply
I	21	16	25	13	11
II	17	18	14	23	13
III	32	27	18	41	19
Demand	6	10	12	15	

Solution: Since $\sum a_i = \sum b_j = 43$, the given transportation problem is balanced. There exists a basic feasible solution to this problem.

By Vogel's approximation method, the initial solution is shown in the following table.

21	16	25	13	(3) - - -
			11	
17	18	14	23	(3) (3) (3) (3)
6	3		4	
32	27	18	41	(9) (9) (9) (9)
	7	12		

(4)	(2)	(4)	(10)
(15)	(9)	(4)	(18)
(15)	(9)	(4)	
	(9)	(4)	

That is

21	16	25	13 11
17 6	18 3	14	23 4
32	27 7	18 12	41

From this table, we see that the number of non-negative independent allocations is $(m + n - 1) = (3+4-1) = 6$. Hence the solution is non-degenerate basic feasible.

∴ The initial transportation cost.

$$= \text{Rs. } 13 \times 11 + 17 \times 6 + 18 \times 3 + 23 \times 4 + 27 \times 7 + 18 \times 12$$

$$= \text{Rs. } 796/-$$

To find the optimal solution

Consider the above transportation table. Since $m+n-1=6$, we apply MODI method,

Now we determine a set of values u_i and v_j for each occupied cell (i,j) by using the relation $c_{ij} = u_i + v_j$. As the 2nd row contains maximum number of allocations, we choose $u_2=0$.

Therefore

$$C_{21} = u_2 + v_1 \Rightarrow 17 = 0 + v_1 \Rightarrow v_1 = 17$$

$$C_{22} = u_2 + v_2 \Rightarrow 18 = 0 + v_2 \Rightarrow v_2 = 18$$

$$C_{24} = u_2 + v_4 \Rightarrow 23 = 0 + v_4 \Rightarrow v_4 = 23$$

$$C_{14} = u_1 + v_4 \Rightarrow 13 = u_1 + 23 \Rightarrow u_1 = -10$$

$$C_{32} = u_3 + v_2 \Rightarrow 27 = u_3 + 18 \Rightarrow u_3 = 9$$

$$C_{33} = u_3 + v_3 \Rightarrow 18 = 9 + v_3 \Rightarrow v_3 = 9$$

Thus we have the following transportation table:

21	16	25	13 11
----	----	----	-----------------

 $u_1 = -10$

17	18	14	23	$u_2 = 0$
6	3		4	
32	27	18	41	$u_3 = 9$
	7	12		

$$v_1 = 17 \quad v_2 = 18 \quad v_3 = 9 \quad v_4 = 23$$

Now we find $u_i + v_j$ for each unoccupied cell (i,j) and enter at the upper right corner of the corresponding unoccupied cell(i,j).

Then we find the cell evaluations $d_{ij} = c_{ij} - (u_i + v_j)$ (ie., upper left corner – upper right corner) for each unoccupied cell (i,j) and enter at the lower right corner of the corresponding unoccupied cell (i,j).

21	7	16	8	25	-1	13	$u_1 = -10$
	14		8		26	11	
17		18		14	9	23	$u_2 = 0$
6		3			5	4	
32	26	27		18		41	$u_3 = 9$
	6	7		12		9	

$$v_1 = 17 \quad v_2 = 18 \quad v_3 = 9 \quad v_4 = 23$$

Since all $d_{ij} > 0$, with $d_{32} = 0$, the current solution is optimal and unique.

∴ The optimum allocation schedule is given by $x_{14} = 11$, $x_{21} = 6$, $x_{22} = 3$, $x_{24} = 4$, $x_{32} = 7$, $x_{33} = 12$, and the optimum (minimum) transportation cost
 = Rs. $13 \times 11 + 17 \times 6 + 18 \times 3 + 23 \times 4 + 27 \times 7 + 18 \times 12$
 = Rs. 796/-

Example 2:

Obtain an optimum feasible solution to the following transportation problem:

		To			Available
From		7	3	2	2
		2	1	3	3
		3	4	6	5
Demand		4	1	5	10

Solution:

Since $\sum a_i = \sum b_j = 43$, the given transportation problem is balanced. ∴ There exists a basic feasible solution to this problem.

By Vogel's approximation method, the initial solution is as shown in the following table:

7	3	2	(1) (5)
		2	
2	1	3	(1) (1) (1)
	1	2	
3	4	6	(1) (3) (3)
4		1	

(1) (2) (1)
 (1) (1) (1)
 (1) (3) (3)

That is

7	3	2
		2
2	1	3
	1	2
3	4	6
4		1

From this table we see that the number of non-negative allocation is

$$m + n - 1 = (3 + 3 - 1) = 5.$$

Hence the solution is non-degenerate basic feasible

$$\begin{aligned} \therefore \text{The initial transportation cost} &= \text{Rs. } 2 \times 2 + 1 \times 1 + 3 \times 2 + 3 \times 4 + 6 \times 1 \\ &= \text{Rs. } 29/- \end{aligned}$$

For optimality: since the number of non – negative independent allocation is $m + n - 1$, we apply MODI method.

Since the third column contains maximum number of allocations, we choose $v_3 = 0$.

Now we determine a set of values u_i and v_j by using the occupied cells and the relation $c_{ij} = u_i + v_j$.

That is

7	-1	3	0	2	$u_1 = 2$
				2	
2		1		3	$u_2 = 3$
		1		2	

3	4	6	
	4		1
$v_1 = -3$	$v_2 = -2$	$v_3 = 0$	$u_3 = 6$

Now we find $u_i + v_j$ for each unoccupied cell (i, j) and enter at the corresponding unoccupied cell (I,j).

Then we find the cell evaluations $d_{ij} = c_{ij} - (u_i + v_j)$ for each unoccupied cell (i, j) and enter at the lower right corner of the corresponding unoccupied cell (i, j).

Thus we get the following table

7	-1	3	0	2	
	8		3		2
2	0	1		3	
	2		1		2
3		4	4	6	
	4		0		1
$v_1 = -3$	$v_2 = -2$	$v_3 = 0$			$u_3 = 6$

Since all $d_{ij} > 0$, with $d_{32} = 0$, the current solution is optimal and there exists an alternative optimal solution.

∴ The optimum allocation schedule is given by $x_{13} = 2$, $x_{32} = 1$, $x_{23} = 2$, $x_{31} = 4$, $x_{33} = 1$, and the optimum (minimum) transportation cost

$$= \text{Rs. } 2 \times 2 + 1 \times 1 + 3 \times 2 + 3 \times 4 + 6 \times 1 = \text{Rs. } 29/-$$

Example 3: Find the optimal transportation cost of the following matrix using least cost method for finding the critical solution.

		A	B	C	D	E	Available
Factory	P	4	1	2	6	9	100
	Q	6	4	3	5	7	120
	R	5	2	6	4	8	120

Demand

40	50	70	90	90
----	----	----	----	----

Solution:

Since $\sum a_i = \sum b_j = 340$, the given transportation problem is balanced. \therefore There exists a basic feasible solution to this problem.

By using Least cost method, the initial solution is shown in the following table:

4	1	2	6	9
	50	50		
6	4	3	5	7
10		20		90
5	2	6	4	8
30			90	

\therefore The initial transportation cost = Rs. $1 \times 50 + 2 \times 50 + 6 \times 10 + 3 \times 20 + 7 \times 90$

$+ 5 \times 30 + 4 \times 90$

= Rs. 1410/-

For optimality: Since the number of non – negative independent allocations is $(m + n - 1)$, we apply MODI method:

4	5	1	2	6	4	9	6	$u_1 = -1$
	-1	50	50		2		3	
6		4	2	3	5	5	7	$u_2 = 0$
10			2	20		0	90	
5		2	1	6	2	4	8	$u_3 = -1$
30			1		4	90	2	

$$v_1 = 6 \quad v_2 = 2 \quad v_3 = 3 \quad v_4 = 5 \quad v_5 = 7$$

Since $d_{11} = -1 < 0$, the current solution is not optimal.

Now let us form a new basic feasible solution by giving maximum allocation to the cell (i,j) for which d_{ij} is most negative by making an occupied cell empty. Here the cell (1,1) having the negative value $d_{11}=-1$. We draw a closed loop consisting of horizontal and vertical lines beginning and ending at this cell (1,1) and having its other corners at some occupied cells. Along this closed loop indicate $+\theta$ and $-\theta$ alternatively at the corners. We have

4	1	2	6	9
$+\theta$	50	50	$-\theta$	
6	4	3	5	7
10		20		90
$-\theta$		$+\theta$		
5	2	6	4	8
30			90	

From the two cells (1,3), (2,1) having $+\theta$, we find that the minimum of the allocations 50,10 is 10. Add this cells with $+\theta$ and subtract this 10 to the cells with $+\theta$.

Hence the new basic feasible solution is displayed in the following table:

4	1	2	6	9
10	50	40		
6	4	3	5	7
		30		90
5	2	6	4	8
30			90	

We see that the above table satisfies the rim conditions with $(m + n - 1)$ non-negative allocations at independent position. So we apply MODI method.

4	1	2	6	3	9	6	$u_1 = 0$
10	50	40		3		3	
6	5	4	3	5	4	7	$u_2 = 1$
	1	2	30		1	90	
5	2	2	6	3	4	8	$u_3 = 1$
30		0		3	90	1	
$v_1 = 4$	$v_2 = 1$	$v_3 = 2$	$v_4 = 3$	$v_5 = 6$			

Since all $d_{ij} > 0$, with $d_{32} = 0$, the current solution is optimal and there exists an alternative optimal solution.

The optimum allocation schedule is given by $x_{11}=10$, $x_{12}=50$, $x_{13}=40$, $x_{23}=30$, $x_{25}=90$, $x_{31}=30$, $x_{34}=90$ and the optimum (minimum) transportation cost.

= Rs. $4 \times 10 + 1 \times 50 + 2 \times 40 + 3 \times 30 + 7 \times 90 + 5 \times 30 + 4 \times 90$.

= Rs. 1400/-

Degeneracy in Transportation Problems

In transportation problems, whenever the number of non-negative independent allocations is less than $m + n - 1$, the transportation problem is said to be **degenerate** one. Degeneracy may occur either at the initial stage or at an intermediate stage at some subsequent iteration.

To resolve degeneracy, we allocate an extremely small amount (close to zero) to one or more empty cells of the transportation table (generally minimum cost cells if possible), so that the total number of occupied cells becomes $(m + n - 1)$ at independent positions. We denote this small amount by ϵ (epsilon) satisfying the following conditions:

- (i) $0 < \epsilon < x_{ij}$, for all $x_{ij} > 0$
- (ii) $x_{ij} \pm \epsilon = x_{ij}$, for all $x_{ij} > 0$

The cells containing ϵ are then treated like other occupied cells and the problems is solved in the usual way. The ϵ 's are kept till the optimum solution is attained. Then we let each $\epsilon \rightarrow 0$.

Example 1: find the non-degenerate basic feasible solution for the following transportation problems using

- (i) North west corner rule
- (ii) Least cost method
- (iii) Vogel's approximation method.

		To				supply
		10	20	5	7	
From	13	9	12	8		20
	4	5	7	9		30
	14	7	1	0		40
	3	12	5	19		50
	Demand	60	60	20	10	

Solution: Since $\sum a_i = \sum b_j = 150$, the given transportation problems is balanced.

\therefore There exists a basic feasible solution to this problem.

- (i) The starting solution by NWC rule is an shown in the following table.

10	20	5	7
10			
13	9	12	8
20			
4	5	7	9
30			
14	7	1	0
	40		

3	12	5	19
	20	20	10

Since the number of non-negative allocations at independent positions is 7 which is less than $(m + n - 1) = (5 + 4 - 1) = 8$, this basic feasible solution is a degenerate one.

To resolve this degeneracy, we allocate a very small quantity ϵ to the unoccupied cell (5,1) so that the number of occupied cells becomes $(m+n-1)$ ($m + n - 1$). Hence the non-degenerate basic feasible solution is as shown in the following table.

	10	20	5	7
	10			
	13	9	12	8
	20			
	4	5	7	9
	30			
	14	7	1	0
		40		
The initial	3	12	5	19
	ϵ	20	20	10

transportation cost = Rs.

$$10 \times 10 + 13 \times 20 + 4 \times 30 + 7 \times 40 + 3 \times \epsilon + 20 \times 20 + 5 \times 20 + 19 \times 10$$

$$= \text{Rs.}(1290 = 3\text{€})$$

$$= \text{Rs. } 1290/- \text{ as } \text{€} \rightarrow 0.$$

(ii) Least cost method: Using this method the starting solution is an shown in the following table:

10	20	5	7
	10		
13	9	12	8
	20		
4	5	7	9
10	20		
14	7	1	0
	10	20	10
3	12	5	19
50			

Since the number of non-negative allocations at independent positions is $(m + n - 1) = 8$, the solution is non-degenerate basic feasible.

$$\begin{aligned} \text{The initial transportation cost} &= \text{Rs. } 20 \times 10 + 9 \times 20 + 4 \times 10 + 5 \times 20 + 7 \times 10 + 1 \times 20 + 0 \times 10 + 3 \times 50 \\ &= \text{Rs. } 760/- \end{aligned}$$

(iii) Vogel's approximation method: The starting solution by this method is an shown in the following table:

10	20	5	7
10			
13	9	12	8
	20		
4	5	7	9
	30		

14	7	1	0
	10	20	10
3	12	5	19
50			

Since the number of non-negative allocations is 7 which is less than $(m + n - 1) = (5 + 4 - 1) = 8$, this basic solution is a degenerate one.

To resolve this degeneracy, we allocate a very small quantity ϵ to the unoccupied cell(5,2) so that the number of occupied cells becomes $(m + n - 1)$. Hence the non-degenerate basic feasible solution is as shown in the following table.

10	20	5	7
10			
13	9	12	8
	20		
4	5	7	9
	30		
14	7	1	0
	10	20	10
3	12	5	19
50	ϵ		

\therefore The initial transportation cost

$$\begin{aligned}
 &= \text{Rs. } 10 \times 10 + 9 \times 20 + 5 \times 30 + 7 \times 10 + 1 \times 20 + 0 \times 10 + 3 \times 50 + 12 \times \epsilon \\
 &= \text{Rs. } (670 + 12\epsilon)
 \end{aligned}$$

$$= \text{Rs.}670/- = \text{as } \epsilon \rightarrow 0.$$

Example 2: Solve the following transportation problems using vogel's method.

	A	B	C	D	E	F	Available
1	9	12	9	6	9	10	5
2	7	3	7	7	5	5	6
3	6	5	9	11	3	11	2
4	6	8	11	2	2	10	9
Requirement	4	4	6	2	4	2	

Solution: Since $\sum a_i = \sum b_j = 22$, the given transportation problem is balanced. \therefore There exists a basic feasible solution to this problem. By Vogel's approximation method, the initial solution is as shown in the following table:

9	12	9	6	9	10
		5			
7	3	7	7	5	5
	4				2
6	5	9	11	3	11
1	ϵ	1			
6	8	11	2	2	10
3			2	4	

Since the number of non-negative allocations is 8 which is less than $(m + n - 1) = (4 + 6 - 1) = 9$, this basic solution is degenerate one.

To resolve degeneracy, we allocate a very small quantity ϵ to the cell (3,2), so that the number of occupied cells becomes $(m + n - 1)$. Hence the non-degenerate basic feasible solution is as shown in the following table.

9	12	9	6	9	10
		5			
7	3	7	7	5	5
	4				2

6 1	5 €	9 1	11	3	11
6 3	8	11	2 2	2 4	10

The initial transportation cost = Rs. $9 \times 5 + 3 \times 4 + 5 \times 2 + 6 \times 1 + 5 \times € + 9 \times 1$
 $+ 6 \times 3 + 2 \times 2 + 2 \times 4$
 $= \text{Rs.}(112+5€) = \text{Rs.}112/-, € \rightarrow 0.$

To find the optimal solution

Now the number of non-negative allocations at independent positions is $(m + n - 1)$. We apply the MODI method.

9 6	12 5	9 5	6 2	9 2	10 7	$u_1 = 0$
3	7		4	7	3	
7 4	3 4	7 7	7 0	5 0	5 2	$u_2 = -2$
3		0	7	5		
6 1	5 €	9 1	11 2	3 2	11 7	$u_3 = 0$
			9	1	4	
6 3	8 5	11 9	2 2	2 4	10 7	$u_4 = 0$
	3	2			3	
$v_1 = 6$	$v_2 = 5$	$v_3 = 9$	$v_4 = 2$	$v_5 = 2$	$v_6 = 7$	

Since all $d_{ij} > 0$ with $d_{23} = 0$, the solution under the test is optimal and an alternative optimal solution is also exists.

\therefore The optimum allocation schedule is given by $x_{14}=5, x_{22}=4, x_{26}=2, x_{31}=1, x_{32}=€$, $x_{33}=1, x_{41}=3, x_{44}=2, x_{45}=4$ and the optimum(minimum) transportation cost is

$$= \text{Rs.} 9 \times 5 + 3 \times 4 + 5 \times 2 + 6 \times 1 + 5 \times € + 9 \times 1 + 6 \times 3 + 2 \times 2 + 2 \times 4$$

$$=Rs. (112+5\epsilon)$$

$$=Rs. 112 \text{ as } \epsilon \rightarrow 0.$$

Example 3: Solve the following transportation problem to minimize the total cost of transportation.

	To				Supply
	1	2	3	4	
From	4	3	2	0	8
	0	2	2	1	10
Demand	4	6	8	6	

Solution: Since $\sum a_i = \sum b_j = 24$, the given transportation problem is balanced. \therefore There exists a basic feasible solution to this problem.

By using Vogel's approximation method, the initial solution is as shown in the following table:

1	2	3	4
	6		
4	3	2	0
		2	6
0	2	2	1
4		6	
4			

Since the number of non-negative allocations is 5, which is less than $(m + n - 1) = (3 + 4 - 1) = 6$, this basic feasible solution is degenerate.

To resolve degeneracy, we allocate a very small quantity ϵ to the cell (1,4), so that the number of occupied cells becomes $(m + n - 1)$. Hence the non-degenerate basic feasible solution is given in the following table

1	2	3	4
	6		
4	3	2	0
		2	6

0	2	2	1
4	€	6	

$$\begin{aligned}
 \therefore \text{The initial transportation cost} &= \text{Rs. } 2 \times 6 + 2 \times 2 + 0 \times 6 + 0 \times 4 + 2 \times € + 2 \times 6 \\
 &= \text{Rs. } (28 + 2€) \\
 &= \text{Rs. } 28/-, \text{ as } € \rightarrow 0.
 \end{aligned}$$

To find the optimum solution:

Now the number of non-negative allocations at independent positions is $(m + n - 1)$. We apply MODI method.

1	0	2	3	4	0	$u_1 = 0$
	1	6		1	4	
4	0	3	2	2	0	$u_2 = 0$
	4		1	2	6	
0		2	2	1	0	$u_3 = 0$
4		€	6		1	
$v_1 = 0 \quad v_2 = 2 \quad v_3 = 2 \quad v_4 = 0$						

Since all $d_{ij} > 0$ the solution under the test is optimal and unique.

\therefore The optimal allocation schedule is given by $x_{12} = 6$, $x_{23} = 2$, $x_{24} = 6$, $x_{31} = 4$, $x_{32} = €$, $x_{33} = 6$ and the optimum (minimum) transportation cost

$$\begin{aligned}
 &= \text{Rs. } 2 \times 6 + 2 \times 2 + 0 \times 6 + 0 \times 4 + 2 \times € + 2 \times 6 \\
 &= \text{Rs. } (28 + 2€) = \text{Rs. } 28, \text{ as } € \rightarrow 0.
 \end{aligned}$$

Example 5:

Solve the following transportation problem to minimize the total cost of transportation.

Destination

		<i>1</i>	<i>2</i>	<i>3</i>	<i>4</i>	<i>supply</i>
<i>1</i>		14	56	48	27	70
<i>Origin 2</i>		82	35	21	81	47
<i>3</i>		99	31	71	63	93
<i>Demand</i>		70	35	45	60	210

Solution:

Since $\sum a_i = \sum b_j = 210$, the given transportation problem is balanced. \therefore There exists a basic feasible solution to this problem.

By using Vogel's approximation method, the initial solution is as shown in the following table:

14 70	56	48	27
82	35	21 45	81 2
99	31 35	71	63 58

Since the number of non-negative allocations is 5, which is less than $(m + n - 1) = (3+4-1) = 6$, this basic feasible solution is degenerate.

To resolve degeneracy, we allocate a very small quantity ϵ to the cell (1,4). So that the number of occupied cells becomes $(m + n - 1)$. Hence the non-degenerate basic feasible solution is given in the following table.

14 70	56	48	27 ϵ
82	35	21 45	81 2
99	31 35	71	63 58

To find the optimum solution:

Now the number of non-negative allocations at independent positions is $(m + n - 1) = 6$. We apply MODI method.

14	70	56	-5	48	-33	27	€	$u_1 = 27$
		61		81				
82	68	35	49	21	45	81	2	$u_2 = 81$
	14		-14					
99	50	31	35	71	3	63	58	$u_3 = 63$
	49			68				
$v_1 = -13$		$v_2 = -32$		$v_3 = 60$		$v_4 = 0$		

Since $d_{22} = -14 < 0$, the solution under the test is not optimal.

Now let us form a new basic feasible solution by giving maximum allocation to the cell (2,2) by making an occupied cell empty. We draw a closed loop consisting of horizontal and vertical lines beginning and ending at this cell (2,2) and having its other corners at some occupied cells. Along this closed loop, indicate $+\theta$ and $-\theta$ alternatively at the corners.

14	70	56		48		27	€
82		35		21	45	81	2
			$+\theta$				$-\theta$
99		31		71		63	
			$-\theta$				$+\theta$
		35				58	

From the two cells (2,4),(3,2) having $-\theta$ we find that the minimum of the allocations 2,35 is 2. Add this 2 to the cells with $+\theta$ and subtract this 2 to the cells with $-\theta$. Hence the new basic feasible solution is given by

14 70	56	48	27 €
82	35 2	21 45	81
99	31 33	71	63 60

We see that the above table satisfies the rim conditions with $(m + n - 1)$ non-negative allocations at independent position. We apply MODI method for optimality.

14 70	56 -5 61	48 -19 81	27 €	$u_1 = -40$
82 54 28	35 2	21 45	81 67 14	$u_2 = 0$
99 50 49	31 33	71 17 54	63 60	$u_3 = -4$
$v_1 = 54 \quad v_2 = 35 \quad v_3 = 21 \quad v_4 = 67$				

Since $d_{ij} > 0$, the solution under the test is optimal.

∴ The optimal allocation schedule is given by $x_{11} = 70$, $x_{14} = €$, $x_{22} = 2$, $x_{23} = 45$, $x_{32} = 33$, $x_{34} = 60$ and the optimum (minimum) transportation cost

$$= \text{Rs. } 14 \times 70 + 27 \times € + 35 \times 2 + 21 \times 45 + 31 \times 33 + 63 \times 60$$

$$= \text{Rs. } 6798/- \text{ as } € \rightarrow 0.$$

Unbalanced Transportation Problems

If the given transportation problems is unbalanced one, i.e., if $\sum a_i \neq \sum b_j$, then convert this into a balanced one by introducing a dummy source or dummy destination with zero cost vector (zero unit transportation costs) as the case may be and then solve by usual method.

When the total supply is greater than the total demand, a dummy destination is included in the matrix with zero cost vectors. The excess supply is entered as a rim requirement for the dummy destination.

When the total demand is greater than the total supply, a dummy source is included in the matrix with zero cost vectors. The excess demand is entered as rim requirements for the dummy source.

Example 1: Solve the transportation problem

		Destination				
		A	B	C	D	supply
Source	1	11	20	7	8	50
	2	21	16	20	12	40
	3	8	12	18	9	70
Demand		30	25	35	40	

Solution: Since the total supply ($\sum a_i = 160$) is greater than the total demand ($\sum b_j = 130$), the given problem is an unbalanced transportation problem. To convert this into a balanced one, we introduce a dummy destination E with zero unit transportation costs and having demand equal to $160 - 130 = 30$ units.

∴ The given problem becomes

		Destination					
		A	B	C	D	E	supply
Source	1	11	20	7	8	0	50
	2	21	16	20	12	0	40
	3	8	12	18	9	0	70
Demand		30	25	35	40	30	160

By using VAM the initial solution is as shown in the following table

11	20	7	8	0
		35	15	

21	16	20	12 10	0 30
8 30	12 25	18	9 15	0

∴ The initial transportation cost

$$= \text{Rs. } 7 \times 35 + 8 \times 15 + 12 \times 10 + 0 \times 30 + 8 \times 30 + 12 \times 25 + 9 \times 15$$

$$= \text{Rs. } 1160/-$$

For Optimality: Since the number non-negative allocations at independent position is $(m + n - 1)$, we apply the MODI method.

11	7	20	11	7	8	0	-4	$u_1 = 8$
	4		9	35	15		4	
21	11	16	15	20	11	12	0	$u_2 = 12$
	10		1		9	10	30	
8	30	12	25	18	8	9	0	$u_3 = 9$
				10	15		3	
$v_1 = -1$	$v_2 = 3$	$v_3 = -1$	$v_4 = 0$	$v_5 = -12$				

Since all $d_{ij} > 0$, the solution under the test is optimum and unique.

∴ The optimum allocation schedule is $x_{13} = 35$, $x_{14} = 15$, $x_{24} = 10$, $x_{25} = 30$, $x_{31} = 30$, $x_{32} = 25$, $x_{34} = 15$

It can be noted that $x_{25} = 30$ means that 30 units are dispatched from source 2 to the dummy destination E. In other words, 30 units are left undispached from source 2.

The optimum (minimum) transportation cost

$$= \text{Rs. } 7 \times 35 + 8 \times 15 + 12 \times 10 + 0 \times 30 + 8 \times 30 + 12 \times 25 + 9 \times 15$$

$$= \text{Rs. } 1160/-$$

Example 2: Solve the transportation problem with unit transportation costs, demands and supplies as given below:

		Destination				
		D ₁	D ₂	D ₃	D ₄	Supply
Source	S ₁	6	1	9	3	70
	S ₂	11	5	2	8	55
	S ₃	10	12	4	7	70
Demand		85	35	50	45	

Solution: Since the total demand ($\sum b_j = 215$) is greater than the total supply ($\sum a_i = 195$), the given problem is unbalanced transportation problem. To convert this into a balanced one, we introduce a dummy source S_4 with zero unit transportation costs and having supply equal to $215 - 195 = 20$ units. \therefore The given problems becomes

		Destination				
		D ₁	D ₂	D ₃	D ₄	Supply
Source	S ₁	6	1	9	3	70
	S ₂	11	5	2	8	55
	S ₃	10	12	4	7	70
	S ₄	0	0	0	0	20
Demand		85	35	50	45	215

As this problem is balanced, there exists a basic feasible solution to this problem. By using Vogel's approximation method, the initial solution is as shown in the following table.

6	1	9	3
65	5		
11	5	2	8
	30	25	
10	12	4	7
		25	45
0	0	0	0
20			

\therefore The initial transportation cost

$$= \text{Rs. } 6 \times 65 + 1 \times 5 + 5 \times 30 + 2 \times 25 + 4 \times 25 + 7 \times 45 + 0 \times 20$$

$$= \text{Rs. } 1010/-$$

For optimality: Since number of non-negative allocations at independent positions is $(m + n - 1)$, we apply the MODI method.

6	65	1	5	9	-2	3	1	$u_1 = 6$
					11		2	
11	10	5	30	2	25	8	5	$u_2 = 10$
	1						3	
10	12	12	7	4	25	7	45	$u_3 = 12$
	-2		5					
0	20	0	-5	0	-8	0	-5	$u_4 = 0$
			5		8		5	
$v_1 = 0$		$v_2 = -5$		$v_3 = -8$		$v_4 = -5$		

Since $d_{31} = -2 < 0$, the solution under the test is not optimal.

Now let us form a new basic feasible solution by giving maximum empty. For this, we draw a closed path consisting of horizontal and vertical lines beginning and ending at this cell (3,1) and having its other corners at some occupied cells. Along this closed loop, indicate $+\theta$ and $-\theta$ alternatively at the corners.

We have,

6	65 $-\theta$	1	5 $+\theta$	9		3	
11		5	30 $-\theta$	2	25 $+\theta$	8	
10	$+\theta$	12		4	25 $-\theta$	7	45

0	0	0	0
20			

From the three cells (1,1), (2,2), (3,3) having $-\theta$, we find that the minimum of the allocations 65,30,25 is 25. Add this 25 to the cells with $+\theta$ and subtract this 25 to this cells with $-\theta$. Finally, the new feasible solution is displayed in the following table.

6	1	9	3
40	30		
11	5	2	8
	5	50	
10	12	4	7
25			45
0	0	0	0
20			

We see that the above table satisfies the rim conditions with $(m + n - 1)$ non-negative allocations at independent positions. Now we check for optimality.

6	1	9	3
40	30	-2	3
		11	0
11	5	2	8
10	5	50	7
1			1
10	12	4	7
25	5	2	45
	7	2	
0	0	0	0
20	-5	-8	-3
	5	8	3

Since all $d_{ij} > 0$ with $d_{14} = 0$, the solution under the test is optimal and an alternative optimal solution exists.

∴ The optimum allocation schedule is given by $x_{13} = 35$, $x_{14} = 15$, $x_{24} = 10$, $x_{25} = 30$, $x_{31} = 30$, $x_{32} = 25$, $x_{34} = 15$, $x_{41} = 20$.

It can be noted that $x_{41}=20$ means that 20 units are dispatched from the dummy source S_4 to the destination D_1 . In other words, 20 units are not fulfilled for the destination D_1 .

The optimum (minimum) transportation cost

$$= \text{Rs. } 6 \times 40 + 1 \times 30 + 5 \times 5 + 2 \times 50 + 10 \times 25 + 7 \times 45 + 0 \times 20$$

$$= \text{Rs. } 960/-$$

Example 3:

Solve the transportation problem with unit transportation costs in rupees, demand and supplies as given below:

		Destination			Supply(units)
		D ₁	D ₂	D ₃	
Origin	A	5	6	9	100
	B	3	5	10	75
	C	6	7	6	50
	D	6	4	10	75
Demand (units)		70	80	120	

Solution: Since the total supply ($\sum a_i = 270$), the given transportation problem is unbalanced.

To convert this into a balanced one, we introduce a dummy source D_4 with zero unit transportation costs and having demand equal to $300 - 270 = 30$ units. ∴ The given problem becomes

Destination

		D ₁	D ₂	D ₃	D ₄	Supply(units)
Origin	A	5	6	9	0	100
	B	3	5	10	0	75
	C	6	7	6	0	50
	D	6	4	10	0	75
Demand (units)		70	80	120	30	300

By using VAM the initial solution is given by

5	6	9 100	0
3 70	5 5	10	0
6	7	6 20	0 30
6	4 75	10	0

Since the number of non-negative allocations is 6, which is less than $(m + n - 1) = 4+4-1 = 7$, this basic feasible solution is degenerate.

To resolve this degeneracy, we allocate a very small quantity ϵ to the cell (2,4), so that the number of occupied cells becomes $(m + n - 1)$. Hence the non-degenerate basic feasible solution is given in the following table.

5	6	9 100	0
3 70	5 5	10	0 ϵ
6	7	6 20	0 30
6	4 75	10	0

Now the number of non-negative allocations at independent positions is $(m + n - 1)$. We apply MODI method.

5	6	6	8	9 100	0	3	$u_1 = 3$
	-1		-2			-3	
3 70	5 5	10	6	0	ϵ		$u_2 = 0$
			4				

6	3	7	5	6	0	
				20	30	$u_3 = 0$
	3		2			
6	2	4		10	5	0
		75				$u_4 = -1$
	4				5	
						1
$v_1 = 3$	$v_2 = 5$	$v_3 = 6$	$v_4 = 0$			

Since there are some $d_{ij} < 0$, the current solution is not optimal.

Since $d_{14} = -3$ is the most negative, let us form a new basic feasible solution by giving maximum allocations to the corresponding cell (1,4) by making an occupied cell empty. We draw a closed loop consisting of horizontal and vertical lines beginning and ending at this cell (1,4) and having its other corners at some occupied cells. Along this closed loop indicate $+\theta$ and $-\theta$ Alternately at the corners.

5	6	9	0	
		100		$+\theta$
		$-\theta$		
3	5	10	0	
70	5			\in
6	7	6	0	
		20	30	
		$+\theta$		$-\theta$
6	4	10	0	
	75			

From the two cells (1, 3), (3, 4) having $-\theta$, we find that the minimum of the allocations 100, 30 is 30. Add this 30 to the cells with $+\theta$ and subtract this 30 to the cells with $-\theta$. Hence the new basic feasible solution is given in the following table.

5	6	9	0
		70	30
3	5	10	0
70	5		€
6	7	6	0
		50	
6	4	10	0
	75		

We see that the above table satisfies the rim conditions with $(m + n - 1)$ non-negative allocations at independent positions. We apply MODI method.

5	3	6	5	9	0		
	2		1	70	30	$u_1 = 0$	
3	70	5	5	10	9	0	$u_2 = 0$
					1		
6	0	7	2	6	50	0	$u_3 = -3$
	6		5				
6	2	4	75	10	8	0	$u_4 = -1$
	4				2	-1	
						1	
$v_1 = 3$	$v_2 = 5$	$v_3 = 9$	$v_4 = 0$				

Since all $d_{ij} > 0$, the current solution is optimal and unique.

The optimum allocation schedule is given by $x_{13} = 70$, $x_{14} = 30$, $x_{21} = 70$, $x_{22} = 5$, $x_{24} = €$, $x_{33} = 50$, $x_{42} = 75$ and the optimum (minimum) transportation cost

$$= \text{Rs. } 9 \times 70 + 0 \times 30 + 3 \times 70 + 5 \times 5 + 0 \times € + 6 \times 50 + 4 \times 75$$

$$= \text{Rs. } 1465/-$$

Maximization case in Transportation Problems

So far we have discussed the transportation problems in which the objectives has been to minimize the total transportation cost and algorithms have been designed accordingly.

If we have a transportation problems where the objective is to maximize the total profit, first we have to convert the maximization problem into a minimization problem by multiplying

all the entries by -1 (or) by subtracting all the entries from the highest entry in the given transportation table. The modified minimization problem can be solved in the usual manner.

Assignment Problem:

Introduction

The assignment problem is a particular case of the transportation problem in which the objective is to assign a number of tasks (Jobs or origins or sources) to an equal number of facilities (machines or persons or destinations) at a minimum cost (or maximum profit).

Suppose that we have ' n ' jobs to be performed on ' m ' machines (one Job to one machine) and our objective is to assign the jobs to the machines at the minimum cost (or maximum profit) under the assumption that each machine can perform each job but with varying degree of efficiencies.

The assignment problem can be stated in the form of $m \times n$ matrix (c_{ij}) called a cost matrix (or) Effectiveness matrix where c_{ij} is the cost of assigning i^{th} machine to the j^{th} job.

	1	2	3	n
1	c_{11}	c_{12}	c_{13}	c_{1n}
2	c_{21}	c_{22}	c_{23}	c_{2n}
Machines 3	c_{31}	c_{32}	c_{33}	c_{3n}
.
.
.
.
m	c_{m1}	c_{m2}	c_{m3}	c_{mn}

Mathematical formulation of an assignment problem.

Consider an assignment problem of assigning n jobs to n machines (one job to one machine). Let c_{ij} be the unit cost of assigning i^{th} machine to the j^{th} job and

$$\text{Let } x_{ij} = \begin{cases} 1, & \text{if } j^{\text{th}} \text{ job is assigned to } i^{\text{th}} \text{ machine} \\ 0, & \text{if } j^{\text{th}} \text{ job is not assigned to } i^{\text{th}} \text{ machine} \end{cases}$$

The assignment model is then given by the following LPP

$$\text{Minimize } Z = \sum_{i=1}^n \sum_{j=1}^n c_{ij} x_{ij}$$

Subject to the constraints

$$\sum_{i=1}^n x_{ij} = 1, \quad j = 1, 2, \dots, n$$

$$\sum_{j=1}^n x_{ij} = 1, \quad i = 1, 2, \dots, n$$

and $x_{ij} = 0$ (or) 1.

Difference between the transportation problem and the assignment problem.

<i>Transportation problem</i>	<i>Assignment problem</i>
(a) Supply at any source may be any positive quantity a_i	Supply at any source (machine) will be 1 i.e., $a_i = 1$.
(b) Demand at any destination may be any positive b_j	Demand at any destination (job) will be 1 i.e., $b_j = 1$.
(c) One or more source to any Number of destinations	One source (machine) to only one destination (job).

Assignment Algorithm (or) Hungarian Method.

First check whether the number of rows is equal to the number of columns. If it is so, the assignment problem is said to be **balanced**. Then proceed to step 1. If it is not balanced, then it should be balanced before applying the algorithm.

Step 1: Subtract the smallest cost element of each row from all the elements in the row of the row of the given cost matrix. See that each row contains atleast one zero.

Step 2: Subtract the smallest cost element of each column from all the elements in the column of the resulting cost matrix obtained by step 1.

Step 3: (Assigning the zeros)

- Examine the rows successively until a row with exactly one unmarked zero is found. Make an assignment to this single unmarked zero by encircling it. Cross all other zeros in the column of this enriched zero, as these will not be considered for any future assignment. Continues in this way until all the rows have been examined.
- Examine the columns successively until a column with exactly one unmarked zero is found. Make an assignment to this single unmarked zero by encircling it and cross any other zero in its row. Continue until all the columns have been examined.

Step 4: (Apply optimal Test)

- If each row and each column contain exactly one encircled zero, then the current assignment is optimal.

- (b) If at least one row/column is without an assignment (i.e., if there is at least one row/column is without one encircled zero), then the current assignment is not optimal. Go to step 5.

Step 5: Cover all the zeros by drawing a minimum number of straight lines as follows.

- Mark (✓) the rows that do not have assignment.
- Mark (✓) the columns (not already marked) that have zeros in marked rows.
- Mark (✓) the rows (not already marked) that have assignments in marked columns.
- Repeat (b) and (c) until no more marking is required.
- Draw lines through all unmarked rows and columns. If the number of these lines is equal to the order of the matrix then it is an optimum solution otherwise not.

Step 6: Determine the smallest cost element not covered by the straight lines. Subtract this smallest cost element from all the uncovered elements and add this to all those elements which are lying in the intersection of these straight lines and do not change the remaining elements which lie on the straight lines.

Step 7: Repeat steps (1) to (6). Until an optimum assignment is attained.

Note 1: In case some rows or columns contain more than one zero, encircle any unmarked zero, encircle any unmarked zero arbitrarily and cross all other zeros in its column or row. Proceed in this way until no zero is left unmarked or encircled.

Note 2: The above assignment algorithm is only for minimization problems.

Note 3: If the given assignment problem is of maximization type, convert it to a minimization assignment problem by $\max Z = -\min (-Z)$ and multiply all the given cost elements by -1 in the cost matrix and then solve by assignment algorithm.

Note 4: Sometimes a final cost matrix contains more than required number of zeros at independent positions. This implies that there is more than one optimal solution (multiple optimal solutions) with the same optimum assignment cost.

Example 1:

Consider the problem of assigning five jobs to five persons. The assignment costs are given as follows:

		Job				
		1	2	3	4	5
From	A	8	4	2	6	1
	B	0	9	5	5	4
	C	3	8	9	2	6

D	4	3	1	0	3
E	9	5	8	9	5

Determine the optimum assignment schedule.

Solution: The cost matrix of the given assignment problem is

$$\begin{pmatrix} 8 & 4 & 2 & 6 & 1 \\ 0 & 9 & 5 & 5 & 4 \\ 3 & 8 & 9 & 2 & 6 \\ 4 & 3 & 1 & 0 & 3 \\ 9 & 5 & 8 & 9 & 5 \end{pmatrix}$$

Since the number of rows is equal to the number of columns in the cost matrix, the given assignment problem is balanced.

Step 1: Select the smallest cost element in each row and subtract this from all the elements of the corresponding row, we get the reduced matrix

$$\begin{pmatrix} 7 & 3 & 1 & 5 & 0 \\ 0 & 9 & 5 & 5 & 4 \\ 1 & 6 & 7 & 0 & 4 \\ 4 & 3 & 1 & 0 & 3 \\ 4 & 0 & 3 & 4 & 0 \end{pmatrix}$$

Step 2: select the smallest cost element in each column and subtract this from all the elements of the corresponding column, we get the reduced matrix.

$$\begin{pmatrix} 7 & 3 & 0 & 5 & 0 \\ 0 & 9 & 4 & 5 & 4 \\ 1 & 6 & 6 & 0 & 4 \\ 4 & 3 & 0 & 0 & 3 \\ 4 & 0 & 2 & 4 & 0 \end{pmatrix}$$

Since each row and each column at least one zero, we shall make assignments in the reduced matrix.

Step 3: Examine the rows successively until a row with exactly one unmarked zero is found. Since the 2nd row contains a single zero, encircle this zero and cross all other zeros of its column. The 3rd row contains exactly one unmarked zero, so encircle this zero and cross all other zeros in its column. The 4th row contains exactly one unmarked zero, so encircle this zero and cross all other zeros in its column. The 1st row contains exactly one unmarked zero, so encircle this zero and cross all other zeros in its column. Finally the last row contains exactly one unmarked zero, so encircle this zero and cross all other zeros in its column. Likewise

examine the columns successively. The assignments in rows and columns in the reduced matrix is given by

$$\begin{pmatrix} 7 & 3 & 0 & 5 & (0) \\ (0) & 9 & 4 & 5 & 4 \\ 1 & 6 & 6 & (0) & 4 \\ 4 & 3 & (0) & 0 & 3 \\ 4 & (0) & 2 & 4 & 0 \end{pmatrix}$$

Step 4: Since each row and each column contains exactly one assignment (i.e., exactly one encircled zero) the current assignment is optimal.

∴ The optimal assignment schedule is given by A → 5, B → 1, C → 4, D → 3, E → 2.

The optimum (minimum) assignment cost = (1 + 0 + 2 + 1 + 5) cost units = 9 units of cost.

Example 2:

The processing time in hours for the when allocated to the different machines are indicated below. Assign the machines for the jobs so that the total processing time is minimum.

		Machines				
		M ₁	M ₂	M ₃	M ₄	M ₅
Jobs	J ₁	9	22	58	11	19
	J ₂	43	78	72	50	63
	J ₃	41	28	91	37	45
	J ₄	74	42	27	49	39
	J ₅	36	11	57	22	25

Solution:

The cost matrix of the given problem is

9	22	58	11	19
43	78	72	50	63
41	28	91	37	45
74	42	27	49	39
36	11	57	22	25

Since the number of rows is equal to the number of columns in the cost matrix, the given assignment problem is balanced.

Step 1: select the smallest cost element in each row and subtract this from all the elements of the corresponding row, we get the reduced matrix.

$$\begin{pmatrix} 0 & 13 & 49 & 2 & 10 \\ 0 & 35 & 29 & 7 & 20 \\ 13 & 0 & 63 & 9 & 17 \\ 47 & 15 & 0 & 22 & 12 \\ 25 & 0 & 46 & 11 & 14 \end{pmatrix}$$

Step 2: Select the smallest cost element in each column and subtract this from all the elements of the corresponding column, we get the following reduced matrix.

$$\begin{pmatrix} 0 & 13 & 49 & 0 & 0 \\ 0 & 35 & 29 & 5 & 10 \\ 13 & 0 & 63 & 7 & 7 \\ 47 & 15 & 0 & 20 & 2 \\ 25 & 0 & 46 & 9 & 4 \end{pmatrix}$$

Step 3: Now we shall examine the rows successively. Second row contains a single unmarked zero, encircle this zero and cross all other zeros in its column. Third row contains a single unmarked zero, encircle this zero and cross all other zeros in its column. Fourth row contains a single unmarked zero, encircle this zero and cross all other zero in its column. After this no row is with exactly one unmarked zero. So go for columns.

Examine the columns successively. Fourth column contains exactly one unmarked zero, encircle this zero and cross all other zeros in its row. After examining all the rows and columns. We get

$$\begin{pmatrix} 0 & 13 & 49 & (0) & 0 \\ (0) & 35 & 29 & 5 & 10 \\ 13 & (0) & 63 & 7 & 7 \\ 47 & 15 & (0) & 20 & 2 \\ 25 & 0 & 46 & 9 & 4 \end{pmatrix}$$

Step 4: Since the 5th column do not have any assignment, the current assignment is not optimal.

Step 5: Cover all the zeros by drawing a minimum number of straight lines as follows:

- (a) Mark (✓) the rows that do not have assignment. The row 5 is marked.
- (b) Mark (✓) the columns (not already marked) that have zeros in marked rows. Thus column 2 is marked.
- (c) Mark the rows (not already marked) that have assignment in, marked columns. Thus row 3 is marked.
- (d) Repeat (b) and (c) until no more marking is required. In the present case this repetition is not necessary.
- (e) Draw lines through all unmarked rows (rows 1, 2 and 4). And marked columns (column 2). We get

$$\begin{pmatrix}
 0 & 13 & 49 & 0 & 0 \\
 0 & 35 & 29 & 5 & 10 \\
 13 & 0 & 63 & 7 & 7 \\
 47 & 15 & 0 & 20 & 2 \\
 25 & 0 & 46 & 9 & (4)
 \end{pmatrix}$$

Step 6: Here 4 is the smallest element not covered by these straight lines. Subtract this 4 from all the uncovered element and add this 4 to all those elements which are lying in the intersections of these straight lines and do not change the remaining elements which lie on these straight lines. We get the following matrix.

$$\begin{pmatrix}
 0 & 17 & 49 & 0 & 0 \\
 0 & 39 & 29 & 5 & 10 \\
 9 & 0 & 59 & 3 & 3 \\
 47 & 19 & 0 & 20 & 2 \\
 21 & 0 & 42 & 5 & 0
 \end{pmatrix}$$

Since each row and each column contains at least one zero, we examine the rows and columns successively, i.e., repeat step 3 above, we get

$$\begin{pmatrix}
 0 & 17 & 49 & (0) & 0 \\
 (0) & 39 & 29 & 5 & 10 \\
 9 & (0) & 59 & 3 & 3 \\
 47 & 19 & (0) & 20 & 2 \\
 21 & 0 & 42 & 5 & (0)
 \end{pmatrix}$$

In the above matrix, each row and each column contains exactly one assignment (i.e., exactly one encircled zero), therefore the current assignment is optimal.

∴ The optimum assignment schedule is $J_1 \rightarrow M_4, J_2 \rightarrow M_1, J_3 \rightarrow M_2, J_4 \rightarrow M_3,$

$J_5 \rightarrow M_5$ and the optimum (minimum) processing time

$$= 11+43+28+27+25 \text{ hours} = 134 \text{ hours.}$$

Unbalanced Assignment Models

If the number of rows is not equal to the number columns in the cost matrix of the given assignment problems, then the given assignment problems is said to be unbalanced.

First convert the unbalanced assignment problems in to a balanced one by adding dummy rows or dummy columns with zero cost element in the cost matrix depending upon whether $m < n$ or $m > n$ and then solve by the usual method.

Example 1: A company has four machines to do three jobs. Each job can be assigned to one and only one machine. The cost of each job on each machine is given in the following table.

		Machines			
		1	2	3	4
Jobs	A	18	24	28	32
	B	8	13	17	19
	C	10	15	19	22

What are job assignments which will minimize the cost?

Solution:

The cost matrix of the given assignment problems is

$$\begin{pmatrix} 18 & 24 & 28 & 32 \\ 8 & 13 & 17 & 19 \\ 10 & 15 & 19 & 22 \end{pmatrix}$$

Since the number of rows is less than the number of columns in the cost matrix, the given assignment problems is unbalanced.

To make it a balanced one, add a dummy job D (row) with zero cost elements. The balanced cost matrix is given by

$$\begin{pmatrix} 18 & 24 & 28 & 32 \\ 8 & 13 & 17 & 19 \\ 10 & 15 & 19 & 22 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

Now select the smallest cost element in each row (column) and subtract this from all the elements of the corresponding row (columns), we get the reduced matrix

$$\begin{pmatrix} 0 & 6 & 10 & 14 \\ 0 & 5 & 9 & 11 \\ 0 & 5 & 9 & 12 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

In this reduced matrix, we shall make the assignment in rows and columns having single zero. We have

$$\begin{pmatrix} (0) & 6 & 10 & 14 \\ 0 & 5 & 9 & 11 \\ 0 & 5 & 9 & 12 \\ 0 & (0) & 0 & 0 \end{pmatrix}$$

Since there are some rows and columns without assignment, the current assignment is not optimal.

Cover the all zeros by drawing a minimum number of straight lines. Choose the smallest cost element not covered by these straight lines.

$$\begin{pmatrix} 0 & 6 & 10 & 14 \\ 0 & 5 & 9 & 11 \\ 0 & (5) & 9 & 12 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

Here 5 is the smallest cost element not covered by these straight lines. Subtract this 5 from all the uncovered element, add this 5 to those elements which lie in the intersections of

these straight lines and do not change the remaining element which lie on the straight lines. We get

0	1	5	9
0	0	4	6
0	0	4	7
5	0	0	0

Since each row and each column contains atleast one zero, we shall make assignment in the rows and columns having single zero. We get

(0)	1	5	9
0	(0)	4	6
0	0	4	7
5	0	(0)	0

Since there are some rows and columns without assignment, the current assignment is not optimal.

Cover all the zeros by drawing a minimum number of straight lines.

0	1	5	9
0	0	4	6
0	0	(4)	7
5	0	0	0

Choose the smallest cost element not covered by these straight line, subtract this from all the uncovered elements, add this to those elements which are in the intersection of the lines and do not change the remaining elements which lie on these straight lines. Thus we get

0	1	1	5
0	0	0	2
0	0	0	3
9	4	0	0

Since each row and each column contains atleast one zero, we shall make the assignment in the rows and columns having single zero. We get

(0)	1	1	5
0	(0)	0	2
0	0	(0)	3
9	4	0	(0)

Since each row and each column contains exactly one assignment (i.e., exactly one encircled zero) the current assignment is optimal.

∴ The optimum assignment schedule is given by $A \rightarrow 1, B \rightarrow 2, C \rightarrow 3, D \rightarrow 4$ and the optimum (minimum) assignment cost

$$= (18+13+19+0) \text{ cost unit} = 50/- \text{ units of cost}$$

Note 1: For this problem, the alternative optimum schedule is $A \rightarrow 1, B \rightarrow 2, C \rightarrow 3, D \rightarrow 4$, with the same optimum assignment cost = Rs. $(18+17+15+0) = 50/-$ units of cost.

Note 2: Here the assignment $D \rightarrow 4$ means that the dummy Job D is assigned to the 4th Machine. It means that machine 4 is left without any assignment.

Maximization case in Assignment Problems

In an assignment problem, we may have to deal with maximization of an objective function. For example, we may have to assign persons to jobs in such a way that the total profit is maximized. The maximization problems has to be converted into an equivalent minimization problem and then solved by the usual Hungarian Method.

The conversion of the maximization problem into an equivalent minimization problems can be done by any of the following methods:

- (i) Since $\max Z = -\min (-Z)$, multiply all the cost element c_{ij} of the cost matrix by -1 .
- (ii) Subtract all the cost elements c_{ij} of the cost matrix from the highest cost element in that cost matrix.

Example:

Solve the assignment problem for maximization given the profit matrix (profit in rupees).

	Machines			
	P	Q	R	S
A	51	53	54	50
B	47	50	48	50

Jobs C	49	50	60	61
D	63	64	60	60

Solution:

The profit matrix of the given assignment problem is

$$\begin{pmatrix} 51 & 53 & 54 & 50 \\ 47 & 50 & 48 & 50 \\ 49 & 50 & 60 & 61 \\ 63 & (64) & 60 & 60 \end{pmatrix}$$

Since this is a maximization problem, it can be converted into an equivalent minimization problem by subtracting all the profit elements in the profit from the highest profit element 64 of this profit matrix. Thus the cost matrix of the equivalent minimization problem is

$$\begin{pmatrix} 13 & 11 & 10 & 14 \\ 17 & 14 & 16 & 14 \\ 15 & 14 & 4 & 3 \\ 1 & 0 & 4 & 4 \end{pmatrix}$$

Select the smallest cost in each row and subtract this from all the cost elements of the corresponding row. We get

$$\begin{pmatrix} 3 & 1 & 0 & 4 \\ 3 & 0 & 2 & 0 \\ 12 & 11 & 1 & 0 \\ 1 & 0 & 4 & 4 \end{pmatrix}$$

Select the smallest cost element in each column and subtract this from all the cost elements of the corresponding column. We get

$$\begin{pmatrix} 2 & 1 & 0 & 4 \\ 2 & 0 & 2 & 0 \\ 11 & 11 & 1 & 0 \\ 0 & 0 & 4 & 4 \end{pmatrix}$$

Since each row and each column contains atleast one zero, we shall make the assignment in rows and columns having single zero. We get

2	1	(0)	4
2	(0)	2	0
11	11	1	(0)
(0)	0	4	4

Since each row and each column contains exactly one encircled zero, the current assignment is optimal.

∴ The optimum assignment schedule is given by $A \rightarrow R$, $B \rightarrow Q$, $C \rightarrow S$, $D \rightarrow P$ and the optimum (maximum) profit

$$= \text{Rs. } (54 + 50 + 61 + 63)$$

$$= \text{Rs. } 228/-$$

POSSIBLE QUESTIONS**PART-B(5× 2 =10)**

1. What do you understand by transportation problem?
2. Define feasible solution of a transportation problem.
3. What is the optimality test used while solving an Assignment Problem using Hungarian method?
1. What is an assignment problem? Give two applications.
2. What is the optimality test used while solving an Assignment Problem using Hungarian method?
3. Define feasible solution of a transportation problem.
4. What do you understand by transportation problem?

PART-C (5× 4 =20)

1. a) Find the initial basic feasible solution for the following transportation problem by VAM.

		Distribution centres				Availability
		D_1	D_1	D_1	D_1	
Origin	S_1	11	13	17	14	250
	S_2	16	18	14	10	300
	S_3	21	24	13	10	400
Requirements		200	225	275	250	

- 2). Find the starting solution of the following transportation model

1	2	6	7
0	4	2	12
3	1	5	11
10	10	10	

Using (i) North West Corner rule
(ii).Least Cost method

- 3) Explain the procedure of Hungarian method to solve Assignment Problem.
4. Solve the following transportation problems using vogel's method.

		A	B	C	D	E	F	Available
Factory	1	9	12	9	6	9	10	5
	2	7	3	7	7	5	5	6
	3	6	5	9	11	3	11	2
	4	6	8	11	2	2	10	9
Requirement		4	4	6	2	4	2	

PART- D (1× 10 =10)

- 1) Solve the assignment problem for maximization given the profit matrix (profit in rupees).
Machines

	P	Q	R	S
A	51	53	54	50
B	47	50	48	50
Jobs C	49	50	60	61
D	63	64	60	60

2) Solve the transportation problem with unit transportation costs, demands and supplies as given below:

		Destination				
		D ₁	D ₂	D ₃	D ₄	Supply
	S ₁	6	1	9	3	70
Source	S ₂	11	5	2	8	55
	S ₃	10	12	4	7	70
	Demand	85	35	50	45	

3. Determine basic feasible solution to the following transportation problems using North West Corner Rule:

		Sink					Supply
		A	B	C	D	E	
Origin	P	2	11	10	3	7	4
	Q	1	4	7	2	1	8
	R	3	9	4	8	12	9
Demand		3	3	4	5	6	



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DEPARTMENT OF MATHEMATICS
Multiple Choice Questions (Each Question Carries One Mark)

Subject Name: REAL ANALYSIS Subject Code: 17MMF002

UNIT IV

	OPTION1	OPTION2	OPTION3	OPTION4	ANSWERS
Sequence $\{1/n\}$ is	unbounded & convergent	decreasing sequence	monotonic sequence	oscillating sequence	bounded
The sequence $\{1, 0, 1, 0, 1, \dots\}$ is	increasing sequence	unbounded	bounded below	bounded above	divergent
Every convergent sequence is	bounded	convergent	unbounded	both converges and diverges	bounded
The series $1+3+5+7+\dots$	divergent	convergent	bounded	divergent	comparison test
Cauchy sequence is	unbounded & convergent	root test	ratio test	Leibniz test	
Which one of the following test does not give absolute convergence series	comparison test	$\lim \sup (x_n + y_n) > \lim \sup x_n + \lim \sup y_n$	$\lim \sup (x_n + y_n) = \lim \sup x_n + \lim \sup y_n$	$\lim \sup (x_n + y_n) < \lim \sup x_n + \lim \sup y_n$	$\lim \sup (x_n + y_n) \leq \lim \sup x_n + \lim \sup y_n$
If $\langle x_n \rangle$ and $\langle y_n \rangle$ sequence of real number	$\lim \sup (x_n + y_n) \leq \lim \sup x_n + \lim \sup y_n$	exactly two constant sub sequence	exactly three constant sub sequence	exactly four constant sub sequence	exactly two constant sub sequence
The sequence $\langle (-1)^n \rangle$ has	exactly one constant sub sequence	bounded	having a subsequence converging to 3	convergent	bounded
Let $\langle a_n \rangle$ be least power of 2 that divides n then $\langle a_n \rangle$ is	divergent to infinity	convergent but not absolutely convergent	absolutely divergent	divergent	convergent but not absolutely convergent
A conditionally converges series is a series which is	absolutely convergent	bounded	not necessarily bounded	neither bounded nor unbounded	bounded
The set of limit points of a bounded sequence is	unbounded	sequence of rational numbers	sequence of irrational numbers	bounded sequence of rational numbers	sequence of real numbers
Cauchy sequence is convergent if it is a	sequence of real numbers	convergent sequence	bounded sequence	unbounded sequence	divergent sequence
If a sequence is not a Cauchy sequence then it is	divergent sequence	bounded	not necessarily bounded	neither bounded nor unbounded	bounded
The set of limit points of a bounded sequence is	unbounded		2	3	4
If $\{x_n\}$ and $\{x_{n+1}\}$ then the sequence $\{x_n\}$ converges to					2
Every Cauchy sequence contains	convergent subsequence	need not be convergent	may be convergent	divergent subsequence	convergent subsequence
The series $\sum_{n=1}^{\infty} \frac{1}{n!}$ is	bounded	infinite limit	unique limit	no limit	unique limit
Every convergent sequence is bounded and it has	finite limit	If it is not Riemann integrable on $[a, b]$	If it is Riemann integrable on R	If it is integrable on R	If it is Riemann integrable on $[a, b]$
If $f: [a, b] \rightarrow R$ is continuous and monotonic functions then	If it is Riemann integrable on $[a, b]$	Q	(S)	(S)	(S)
The notation of a sequence is	S	divergent	Uniformly Convergent	does not Uniformly Convergent	Uniformly Convergent
The series $\sum_{n=1}^{\infty} \frac{1}{n!}$ is	convergent	need not be measurable	measurable	may be measurable	measurable
Union of two measurable sets is	not measurable	infinity			2
Cantor ternary set is measurable and its measure is	0	uncountable	bounded	un bounded	2
A set without measure different from zero is	countable	only one subcover	finite subcover	no subcover	uncountable
A set A is said to be compact if it has a	many subcover	closed	open as well as closed	neither open nor closed	finite subcover
The empty set \emptyset and whole set X	open	one limit	many limit	no limit point	open as well as closed
Finite sets in a metric space have	more than one limit	it is set	If it is a limit point	it is empty	no limit point
A subset A of R is connected if and only if	It is an interval	closed	semi open	semi closed	it is an interval
In a metric space every singleton set $\{p\}$ is	open	$B - A$ is semi open set	$B - A$ is closed set	$B - A$ is empty set	closed
If A is open set and B is closed set then	$B - A$ is open set	unbounded	totally bounded	bounded below	$B - A$ is closed set
A sequentially compact metric space is	bounded	non negative finite number	extended real number	extended rational number	totally bounded
The total Variation on $[a, b]$ is	non positive finite number	it is not of bounded variation of $[a, b]$	it is not of bounded variation of R	it is of bounded variation of R	non negative finite number
If f is absolutely continuous on $[a, b]$	it is of bounded variation of $[a, b]$	it is always a bounded variation	its never a function of bounded variation	may be a function of bounded variation	it is of bounded variation of $[a, b]$
A continuous function is	may or may not be a function of bounded variation	almost one cluster point	atleast one cluster point	unique cluster point	may or may not be a function of bounded variation
Every infinite sequence $\{x_n\}$ in X has	more cluster point	A is complete metric space	undefined	need not be a complete metric space	atleast one cluster point
If A is closed subset of a complete metric space	A is incomplete metric space	any finite subcollection of F has empty intersection	any finite subcollection of F has empty set	any finite subcollection of F has non-empty intersection	A is incomplete metric space
A collection F of sets have finite intersection property if	A is incomplete	A is complete	complement of A is closed	complement of A is open	any finite subcollection of F has empty intersection
If A is an open subset of complete metric space X then	A is incomplete	every sequence in X is divergent	every Cauchy sequence in X is convergent	every Cauchy sequence in X is divergent	complement of A is closed
A metric space (X, p) is complete if	every sequence in X is convergent	closed set	empty set	non empty set	every Cauchy sequence in X is convergent
The union of any finite collection of non empty closed set is	open set	closed	open and closed	does not exist	closed set
The empty set \emptyset of a metric space is	open	not countable	may be countable	need not be a countable	open and closed
The set $[0, 1]$ is	countable	C is of measure zero	C is uncountable and of measure zero	C is uncountable and of positive measure	not countable
Let C be the Cantor's middle third set then	C is not countable	3	2	3	C is uncountable and of measure zero
The set of rational numbers lebesgue outer measure is	finite subcovering of F	0	no finite subcovering of F	no infinite subcovering of F	0
If F is a closed and bounded set of real numbers then each open covering is	finite subcovering of F	it is uncountable	it is dense	it is perfect set	finite subcovering of F
What is not correct about cantor ternary set	it is closed	not measurable	neither measurable nor not measurable	need not be measurable	it is dense
If f is a measurable function and $f \neq g$ almost everywhere, then g is	measurable	sum of an end points of the interval	Product of an end points of the interval	division of an end points of the interval	measurable
The length of an interval I is	difference of an end points of the interval	every interval is measurable	every open set in R is measurable	every closed set in R is measurable	difference of an end points of the interval



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Subject : REAL ANALYSIS
SUBJECT CODE: 17MMP102

SEMESTER: I
CLASS : I M.Sc

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UNIT V

Implicit functions and extremum problems:

Introduction – Functions with non zero Jacobian determinant – Inverse function theorem –
Implicit function theorem – Extrema of real valued functions of one variable and several
variables.

TEXT BOOK

1. Rudin. W., (1976) .Principles of Mathematical Analysis, Mcgraw Hill, New york .

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1. Tom .M. Apostol., (2002). Mathematical Analysis, Second edition, Narosa Publishing House, New Delhi.
2. Gupta.S.L. and Gupta.N.R.,(2003).Principles of Real Analysis, Second edition, Pearson Education Pvt.Ltd, Singapore.
3. Royden .H.L., (2002). Real Analysis, Third edition, Prentice hall of India,New Delhi.
4. Sterling. K. Berberian., (2015).A First Course in Real Analysis, Springer Pvt Ltd, New Delhi.

2. The definition of a local maximum or minimum is as follows:

Def. f has a local maximum (minimum) at a point $x^0 \in X$ if $\exists N(x^0)$ such that for all $x \in N(x^0)$, $f(x) - f(x^0) < 0$. ($f(x) - f(x^0) > 0$ for a minimum.)

3. The generalization of the first order condition is as follows:

Proposition 1. If a differentiable function f has a maximum or a minimum at $x^0 \in X$, then $f_i(x^0) = 0$, for all i .

The n equations generated by setting each partial derivative equal to zero represent the first order conditions. If a solution exists, then they may be solved for the n solution values x_i^0 .

4. As in the case of n choice variables, there are second order conditions which determine whether a critical point is a maximum or a minimum. The complication is that there is no longer one second order derivative which can be checked for negativity or positivity. In fact, there are n^2 such derivatives, $f_{ij}(x^0)$, $i, j = 1, \dots, n$. The relevant second order condition for a maximum is that

$$(*) \quad d^2f(x^0) = \sum_{i=1}^n \sum_{j=1}^n f_{ij}(x^0) dx_i dx_j < 0 \text{ for all } (dx_1, \dots, dx_n) \neq 0.$$

This condition is that the quadratic form $\sum_{i=1}^n \sum_{j=1}^n f_{ij}(x^0) dx_i dx_j$ is *negative definite*.

In (*), the discriminant is the Hessian matrix of f (the objective function). As discussed above, the rather cumbersome (*) condition is equivalent to a fairly simple sign condition. This is as follows:

$$(SOC) \text{ (max)} \quad |PM_i| \text{ of } H = \begin{bmatrix} f_{11} & \cdot & \cdot & \cdot & f_{1n} \\ \cdot & & & & \cdot \\ \cdot & & & & \cdot \\ \cdot & & & & \cdot \\ f_{n1} & \cdot & \cdot & \cdot & f_{nn} \end{bmatrix}, \text{ evaluated at } x^0, \text{ have signs } (-1)^i.$$

The analogous conditions for a minimum are that

$$(**) \quad d^2f(x^0) = \sum_{i=1}^n \sum_{j=1}^n f_{ij}(x^0) dx_i dx_j > 0 \text{ for all } (dx_1, \dots, dx_n) \neq 0,$$

meaning that $d^2f(x^0)$ is positive definite, and this condition is equivalent to

$$(SOC) \text{ (min)} \quad |PM_i| \text{ of } H = \begin{bmatrix} f_{11} & \cdot & \cdot & \cdot & f_{1n} \\ \cdot & & & & \cdot \\ \cdot & & & & \cdot \\ \cdot & & & & \cdot \\ f_{n1} & \cdot & \cdot & \cdot & f_{nn} \end{bmatrix}, \text{ evaluated at } x^0, \text{ have positive signs}$$

If f satisfies the SOC for a maximum globally, then f is strictly concave. If it satisfies the SOC for a minimum globally, then f is strictly convex. For an n variable function, the definition of strict concavity reads the same: $f(\alpha x + (1-\alpha)x') > \alpha f(x) + (1-\alpha)f(x')$, $x \neq x'$, $\alpha \in (0,1)$.

Proposition 2. If at a point x^0 we have

(i) $f_i(x^0) = 0$, for all i , and

(ii) SOC for a maximum (minimum) is satisfied at x^0 ,

Then x^o is a local maximum (minimum). If in addition the SOC is met for all $x \in X$ or if f is strictly concave (convex), then x^o is a unique global maximum (minimum).

Examples: #1 Maximizing a profit function over two strategy variables. Let profit be a function of the two variables x_i , $i=1,2$. The profit function is $\pi(x_1, x_2) = R(x_1, x_2) - \sum r_i x_i$, where r_i is the unit cost of x_i and R is revenue. We wish to characterize a profit maximal choice of x_i . The problem is written as

$$\text{Max}_{\{x_1, x_2\}} \pi(x_1, x_2).$$

The FOC are

$$\pi_1(x_1, x_2) = 0$$

$$\pi_2(x_1, x_2) = 0.$$

The second order conditions are

$$\pi_{11} < 0, \pi_{11}\pi_{22} - \pi_{12}^2 > 0 \text{ (recall Young's Theorem } \pi_{ij} = \pi_{ji}).$$

The effect of a change in r_1 can be determined by differentiating the FOC with respect to r_1 . We obtain

$$H \begin{bmatrix} \partial x_1 / \partial r_1 \\ \partial x_2 / \partial r_1 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \text{ where } H \text{ is the relevant Hessian.}$$

Using Cramer's rule,

$$\partial x_1 / \partial r_1 = \frac{\begin{vmatrix} 1 & \pi_{12} \\ 0 & \pi_{22} \end{vmatrix}}{|H|} = \pi_{22} / |H| < 0.$$

Likewise

$$\partial x_2 / \partial r_1 = \frac{\begin{vmatrix} \pi_{11} & 1 \\ \pi_{21} & 0 \end{vmatrix}}{|H|} = -\pi_{21} / |H|.$$

The sign of π_{21} is positive if 1 and 2 are complements in profit and it is negative if they are substitutes.

#2. Min $x^2 + xy + 2y^2$. The FOC are

$$2x + y = 0,$$

$$x + 4y = 0.$$

Solving for the critical values $x = 0$ and $y = 0$. $f_{11} = 2$, $f_{12} = 1$ and $f_{22} = 4$. The Hessian is

$$H = \begin{bmatrix} 2 & 1 \\ 1 & 4 \end{bmatrix}, \text{ with } f_{11} = 2 > 0 \text{ and } |H| = 8 - 1 = 7 > 0.$$

Thus, (0,0) is a minimum. Further, it is global, because the Hessian sign conditions are met for any x, y .

Existence

In the case of a function of many variables, we want to generalize our existence argument above. To do this, we must introduce a few concepts.

Def. 1. A set $X \subset \mathbb{R}^N$ is said to be *open* if for all $x \in X \exists N(x)$ such that $N(x) \subset X$. The set X is said to be *closed* if its complement is open.

Def. 2. A set $X \subset \mathbb{R}^N$ is said to be *bounded* if the distance between any two of its points is finite.

That is, $\left[\sum_{i=1}^n (x_i - x'_i)^2 \right]^{1/2} < \infty$, for all $x, x' \in X$.

Def. 3 A set $X \subset \mathbb{R}^N$ is said to be *compact* if it is both closed and bounded.

We can now state a basic existence result.

Proposition. Let $f: X \rightarrow \mathbb{R}$, where X is a subset of \mathbb{R}^N . If X is compact and f is continuous, then f has a maximum and a minimum on X . If X is both compact and convex and f is strictly concave, then f has a unique maximum on X . If X is both compact and convex and f is strictly convex, then f has a unique minimum on X .

This proposition does not distinguish between boundary optima and interior optima. As in the case of a function of single variable, the results can be used to show the existence of interior optima by showing that boundary optima are dominated. The technique is as described above.

Suppose that there are $m < n$ constraints $g_j(x) = 0$, $j = 1, \dots, m$. The Lagrangian is written as $L(\lambda_1, \dots, \lambda_m, x_1, \dots, x_n) = f(x) + \sum \lambda_j g_j(x)$. The FOC are that the derivatives of L in x_i and λ_j vanish:

$$f_i + \sum \lambda_j \partial g_j / \partial x_i = 0, \text{ for all } i,$$

$$g_j(x) = 0, \text{ for all } j.$$

The bordered Hessian becomes

$$|\overline{H}| = \begin{bmatrix} 0 & J_g \\ J_g^* & L_{ij} \end{bmatrix},$$

where J_g is the Jacobian of the constraint system in x , and $[L_{ij}]$ is the Hessian of the function L in x . The condition $m < n$ must be met, and the sign conditions for a maximum and a minimum are written in terms of the principle minors of the above bordered Hessian. For a maximum, the condition is

(SOC)(max) PM_i of $|\overline{H}|$ of order $i > 2m$ has sign $(-1)^r$, where r is the order of the largest order square $[L_{ij}]$ embedded in $|PM_i|$.

For a minimum, the condition is

(SOC)(min) $|PM_i|$ of $|\overline{H}|$ of order $i > 2m$ has sign $(-1)^m$.

Examples:

#1. Find the critical 4-tuple for the function $y = f(x_1, x_2, x_3)$, $f: \mathbb{R}_+^3 \rightarrow \mathbb{R}$

$$y = x_1 x_2 x_3$$

subject to

$$x_1^2 + x_2^2 - 1 = 0$$

$$x_1 - x_3 = 0.$$

Restrict your choice of x_i , $i = 1, 2, 3$, to the positive reals.

$$L = x_1 x_2 x_3 + \lambda_1 (x_1^2 + x_2^2 - 1) + \lambda_2 (x_1 - x_3)$$

$$(1) \quad x_2 x_3 + 2\lambda_1 x_1 + \lambda_2 = 0$$

$$(2) \quad x_1 x_3 + 2\lambda_1 x_2 = 0$$

$$(3) \quad x_1 x_2 - \lambda_2 = 0$$

$$(4) \quad x_1^2 + x_2^2 - 1 = 0$$

$$(5) \quad x_1 - x_3 = 0$$

From (5)

$$x_1 = x_3$$

From (3)

$$x_1 x_2 = x_3 x_2 = \lambda_2$$

Solve for λ_1 from 2

$$x_1 x_3 + 2\lambda_1 x_2 = 0$$

$$x_1^2 = -2\lambda_1 x_2$$

$$\therefore \lambda_1 = \frac{-x_1^2}{2x_2}$$

Go to (1)

$$x_1 x_2 + 2 \left(\frac{-x_1^2}{2x_2} \right) x_1 + x_1 x_2 = 0$$

$$2x_1 x_2 - \frac{x_1^3}{x_2} = 0$$

Multiply by x_2/x_1 (both sides)

$$(\alpha) \quad 2x_2^2 - x_1^2 = 0$$

Use (4)

$$(\beta) \quad x_1^2 + x_2^2 - 1 = 0$$

Solve x_1, x_2 :

$$(\alpha) \quad x_1^2 = 2x_2^2$$

$$\therefore \quad 2x_2^2 + x_2^2 = 1$$

$$3x_2^2 = 1$$

$$x_2^2 = 1/3$$

$$x_2 = \sqrt{1/3}$$

but

$$x_1^2 = 2x_2^2 = 2(1/3) = 2/3$$

$$x_1 = \sqrt{2/3} \quad x_3 = \sqrt{2/3}$$

and

$$y^0 = (1/3)^{1/2}(2/3)^{1/2}(2/3)^{1/2}$$

$$y^0 = \frac{2}{3\sqrt{3}}$$

$$\therefore \quad (\sqrt{2/3}, \sqrt{1/3}, \sqrt{2/3}, \frac{2}{3}\sqrt{3}).$$

The Problem's Bordered Hessian.

$$[H] = \begin{bmatrix} 0 & 0 & 2x_1 & 2x_2 & 0 \\ 0 & 0 & 1 & 0 & -1 \\ 2x_1 & 1 & 2\lambda_1 & x_3 & x_2 \\ 2x_2 & 0 & x_3 & 2\lambda_1 & x_1 \\ 0 & -1 & x_2 & x_1 & 0 \end{bmatrix}$$

Remark 2: The FOC (1) and (2) are necessary conditions only if the Constraint Qualification holds. This rules out particular irregularities by imposing restrictions on the boundary of the feasible set. These irregularities would invalidate the FOC (1) and (2) should the solution occur there. Let x^o be the point at which (1) and (2) hold and let index set $k = 1, \dots, K$ represent the set of g_j which are satisfied with equality at x^o . Then the matrix

$$J = \begin{vmatrix} \frac{\partial g_1(x^o)}{\partial x_1} & \cdot & \cdot & \cdot & \frac{\partial g_1(x^o)}{\partial x_n} \\ \cdot & \cdot & & & \cdot \\ \cdot & & \cdot & & \cdot \\ \cdot & & & \cdot & \cdot \\ \frac{\partial g_K(x^o)}{\partial x_1} & & & & \frac{\partial g_K(x^o)}{\partial x_n} \end{vmatrix}$$

Has rank $K \leq n$. That is the gradient vectors of the set of equality constraints are linearly independent.

Proof. 1. Choose $\delta_1 > 0$ and $\epsilon_1 > 0$ so that $F_y(x, y) > 0$ for $\|x - a\| < \delta_1$, $|y - b| < \epsilon_1$. Since $F(a, b) = 0$ and $F(a, y)$ is strictly increasing in y , $F(a, b + \epsilon_1/2) > 0$ and $F(a, b - \epsilon_1/2) < 0$. Let $\epsilon = \epsilon_1/2$ and choose $\delta < \delta_1$ so that $F(x, b + \epsilon) > 0$ and $F(x, b - \epsilon) < 0$ if $\|x - a\| < \delta$. These dimensions define B . For fixed x with $\|x - a\| < \delta$, since $F(x, b - \epsilon) < 0$, $F(x, b + \epsilon) > 0$, and $F(x, y)$ is strictly increasing in y , the intermediate value theorem implies that there is a unique y with $|y - b| < \epsilon$ such that $F(x, y) = 0$. The uniquely determined y defines a function $f(x)$. This proves the first statement.

2. We prove that f is continuous at a . Let $e > 0$ be given. Assume that $e < \epsilon$. Then by the proof of the first statement, there is a $d > 0$ (we may choose $d < \delta$) so that the uniquely defined $f(x)$ in $\{\|x - a\| < d\}$ satisfies $|f(x) - b| < d$. This proves continuity at a . We can repeat this argument at any point $(a_1, f(a_1)) \in B$, proving that f is continuous on $\{\|x - a\| < \epsilon\}$.

3. By differentiability

$$\begin{aligned} 0 = F(x, f(x)) &= F(a, b) + \sum_j P_j(x, f(x))(x - a_j) + Q(x, f(x))(f(x) - f(a)) \\ &= \sum_j P_j(x, f(x))(x - a_j) + Q(x, f(x))(f(x) - f(a)), \end{aligned}$$

where $P_j(x, f(x)), Q(x, f(x))$ are continuous at a . Rewrite this as

$$Q(x, f(x))(f(x) - f(a)) = - \sum_j P_j(x, f(x))(x - a_j).$$

Since $Q(x, f(x))$ is continuous at a and $Q(a, f(a)) = f_y(a, b) > 0$, $Q(x, f(x)) > 0$ for x near a and we can divide by it to get

$$f(x) = f(a) + - \sum_j \frac{P_j(x, f(x))}{Q(x, f(x))}(x - a_j).$$

Each term $\frac{P_j(x, f(x))}{Q(x, f(x))}$ is continuous at a so f is differentiable at a . Moreover

$$f_{x_j}(a, b) = - \frac{F_j(a, b)}{F_y(a, b)}.$$

You might like this bad notation:

$$\frac{\partial y}{\partial x_j} = - \frac{\partial F_{x_j}}{\partial F_y}.$$

□

The geometrical significance of the Jacobian determinant is outlined here. Consider a transformation of a single rectangular Cartesian coordinate x to a new coordinate ξ . The line element dx is transformed to the new coordinate via

$$dx = \frac{dx}{d\xi} d\xi$$

In this case, the Jacobian determinant is simply the derivative $\frac{dx}{d\xi}$.

Now, consider an area element $dx dy$. For convenience in later generalization, we label the coordinates (x_1, x_2) . Therefore, $x_1 = x$, and $x_2 = y$. Let us make a transformation to a new set of coordinates (ξ_1, ξ_2) . The area element transforms as follows.

$$dx_1 dx_2 = \left| \frac{\partial x_1}{\partial \xi_1} \frac{\partial x_2}{\partial \xi_2} - \frac{\partial x_2}{\partial \xi_1} \frac{\partial x_1}{\partial \xi_2} \right| d\xi_1 d\xi_2$$

How did we obtain the above result? First, consider a differential change in the new variable ξ_1 while keeping the variable ξ_2 fixed. The components of the infinitesimal vector resulting from this change are

$$\left(\frac{\partial x_1}{\partial \xi_1} d\xi_1, \frac{\partial x_2}{\partial \xi_1} d\xi_1 \right)$$

In a like manner, we can write the components of the vector obtained by making a differential change in the second variable ξ_2 while keeping the variable ξ_1 fixed.

$$\left(\frac{\partial x_1}{\partial \xi_2} d\xi_2, \frac{\partial x_2}{\partial \xi_2} d\xi_2 \right)$$

These two vectors need not be orthogonal in general. Therefore, we need a result for the area of a parallelogram whose sides are the differential vectors written above. We know that this is the magnitude of the vector product (cross product) of the two vectors. This is the result given above for the area element.

POSSIBLE QUESTIONS**UNIT-V****PART-B ($5 \times 6 = 30$)**

1. State and prove implicit function theorem.
2. Find implicit function defined by relation $x + y + z - xyz = 0$ near point $(0,0,0)$ and find derivative with respect to x and y using implicit function theorem.
3. Prove that functions $u = x + y - z$, $v = x - y + z$ and $w = x^2 + y^2 + z^2 - 2yz$ are not independent of one another. Also find relation between them.
4. State and prove functions with non Zero Jacobian determinant.
5. Define saddle point with example and Define Jacobian determinant.
6. State and prove Inverse Function Theorem.
7. prove that if $f = u + iv$ is a complex valued function with derivative at a point z in \mathbb{C} , then $J_f(z) = |f'(z)|^2$.

PART- C ($1 \times 10 = 10$)

1. Let A be an open subset of \mathbb{R}^n and assume that $f: A \rightarrow \mathbb{R}^n$ has continuous partial derivatives on A . If $J_f(x) \neq 0$ for every x in A , then prove f is an open mapping.
2. Examine the following functions extreme values
 - (i) $f(x, y) = y^2 + 4xy + 3x^2 + x^3$
 - (ii) $f(x, y, z) = 2xy^2 + 3y^2 + 4z^2 - 3xy + 8z$
3. Prove: For some integer $n \geq 1$, Let f have continuous n^{th} derivative in the open interval (a, b) . Suppose for some interior point c in (a, b) , we have $f'(c), f''(c), \dots, f^{(n-1)}(c) = 0$ such that $f^{(n)}(c) \neq 0$. Then for n even, f has a local minimum at c if $f^{(n)}(c) > 0$ & a local maximum at c if $f^{(n)}(c) < 0$. If n is odd, there is neither a local minimum nor a local maximum at c .
4. State and prove Inverse function Theorem



KARPAGAM ACADEMY OF HIGHER EDUCATION
(Deemed to be University Established Under Section 2 of UGC Act 1956)
Pollachi Main Road, Echamart (Po).

Coinabition: -641 02

DEPARTMENT OF MATHEMATICS

Multiple Choice Questions (Each Question Carries One Mark)

Subject Name: REAL ANALYSIS	Subject Code: 17ANMP02						
UNIT-V							
The outer measure of its interval is its	value	length	size of the interval	modulus of its value	length		
Every increasing function of its	bounded variation	variational value	bounded	unbounded variation	bounded variation		
$v(x)$ is monotonic-----function	increasing	decreasing	either increasing or decreasing	neither increasing nor decreasing	increasing		
A function f of bounded variation is the expressible as a	difference of two monotonic increasing functions	sum of two monotonic increasing functions	division of two monotonic increasing functions	product of two monotonic increasing functions	difference of two monotonic increasing functions		
Outer Lebesgue measure is also known as	Lebesgue exterior measure	Lebesgue measure	Lebesgue interior measure				
If set A is Lebesgue measurable and $m^*(A \cap B) =$ -----	0	1	2	3	0		
The sets S_1, S_2, \dots, S_n are called the ----- of the partition P	components	parts	partition	combined	components		
The refinement of P is denoted by	p_1	p^*	p^*	p^*			
Every singleton set is ----- set	disconnected set	connected	measurable	unmeasurable	connected		
Intersection of finite number of open set is	closed	open	neither open nor not closed	semiclosed	open		
Union of finite number of closed set is	closed	open	neither open nor not closed	semi closed	closed		
Every closed interval is	closed	compact	open as well as closed	not compact	compact		
In a metric space (X, d) a non-empty X is	closed	compact	open as well as closed	not compact	compact		
Every infinite set A has a -----	no limit point	neighbourhood	limit point	need not be a limit point	limit point		
If f is a continuous mapping of a compact metric space X into M space y	f is uniformly continuous	not uniformly continuous	continuous	discontinuous	f is uniformly continuous		
If E_1 and E_2 are Lebesgue measurable set then	E_1 union E_2 is also Lebesgue measurable sets	$E_1 \cap E_2$	$E_1 \cup E_2$	$E_1 \cap E_2$	E_1 union E_2 is also Lebesgue measurable sets		
Let $\{F_\alpha\}$ be an open covering of A , then	there exist countable collection of F which covers A	there exist uncountable collection of F which covers A	there exist collection of F which covers A	there exist Uncountable collection of F which covers A	there exist countable collection of F which covers A		
If f is continuous Real-valued function of Compact metric space then	f is bounded	f is unbounded	f is constant	f is a function	f is bounded		
If f is an open covering of a closed and bounded set A then	there exist a finite subcollection of F which covers A	a finite subcollection of F which covers A	there exist a infinite subcollection of F which covers A	there exist a finite collection of F which covers A	there exist a finite subcollection of F which covers A		
Any Countable set of points on the real axis line	no measure	measure is 1	measure zero	not measurable	measure zero		
$\int_0^1 x dx$	$\frac{1}{2}$	$\frac{1}{2}$	4	2	$\frac{1}{2}$		
A real valued function g is ----- function if f is Lebesgue measurable	single	single	major	single			
Every open and closed set is ----- measurable	zero	Lebesgue	not measurable	need not be measurable	Lebesgue		
Every ----- set is Lebesgue measurable	carior set	empty set	Heine borel	non-empty set	Heine borel		
Which one of the following is true ?	Family M of Lebesgue measurable sets is an algebra of sets	Family M of Lebesgue measurable sets is not algebra of sets	Family M of Lebesgue measurable sets	Lebesgue measurable sets	Family M of Lebesgue measurable sets is an algebra of sets		
If Lebesgue outer measure of a set $E, m^*(E) = 0$ then	E is measurable	E is not measurable	E is need not be measurable	neither measurable nor not measurable	E is measurable		
Which one of the following is true ?	Lebesgue exterior measure m^* 's translation is invariant	Lebesgue exterior measure m^* 's translation is not invariant	need not be measurable	Lebesgue interior measure m^* 's translation is invariant	Lebesgue exterior measure m^* 's translation is invariant		
If A is countable, then	$m^*(A) = 0$	$m^*(A) = \infty$	$m^*(A) = 1$	$m^*(A) = 0$	$m^*(A) = 0$		
If $\{a_n\} \rightarrow \infty$ monotonic then	f is of bounded variation	f is unbounded	the set of discontinuity of f are uncountable	the set of continuity of f are uncountable	f is of bounded variation		
Every sequence x_n in a metric space X is convergent then every Cauchy's sequence	convergent	divergent	constant	convergent as well as divergent	convergent		
If function f is the difference of two monotonic real valued functions on $[a, b]$ then	f is of bounded variation	f is of not bounded variation	f is of need not bounded variation	f is of may be bounded variation	f is of bounded variation		
If x is an accumulation point of $S \subset \mathbb{R}$	every open ball $B(x)$ contains x	every open ball $B(x)$ contains finitely many points	every open ball $B(x)$ does not contain many of points	every open ball $B(x)$ contains few points	every open ball $B(x)$ contains x		
If a set $S \subset \mathbb{R}^n$ contains all its adherent points then S is	closed	open	neither open nor not closed	not closed	closed		
If $R = \{x \in \mathbb{R}^n : x \text{ is open then } S \subset R\}$ is	open	closed	not open	not closed	closed		
Let X is metric space if X is sequentially compact then X is	unbounded	not compact	compact	bounded	compact		
A metric space X has the Borel-Lebesgue property, if X is	sequentially compact	not compact	not compact	compact	sequentially compact		
If we take $g(x) = x$ and $f(x) = x^2$ in general mean value theorem we obtain	Lagrange's mean value theorem	Cauchy's mean value theorem	Rolle's theorem	Taylor's theorem	Cauchy's mean value theorem		
Any closed interval with usual metric is	compact	not compact	need not be compact	sequentially compact	not compact		
The Euclidean Space \mathbb{R}^n is	not separable	separable	connected	disconnected	separable		
Every dense subset is ----- in \mathbb{R}	countable	uncountable	bounded	unbounded	countable		
The usual metric space (\mathbb{R}, d) is	separable	not separable	not compact	compact	compact		
The set of rational numbers Lebesgue outer measure is π	1	0	3	4	0		
Every measurable set is nearly a finite union of	set	open set	closed	intervals	intervals		
Every convergent sequence of measurable functions is nearly	convergence	divergence	convergence	absolutely convergence	uniformly convergence		

Reg No : -----
[17MBAP105]

KARPAGAM UNIVERSITY
Karpagam Academy of Higher Education
COIMBATORE – 641021
First Semester

II Internal Test Sep'17

Quantitative Methods For Management

Class: I MBA

Time :2 Hours

Date : 27.10.17

Maximum: 50 Marks

PART – A (15 x 1 = 15 marks)

ANSWER ALL THE QUESTIONS

1. Two cards are drawn from a pack of 52 cards. Find the probability that both are red cards-----
a) $26C_2$ b) $52C_4$ c) $52C_2$ d) $26C_3$
2. The mean of Binomial distribution is measured by
a) np b) npq c) pq d) nq.
3. Optimization means-----
a) Maximization of profit or minimization of cost
b) minimization
c) Maximization
d) cost
4. An assignment problem is a particular case of -----
a) Transportation problem b) LPP
c) Network problem d) Integer programming problem
5. Assignment problem is a -----form of a transportation problem.
a) non-degenerate b) feasible
c) degenerate d) infeasible
6. In Birth–death model, the probability distribution of queue length is given by -----.
a) $\rho^n / (1-\rho)$ b) $\rho^2 / (1-\rho)$ c) $\rho / (1-\rho)$ d) $(1-\rho) / \rho^n$
7. The probability of an empty system is given by -----.
a) $1 - (\lambda / \mu)$ b) $\lambda / (\mu - \lambda)$ c) $\lambda / \mu (\mu - \lambda)$ d) λ / μ
8. A _____ is a decision of the player to always select the same strategy.
a) competitor b) Mixed Strategy
c) pure Strategy d) strategy
9. The variance of a binomial distribution is measured by
a) np b) np(1 - p) c) Pq d) Nq
10. A BFS for transportation problem must have exactly----- non-negative allocation
a) m+n b) m+n-2 c) m+n-1 d) m+n+1
11. The arriving people in a queueing system are called -----
a) Input b) servers c) customers d) queue
12. A loss is considered as a _____ gain.
a) Positive b) negative c) finite d) infinite
13. If binomial distribution is symmetrical if p=q=?
a) 1 b) 0.4 c) 0 d) 0.5
14. Assignment technique is essentially a ----- technique
a) Maximization b) minimization
c) minimization or maximization d) none
15. The Birth–death model is called -----.
a) M / M / 1 b) M / M / N
c) M / M / ∞ d) M / M / 2

PART-B (3 x 8 =24 Marks)

ANSWER ALL THE QUESTIONS

16. a) Three coins are tossed simultaneously. Find the probability that

- i) no head ii) one head iii) two heads

(OR)

b) Two persons A and B appeared for an interview for a job. The probability of selection of A is $\frac{1}{3}$ and that of B is $\frac{1}{2}$. Find the probability that

- (i) both of them will be selected
(ii) only one of them will be selected

17.a) Solve the transportation problem with unit transportation costs, demands and supplies as given below:

Destination

	D ₁	D ₂	D ₃	D ₄	Supply
S ₁	6	1	9	3	70
Source S ₂	11	5	2	8	55
S ₃	10	12	4	7	70
Demand	85	35	50	45	

(OR)

b) Solve the assignment problem for maximization given the profit matrix (profit in rupees).

		Machines			
		P	Q	R	S
Jobs	A	51	53	54	50
	B	47	50	48	50
	C	49	50	60	61
	D	63	64	60	60

18.a) Find the starting solution of the following transportation model

1	2	6	7
0	4	2	12
3	1	5	11
10	10	10	

Using (i) North West Corner rule
(ii) Least Cost method

(OR)

b) What is the application of the Theory of Game in the contemporary business environment? Explain.

PART- C (1 X 11 = 11 Marks)

COMPULSORY

- 19.) A box containing 100 transistors, 20 of which are Defective, 10 are selected for inspection. Indicate what is the probability that (a) all 10 are defective
(b) all 10 are good (c) at least one is defective
(d) at most 3 are defective
(i) Write the properties of Binomial Distribution.
(ii) Comment on the following "The mean of a binomial distribution is 5 and its variance is 9."

Reg. No -----
(17MMP102)

KARPAGAM UNIVERSITY
KARPAGAM ACADEMY OF HIGHER EDUCATION
COIMBATORE-21
DEPARTMENT OF MATHEMATICS

First Semester

I Internal Test – Aug'17

REAL ANALYSIS

Class : I M.Sc (MATHEMATICS) Time:2 hours
Date : 29.08.17 (FN) Max Marks: 50 Marks

PART – A (20 x 1 = 20 Marks)

ANSWER ALL THE QUESTIONS:

- If P is a partition of closed interval, then $\|P\|$ is the length of
(a) largest subinterval (b) smallest subinterval
(c) last subinterval (d) first subinterval
- The least upper bound is called -----
(a) bounded above (b) bounded below
(c) Supremum (d) infimum
- A continuous mapping r of $[a,b]$ is said to be closed if.....
(a) $r(a)=r(b)$ (b) $r(a)<r(b)$ (c) $r(a)>r(b)$ (d) $r(a)=r(b)=0$
- A function which is Riemann Steiltjes integrable w.r.t S can be denoted as
(a) $f \in R(S)$ (b) $f \notin R(S)$ (c) $f \in R$ (d) $f \notin R$
- A set S is closed if and only if $S =$ -----
(a). int S (b) closure of S (c) limit (d) 0
- The set N of natural numbers is -----
(a) bounded (b) not bounded (c) countable (d) uncountable
- If f is Riemann Steiltjes integrable w.r.t S , then its is also similar
(a) $|f|$ (b) f' (c) $|f|/f$ (d) $f/|f|$
- If $f: [a,b] \rightarrow \mathbb{R}$ is continuous and monotonic functions then
(a) f is Riemann integrable (b) Infinite limit
(c) Unique limit (d) No limit
- The series $1+3+5+7+\dots$ -----
(a) divergent (b) Convergent
(c) Bounded (d) conditionally convergent
- The series $1+r+r^2+r^3+\dots$ is oscillatory if
(a) $r=1$ (b) $r=0$ (c) $r=-1$ (d) $r=2$
- If $S = [0, 1)$ then $\sup S =$ -----
(a) $(0, 1)$ (b) $[0,1]$ (c) 0 (d) 1
- If f and g are Riemann Steiltjes integrable w.r.t S , then..... is also.
(a) $f * g$ (b) $f + g/g$
(c) $f - g/g$ (d) f/g
- If a sequence $\{a_n\}$ converges to a real number then the given sequence is-----
(a) unbounded sequence (b) convergent
(c) divergent & bounded (d) bounded
- is a range of a function from $N \times N$ to \mathbb{R} .
(a) Double sequence (b) Double series
(c) finite sequence (d) finite series
- The n^{th} partial sum of series $\sum f_n$ is given by-----
(a) $f_1 - f_2 + \dots + f_n$
(b) $f_1 - f_2 - \dots - f_n$
(c) $f_1 + f_2 - \dots - f_n$
(d) $f_1 + f_2 + \dots + f_n$

16. If an unbounded sequence -----

- (a) has a limit point
- (b) does not have a limit point
- (c) has a unique limit point
- (d) may or may not have a limit point

17. Double sequence of real numbers can be denoted by.....

- (a) $(x_{m,n})$ (b) $\sum x_{m,n}$ (c) $\sum x_m$ (d) (x_m)

18. The sequence $\{1/n\}$ is -----

- (a) convergent (b) increasing sequence
- (c) monotonic sequence (d) oscillating sequence

19. The set of limit points of a bounded sequence is

- (a) unbounded (b) not necessarily bounded
- (c) bounded (d) neither bounded nor unbounded

20. Double series of real numbers depends on -----

- (a) a single parameter (b) 2 parameters
- (c) 3 parameters (d) 4 parameters

PART – B (3 x 2 = 6 Marks)

ANSWER ALL THE QUESTIONS

21. Show that Riemann Stieltjes integral can be reduced to a finite sum.

22. State the Root test for convergence of series

23. State and prove linearity property in Riemann Stieltjes integrals.

PART – C (3x8 = 24 Marks)

ANSWER ALL THE QUESTIONS

24. a) If $f \in R(\alpha)$ on $[a, b]$ then $\alpha \in R(f)$ on $[a, b]$, prove that

$$\int_a^b f(x) d\alpha(x) + \int_a^b \alpha(x) df(x) = f(b)\alpha(b) - f(a)\alpha(a)$$

(OR)

b) For any $f \in R(\alpha)$ on $[a, b]$ and $g \in R(\alpha)$ on $[a, b]$ then $c_1 f + c_2 g \in R(\alpha)$.

$$\text{we have } \int_a^b (c_1 f + c_2 g) d\alpha = c_1 \int_a^b f d\alpha + c_2 \int_a^b g d\alpha.$$

25. a) State and prove Riemann – Stieltjes condition.

(OR)

b) Assume that α is increasing on $[a, b]$ then prove that the following are equivalent

- (i) $f \in R(\alpha)$ on $[a, b]$
- (ii) f satisfies Riemann condition w.r.to α on $[a, b]$
- (iii) $I_-(f, \alpha) = I_+(f, \alpha)$

(OR)

26. a) Let $\sum a_n$ be an absolutely convergent series f having sum S . then every rearrangement of $\sum a_n$ also converges absolutely f has sum S .

(OR)

b) State and prove Ratio Test Theorem.

Reg. No.....

[15MMP102]

KARPAGAM UNIVERSITY

Karpagam Academy of Higher Education
(Established Under Section 3 of UGC Act 1956)
COIMBATORE – 641 021
(For the candidates admitted from 2015 onwards)

M.Sc., DEGREE EXAMINATION, NOVEMBER 2015

First Semester

MATHEMATICS

REAL ANALYSIS

Time: 3 hours

Maximum : 60 marks

PART – A (20 x 1 = 20 Marks) (30 Minutes)
(Question Nos. 1 to 20 Online Examinations)

(Part - B & C 2 ½ Hours)

PART B (5 x 6 = 30 Marks)
Answer ALL the Questions

21. a) If $f \in R(\alpha)$ on $[a, b]$ and $f \in R(\beta)$ on $[a, b]$ then $f \in R(c_1 \alpha + c_2 \beta)$ on $[a, b]$ we have $\int_a^b f d(c_1 \alpha + c_2 \beta) = c_1 \int_a^b f d\alpha + c_2 \int_a^b f d\beta$.
Or
b) State and prove change of variable in Riemann – Stieltjes integral.
22. a) State and prove Merten's Theorem
Or
b) State and prove Ratio Test Theorem.
23. a) State and prove Cauchy's condition for Uniform convergence.
Or
b) Assume that $\sum f_n(x) = f(x)$ (uniformly on S) if each f_n is continuous at a point x_0 of S then f is also continuous at x_0 .
24. a) State and prove Lesgue dominated convergence Theorem
Or
b) Assume f is Riemann integrable on $[a, b] \forall b \geq a$ and assume there is a positive constant M such that $\int_a^b |f(x)| dx \leq M \forall b \geq a$.

25. a) Let A be an open subset of \mathbb{R}^n and assume that $f: A \rightarrow \mathbb{R}^n$ has continuous partial Derivatives $D_j f_j$ on A . If f is 1-1 on A and if $J_f(x) \neq 0 \forall x \in A$, then $f(A)$ is an open.
Or

- b) State and prove functions with non Zero Jacobian determinant.

PART C (1 x 10 = 10 Marks)
(Compulsory)

26. State and prove Rearrangement Theorem for double sequence.

Reg. No.....

[14MMP102]

KARPAGAM UNIVERSITY

(Under Section 3 of UGC Act 1956)

COIMBATORE - 641 021

(For the candidates admitted from 2014 onwards)

M.Sc. DEGREE EXAMINATION, NOVEMBER 2014

First Semester

MATHEMATICS

REAL ANALYSIS

Time: 3 hours

Maximum : 60 marks

PART - A (10 x 2 = 20 Marks)

Answer any TEN Questions

1. What do you mean by Partition?
2. Write about refinement.
3. Define Riemann stieljes sum.
4. Define rearrangement of a series.
5. Define double series.
6. What do you mean by Cesaro summability?
7. What do you say about limit function?
8. When a series is said to be converges uniformly?
9. Define mean convergence.
10. Define Lebesgue intergrable.
11. When we say that a function f is measurable?
12. State Lebesgue bounded convergence theorem.
13. Write a brief note on Jacobian determinant.
14. Define open mapping.
15. Define quadratic form.

PART B (5 X 8= 40 Marks)

Answer ALL the Questions

16. a) If $f \in R(\alpha)$ on $[a, b]$, then $\alpha \in R(f)$ on $[a, b]$ and prove that

$$\int_a^b f(x) d\alpha(x) + \int_a^b \alpha(x) df(x) = f(b)\alpha(b) - f(a)\alpha(a)$$

Or

1

- b) Let α be increasing on $[a, b]$. Then show the following

i) If P' is finer than P , prove that

$$U(P, f, \alpha) \leq L(P, f, \alpha) \text{ and } L(P', f, \alpha) \geq L(P, f, \alpha)$$

ii) For any two partitions P_1 and P_2 , prove that $L(P_1, f, \alpha) \leq U(P_2, f, \alpha)$

17. a) State and prove Riemann's theorem on conditionally convergent series.

Or

b) If $\sum a_n$ converges absolutely, then show that the subseries $\sum b_n$ also

converges absolutely. Also, $|\sum_{n=1}^{\infty} b_n| \leq \sum_{n=1}^{\infty} b_n \leq \sum_{n=1}^{\infty} a_n$

18. a) State and prove Dirichlet's test for uniform convergence

Or

b) Assume that $\sum f_n(x) = f(x)$ (uniformly on S) if each f_n is continuous at a point x_0 of S , then show that f is also continuous at x_0 .

19. a) Assume $f \in L(I)$ and $g \in L(I)$. Then show that

i. $(af + bg) \in L(I)$ for ever real a and b , and $\int (af + bg) = a \int f + b \int g$

ii. $\int f \geq 0$ if $f(x) \geq 0$ a.e. on I .

iii. $\int f \geq \int g$ if $f(x) \geq g(x)$ a.e. on I .

iv. $\int f = \int g$ if $f(x) = g(x)$ a.e. on I .

Or

b) State and prove Levi theorem for sequences of Lebesgue-integrable functions.

20. Compulsory : -

State and prove implicit function theorem.

2

Reg. No.....

[12MCP102]

KARPAGAM UNIVERSITY

(Under Section 3 of UGC Act 1956)

COIMBATORE - 641 021

(For the candidates admitted from 2012 onwards)

M.Sc. DEGREE EXAMINATION, NOVEMBER 2012

First Semester

MATHEMATICS

REAL ANALYSIS

Time: 3 hours

Maximum : 100 marks

PART - A (15 x 2 = 30 Marks)

Answer ALL the Questions

1. Write about the family of all functions of bounded variation?
2. What do you mean by compact interval?
3. What is called Partition of $[a, b]$?
4. Describe about limit superior?
5. Define about Strictly finer?
6. Write about lower stieltes integral?
7. Describe about Infinite series?
8. Explain about Supremum?
9. Describe about sum of the series?
10. Describe about Converge pointwise?
11. Explain about Converges uniformly on sequence of functions?
12. Mention about power series?
13. Define about Directional derivative?
14. Describe about Gradient vector?
15. Mention about Chain rule?

PART B (5 X 14= 70 Marks)

Answer ALL the Questions

16. a) i) Let $f: [a, b] \rightarrow \mathbb{R}^n$ and $g: [c, d] \rightarrow \mathbb{R}^n$ be two paths in \mathbb{R}^n , each of which is one to one on its domain. Then prove that f and g are equivalent iff they have same graph.
 - ii) If f is an increasing function on $[a, b]$ and $c \in (a, b)$, then both $f(c-)$ and $f(c+)$ exist and satisfy the inequalities $f(c-) \leq f(c) \leq f(c+)$. Then $f(a+)$ and $f(b-)$ exist and $f(a) \leq f(a+)$ and $f(b-) \leq f(b)$.
- (Or)

- b) Let f be of bounded variation on $[a, b]$. If $x \in [a, b]$, let $V(x) = V_f(a, x)$ and put $V(a) = 0$. Then show that every point of continuity of f is also a point of continuity of V and Converse is also true.

17. a) i) Let α be of bounded variation on $[a, b]$ and assume that $f \in R(\alpha)$ on $[a, b]$. then $f \in R(\alpha)$ On every subinterval $[c, d]$ of $[a, b]$.
 - ii) Assume that α be an increasing sequence on $[a, b]$, if $f \in R(\alpha)$ on $[a, b]$. If $g \in R(\alpha)$ on $[a, b]$ then the product $f \cdot g \in R(\alpha)$ on $[a, b]$.
- (Or)

- b) State and prove Riemann conditions.

18. a) The infinite product $\prod U_n$ converges iff for all $\epsilon > 0$ there exist an integer N , Such that for all $n > N \Rightarrow |U_{n+1} U_{n+2} \dots U_{n+k}| < |U_{n+1}| |U_{n+2}| \dots |U_{n+k}|$.
- (Or)

- b) State and prove Merten's theorem.

19. a) Let α be of bounded variation on $[a, b]$. Assume that each term of the sequence $\{f_n\}$ is a real valued function such that $f_n \in R(\alpha)$ on $[a, b]$ for each $n = 1, 2, \dots$. Assume That $f_n \rightarrow f$ uniformly on $[a, b]$ and define $g_n(x) = \int_a^x f_n(t) d\alpha(t)$.

- i) $f \in R(\alpha)$ on $[a, b]$.

- ii) $g_n \rightarrow g$ uniformly on $[a, b]$ where $g(x) = \int_a^x f(t) d\alpha(t)$.

(Or)

- b) i) State and prove Cauchy condition for uniform convergence.

- ii) Assume that $f_n \rightarrow f$ uniformly on S . If each f_n is continuous at a point c of S , then show that the limit function f is also continuous at c .

20. a) i) Let A be an open subset of \mathbb{R}^n and assume that $f: A \rightarrow \mathbb{R}^n$ is continuous and has finite partial derivatives $D_j f_i$ on A . If f is one to one on A and if $J_f(x) \neq 0$ for each x in A , then prove that $f(A)$ is open.

- ii) Assume that $f = (f_1, f_2, \dots, f_n)$ has continuous partial derivatives $D_j f_i$ on an open set S in \mathbb{R}^n , nad that the Jacobian determinant $J_f(a) \neq 0$ for some point a in S . Then there is an n -ball $B(a)$ on which f is one- to-one.

(Or)

- b) State and prove second derivative test for extrema.